

Free University of Bozen-Bolzano – Faculty of Computer Science  
 Master of Science in Computer Science  
 Theory of Computing – A.A. 2004/2005  
 Midterm exam – 24/11/2004

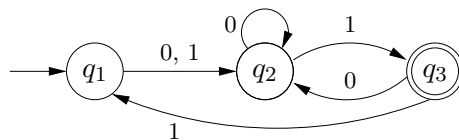
*Duration: 90 minutes*

**Problem 1** [6 points] Decide which of the following statements is TRUE and which is FALSE. You must give a brief explanation of your answer to receive full credit.

- (a) For all languages  $L_1$  and  $L_2$ , it holds that  $L_1^* \cup L_2^* = (L_1 \cup L_2)^*$ .
- (b) If  $L_1$  is regular and  $L_2$  is non-regular, then  $L_1 \cap L_2$  must be regular.
- (c) If  $L$  is not of type 2 (i.e., not context-free), then it is not of type 3 (i.e., not regular).
- (d) If the language  $L^*$  is regular, then  $L$  must be regular.

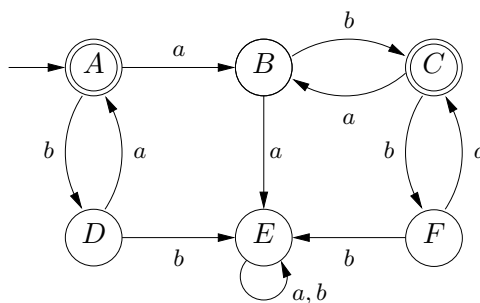
**Problem 2** [3 points] Consider the regular expression  $E = 1 \cdot (0^* + 0 \cdot 1)^*$ . Construct an  $\epsilon$ -NFA  $A$  such that  $\mathcal{L}(A) = \mathcal{L}(E)$ . Illustrate the steps of the algorithm you have followed to construct  $A$ .

**Problem 3** [6 points] Consider the following DFA  $A$  over  $\{0, 1\}$ :



Construct a regular expression  $E$  such that  $\mathcal{L}(E) = \mathcal{L}(A)$ . Illustrate the steps of the algorithm you have followed to construct  $E$ .

**Problem 4** [6 points] Consider the following DFA  $A$  over  $\{a, b\}$ :



Construct a DFA  $A_m$  with minimal number of states such that  $\mathcal{L}(A_m) = \mathcal{L}(A)$ . Illustrate the steps of the algorithm you have followed to construct  $A_m$ .

**Problem 5** [4 points] Show that the language  $L = \{a^i b^j \mid i, j \geq 0, i \neq j\}$  is not regular. [Hint: Exploit in your argument closure properties of regular languages and the known facts that the language  $L_1 = \{a^i b^j \mid i, j \geq 0\}$  is regular and the language  $L_2 = \{a^n b^n \mid n \geq 0\}$  is not regular.]

**Problem 6** [5 points] Consider the grammar  $G = (\{S\}, \{a, b\}, P, S)$ , where  $P$  consists of the following productions

$$S \longrightarrow Sa \mid bSS \mid SSb \mid Sbs \mid a$$

Prove that every string in  $\mathcal{L}(G)$  has more  $a$ 's than  $b$ 's.