Introduction to Databases Exam of 14/09/2021 With Solutions

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Design the Entity-Relationship schema of an application relating to maintenance operations on movable power supplies. For each **power supply**, the code (identifier), the model, the company that built it, and the year of construction are of interest. **Gas power supplies** are a special kind of power supplies, the cost of which is also of interest. For each **company**, the fiscal code (identifier), the city where it is based, and the number of employees are of interest (this last information may not be available). For each **city**, the name (unique within the region), the region, and the number of inhabitants are of interest. Power supplies are subject to maintenance: for each maintenance operation we are interested in the power supply to which it is applied, the date, the labour cost, and the city where the operation took place. A power supply cannot be subject to more than one maintenance operation per day. For a maintenance operation that concerns a gas power supply, we are interested also in the company responsible for the maintenance operation. A maintenance operation may require the replacement of parts. For each **part**, the code (identifier), the unit cost, and the manufacturing company are of interest. If the part is an **engine part**, it is also important to know the reliability index (positive integer). Finally, for each part to be replaced in a maintenance operation, it is of interest to know the quantity that has been replaced (for example, 3 valves may be replaced in a maintenance operation of a power supply), and in the case of an engine part, the discount applied on the unit price.

Problem 1: Conceptual schema – Diagram



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Problem 1: Conceptual schema – External constraints

- For each instance *I* of the schema, for each (M:*m*,P:*p*) ∈ *instances*(*I*,SubjectTo), if *p* ∈ *instances*(*I*,GasPS), then *m* ∈ *instances*(*I*,MaintGasPS) (and hence (M:*m*,P:*p*) ∈ *instances*(*I*,STGasPS)).
- 2) For each instance *I* of the schema, for each (M:*m*,P:*p*) ∈ *instances*(*I*,**Replaces**), if *p* ∈ *instances*(*I*,**EnginePart**), then (M:*m*,P:*p*) ∈ *instances*(*I*,**RepEng**).

Carry out the logical design of the database, producing the complete relational schema with constraints, taking into account the following indications:

- When accessing data about a maintenance operation, we are always interested in knowing the city where it took place and whether it is a maintenance operation concerning a gas power supply, and in this case also the company responsible for it.
- 2. When accessing data about a power supply, we are always interested in knowing the company that built it.

As steps in your design you should produce:

- the restructured ER schema (possibly with external constraints),
- the direct translation into the relational model (possibly with external constraints), and
- the restructured relational schema (again with constraints).

Motivate explicitly how the above indications affect your design.

Problem 2: Restructured conceptual schema – Diagram



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Problem 2: Restructured conceptual schema – External constraints

- For each instance *I* of the schema, for each (M:*m*,P:*p*) ∈ *instances*(*I*,**SubjectTo**), if there is a *g* ∈ *instances*(*I*,**GasPS**) such that (GasPS:*g*,PowerSupply:*p*) ∈ *instances*(*I*,**ISA-G-P**), then there is a (unique) *mg* ∈ *instances*(*I*,**MaintGasPS**) such that (MaintGasPS:*mg*,Maintenance:*m*) ∈ *instances*(*I*,**ISA-G-M**) and (M:*mg*,P:*g*) ∈ *instances*(*I*,**STGasPS**).
- 2) For each instance *I* of the schema, for each (M:*m*,P:*p*) ∈ *instances*(*I*,**Replaces**), if there is an *e* ∈ *instances*(*I*,**EnginePart**) such that (EnginePart:*e*,Part:*p*) ∈ *instances*(*I*,**ISA-E-P**), then (M:*m*,P:*e*) ∈ *instances*(*I*,**RepEng**).
- 3) Constraint resulting from the elimination of ISA between the relationships STGasPS and SubjectTo: For each instance *I* of the schema, if (M:*mg*,P:*g*) ∈ *instances*(*I*,STGasPS), *m* is the unique element of *instances*(*I*,Maintenance) such that (MaintGasPS:*mg*,Maintenance:*m*) ∈ *instances*(*I*,ISA-G-M), and *p* is the unique element of *instances*(*I*,PowerSupply) such that (GasPS:*g*,PowerSupply:*p*) ∈ *instances*(*I*,ISA-G-P), then (M:*m*,P:*p*) ∈ *instances*(*I*,SubjectTo).
- 4) Constraint resulting from the elimination of ISA between the relationships ReplEng and Replaces: For each instance / of the schema, if (M:m,P:e) ∈ instances(/,ReplEng) and p is the unique element of instances(/,Part) such that (EnginePart:e,Part:p) ∈ instances(/,ISA-E-P), then (M:m,P:p) ∈ instances(/,Replaces).

Problem 2: Direct translation (1/2)

```
PowerSupply(code, model, year)
  foreign key: PowerSupply[code] 
BuiltBy[powerSupply]
Company(fiscalCode, #emp*)
  foreign key: Company[fiscalCode] 

BasedIn[company]
Maintenance(date, powerSupply, cost)
  foreign key: Maintenance[powerSupply] 
PowerSupply[code]
  foreign key: Maintenance[date,powerSupply] \subset InCity[date,PowerSupply]
GasPS(code, cost)
  foreign key: GasPS[code] 

PowerSupply[code]
MaintGasPS(date, powerSupply)
  foreign key: MaintGasPS[date,powerSupply] 
Maintenance[date,powerSupply]
  foreign key: MaintGasPS[date,powerSupply] 

RespFor[date,powerSupply]
  foreign key: MaintGasPS[date,powerSupply] 
C STGasPS[date,powerSupply]
City(<u>name</u>, <u>region</u>, #inhabitants)
Part(code, cost)
  foreign key: Part[code] 

Manufacturer[part]
EnginePart(code, relIndex)
  foreign key: EnginePart[code] 

Part[code]
```

Problem 2: Direct translation (2/2)

BuiltBy(powerSupply, company)

foreign key: BuiltBy[powerSupply] ⊂ PowerSupply[code] foreign key: BuiltBy company c Company fiscal Code

BasedIn(company, city, region)

foreign key: BasedIn[company] ⊆ Company[fiscalCode] foreign key: BasedIn[city, region] \subset City[name, region]

RespFor(date, powerSupply, company)

foreign key: RespFor[date,powerSupply]
 MaintGasPS[date,powerSupply] foreign key: RespFor[company] C Company[fiscalCode]

STGasPS(<u>date</u>, <u>powerSupply</u>, gasPS) foreign key: STGasPS[date,powerSupply] ⊆ MaintGasPS[date,powerSupply] foreign key: STGasPS[gasPS] ⊆ GasPS[code]

InCity(date, powerSupply, city, region)

foreign key: InCity[date,powerSupply]
C Maintenance[date,powerSupply] foreign key: $InCity[city, region] \subset City[name, region]$

Replaces(date, powerSupply, part, quantity)

foreign key: Replaces[date,powerSupply] ⊆ Maintenance[date,powerSupply] foreign key: Replaces part - Part code

ReplEng(date, powerSupply, enginePart, discount)

foreign key: ReplEng[date,powerSupply]
_ Maintenance[date,powerSupply] foreign key: ReplEng[enginePart]
C EnginePart[code] foreign key: ReplEng[date.powerSupply.enginePart]

Replaces[date.powerSupply.part]

Manufacturer(part, company)

foreign key: Manufacturer[part] ⊆ Part[code] foreign key: Manufacturer company] C Company[fiscalCode]

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Problem 2: Direct translation – External constraints

- 1) $PROJ_{date,powerSupply}$ (Maintenance $JOIN_{powerSupply=code}$ GasPS) \subseteq MaintGasPS.
- 2) $PROJ_{date,powerSupply,part}(Replaces JOIN_{part=code} EnginePart) \subseteq PROJ_{date,powerSupply,enginePart}(ReplEng).$
- 3) The constraint resulting from the elimination of ISA between the relationships STGasPS and SubjectTo is already implied by the foreign key between the relations STGasPS and MaintGasPS and the foreign key between the relations MaintGasPS and Maintenance (into which the relationship SubjectTo has been merged).
- The constraint resulting from the elimination of ISA between the relationships ReplEng and Replaces has already been specified as foreign key between the corresponding relations ReplEng and Replaces.

Problem 2: Restructuring of the relational schema

- 1. When accessing data about a maintenance operation, we are always interested in knowing the city where it took place and whether it is a maintenance operation concerning a gas power supply, and in this case also the company responsible for it.
- 2. When accessing data about a power supply, we are always interested in knowing the company that built it.

We take into account the above indications in the following way:

- We take into account indication 1 by:
 - Merging into the relation MaintGasPS the relation RespFor, from which and to which MaintGasPS has a foreign key, and then merging the resulting relation into Maintenance. Since there is a foreign key from MaintGasPS to Maintenance but not vice-versa, the latter merging operation might introduce NULL values in Maintenance.
 - 2. Merging into the relation **Maintenance** the relation **InCity**, from which and to which **Maintenance** has a foreign key.
- We take into account indication 2 by merging into the relation PowerSupply the relation BuiltBy, from which and to which PowerSupply has a foreign key.

Problem 2: Restructured relational schema

We specify here only the relations with their constraints that have been changed with respect to the schema obtained through the direct translation.

Maintenance(date, powerSupply, cost, company*, city, region) foreign key: Maintenance[powerSupply] \subseteq PowerSupply[code] foreign key: Maintenance[company] \subseteq Company[fiscalCode] foreign key: Maintenance[date,powerSupply] \subseteq InCity[date,PowerSupply] foreign key: Maintenance[city,region] \subseteq City[name,region]

STGasPS(<u>date</u>, <u>powerSupply</u>, gasPS)

foreign key: STGasPS[date,powerSupply] \subseteq Maintenance[date,powerSupply] foreign key: STGasPS[gasPS] \subseteq GasPS[code]

PowerSupply(<u>code</u>, model, year, company) foreign key: PowerSupply[company] ⊆ Company[fiscalCode]

Moreover, we remove the relations MaintGasPS, RespFor, and BuiltBy.

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Problem 2: Restructuring of the relational schema – External constraints

- 1) SEL_{company IS NULL}(Maintenance JOIN_{powerSupply=code} GasPS) = Ø.
- 2) This constraint is not affected by the restructuring.
- The constraint resulting from the elimination of ISA between the relationships STGasPS and SubjectTo is already specified as foreign key between the relations STGasPS and Maintenance (into which the relationship SubjectTo has been merged).
- 4) This constraint is not affected by the restructuring and stays specified as foreign key between the relations **RepIEng** and **Replaces**.
- 5) The foreign keys MaintGasPS[date,powerSupply] ⊆ STGasPS[date,powerSupply] and STGasPS[date,powerSupply] ⊆ MaintGasPS[date,powerSupply] become the following external constraint:

 $\begin{aligned} \mathsf{PROJ}_{\mathsf{powerSupply}}(\mathsf{SEL}_{\mathsf{company}\,\mathsf{IS}\,\mathsf{NOT}\,\mathsf{NULL}}(\mathsf{Maintenance})) &\subseteq \mathsf{PROJ}_{\mathsf{code}}(\mathsf{GasPS}) \\ \mathsf{PROJ}_{\mathsf{code}}(\mathsf{GasPS}) &\subseteq \mathsf{PROJ}_{\mathsf{powerSupply}}(\mathsf{SEL}_{\mathsf{company}\,\mathsf{IS}\,\mathsf{NOT}\,\mathsf{NULL}}(\mathsf{Maintenance})). \end{aligned}$

Consider a database that includes the relations Graduated and School. The relation Graduated (studId,grade,schoolCode) stores for each graduated student, the student id, the grade obtained at the graduation, and the school where they obtained it. The relation School (schoolCode,city) stores for each school the city.

- 1. We call "exit indicator" of a school the average of the grades obtained by the graduates in that school, and "average exit indicator" the average of the exit indicators of all the schools. Then we call "school quality indicator" of a city the percentage of schools that are located in that city and that have the exit indicator higher than the average exit indicator. Write a *SQL query* that returns for each city its school quality indicator.
- 2. We call a city a "top city" if all its schools have at least one graduate with a grade of 100. Write a query in *relational algebra* that returns the top cities.

Problem 3: Solution (1/2)

Graduated(studId,grade,schoolCode) School(schoolCode,city)

1. We call "exit indicator" of a school the average of the grades obtained by the graduates in that school, and "average exit indicator" the average of the exit indicators of all the schools. Then we call "school quality indicator" of a city the percentage of schools that are located in that city and that have the exit indicator higher than the average exit indicator. Write a *SQL query* that returns for each city its school quality indicator.

Problem 3: Solution (2/2)

Graduated(studId,grade,schoolCode)
School(schoolCode,city)

2. We call a city a "top city" if all its schools have at least one graduate with a grade of 100. Write a query in *relational algebra* that returns the top cities.

```
PROJ<sub>city</sub>(School)

–

PROJ<sub>city</sub>(School JOIN (PROJ<sub>schoolCode</sub>(School) –

PROJ<sub>schoolCode</sub>(SEL<sub>grade = 100</sub>(Graduated))))
```



Consider the restructured conceptual schema shown above, and carry out the direct translation step, illustrating the relational schema complete with constraints resulting from that step.

Problem 4: Solution



 $\begin{array}{l} \mathsf{E}(a, \underline{b}, c, \underline{f}, \underline{g}, \underline{h}) \\ \text{foreign key: } \mathsf{E}[f] \subseteq \mathsf{F}[d] \\ \text{foreign key: } \mathsf{E}[g,h] \subseteq \mathsf{G}[m,h] \\ \text{key: } \mathsf{c} \end{array}$

F(<u>d</u>)

 $\begin{array}{l} G(\underline{m},\,n,\,\underline{h}) \\ \text{ foreign key: } G[h] \subseteq H[o] \end{array}$

 $H(\underline{o}, p)$ inclusion: $H[o] \subseteq G[h]$ Notes:

- Since the primary identifier of entity E involves a role of relationship R, R gets merged into relation E.
- Since the primary identifier of entity **G** involves a role of relationship **Q**, **Q** gets merged into relation **G**.
- We have chosen as names for the attributes that refer to external entities, the name of the entity rather than the name of the attribute (when such attribute is part of the primay identifier of the external entity).