3/11/2010
E 5.1

Turing Machines computing functions:

We consider number-theoretic functions $f: \mathbb{N}^k \to \mathbb{N}$

- natural numbers are represented in unary

n is represented by $\bar{n} = 1^{n+1}$

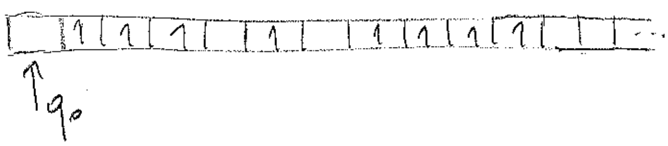
e.g.

n	$\bar{n}$
0	1
1	11
2	111
⋮	⋮

Hence, the input alphabet of a TM is $\Sigma = \{1\}$

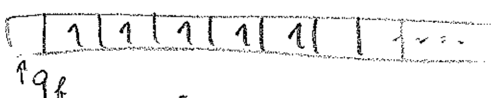
- consecutive numbers on the input are separated by a $\not{b}$
- the TM starts its computation with the head placed on a $\not{b}$ preceding the first number

e.g. for $f(2,0,3)$



- when the computation terminates, the input has been replaced on the tape with the result of the function

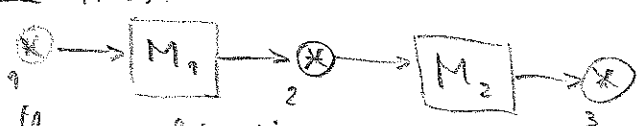
e.g. if $f(2,0,3) = 4$



We make some assumptions to ease composition of TMs:

1) there is a single final state q_f
2) the only transition from q_0 is $\delta(q_0, \not{b}) = (q_i, \not{b}, R)$
3) there are no transitions entering q_0 or of the form $\delta(q_f, \not{b})$
4) the computation loops whenever $f(n)\uparrow$, i.e. $f(n)$ is undefined

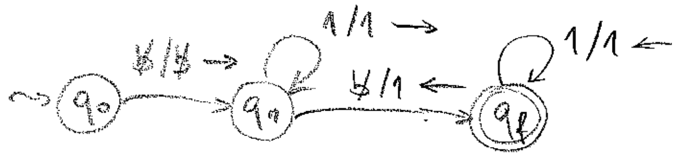
This allows us to sequentially compose TMs: represented by the diagram



1 … initial state of M_1 and of the combination
3 … final — " — M_2 — " —
2 … final state of M_1 is also the initial state of M_2

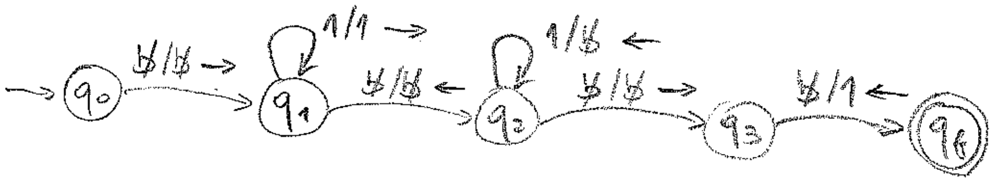
1) Construct a TM computing the successor function $s(n) = n+1$

S:



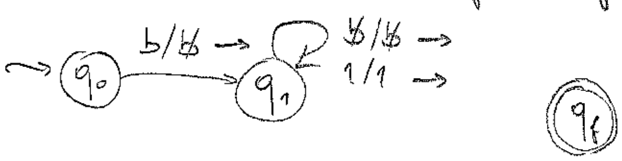
2) Construct a TM computing the zero function $z(n) = 0$

Z :



3) Construct a TM computing the empty function $e(n)\uparrow$ i.e., the function that is undefined for every $n \in \mathbb{N}$

E :

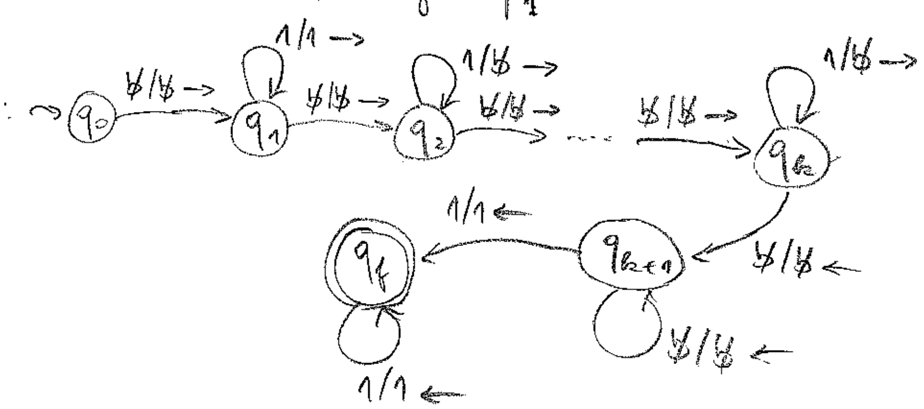


The k-variable projection function $\pi_i^{(k)}$ is defined as

$$\pi_i^{(k)}(n_1, \dots, n_k) = n_i \qquad (\text{for } 1 \le i \le k)$$

4) Construct a TM computing $\pi_1^{(k)}$

$P_1^{(k)}$:

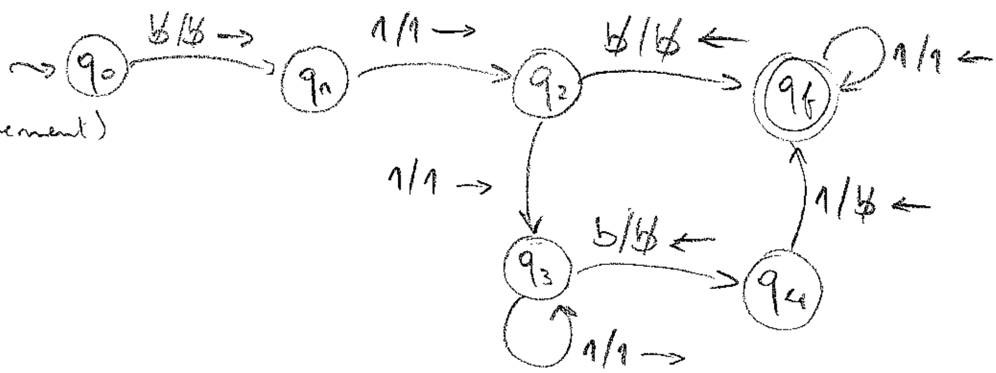


Note $\pi_1^{(1)}$ is also called the identity function $id(n) = n$

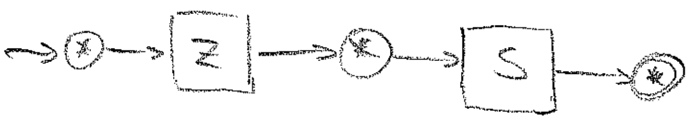
5) Construct a TM computing the predecessor function

$$pred(n) = \begin{cases} 0 & \text{if } n = 0 \\ n-1 & \text{if } n > 0 \end{cases}$$

D :
(for decrement)



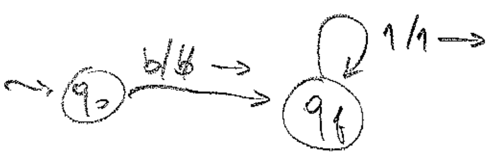
6) Using sequential composition, construct a TM computing the constant function $c(n) = 1$

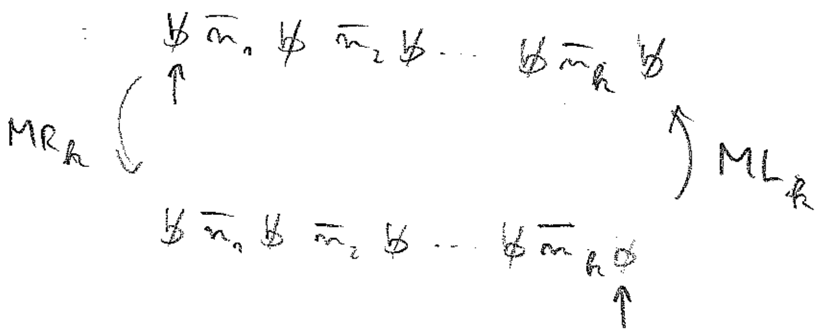


We now define some TMs that can be used as macros:

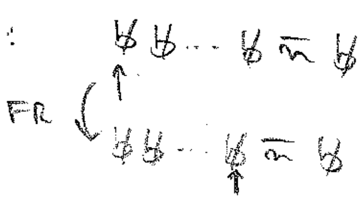
MR_k: move the tape head to the right through k consecutive natural numbers

e.g. MR_1





FR (find right):



FL (find left): $\not b\ \bar{n}\ \not b \cdots \underset{\uparrow}{\not b}$

$\longrightarrow\ \underset{\uparrow}{\not b}\ \bar{n}\ \not b \cdots \not b$

E_k (erase): $\underset{\uparrow}{\not b}\ \bar{n}_1\ \not b\ \bar{n}_2\ \not b \cdots \not b\ \bar{n}_k\ \not b$

$\longrightarrow\ \underset{\uparrow}{\not b}\ \not b \cdots\cdots \not b\ \not b$

CPY_k (copy): $\underset{\uparrow}{\not b}\ \bar{n}_1\ \not b\ \bar{n}_2\ \not b \cdots \not b\ \bar{n}_k\ \not b\ \not b \cdots\cdots \not b\ \not b$

$\longrightarrow\ \underset{\uparrow}{\not b}\ \bar{n}_1\ \not b\ \bar{n}_2\ \not b \cdots \not b\ \bar{n}_k\ \not b\ \bar{n}_1\ \not b\ \bar{n}_2\ \not b \cdots \not b\ \bar{n}_k\ \not b$

$CPY_{k,i}$ (copy through i numbers)

$\underset{\uparrow}{\not b}\ \bar{n}_1\ \not b \cdots \not b\ \bar{n}_k\ \not b\ \bar{n}_{k+1}\ \not b \cdots \not b\ \bar{n}_{k+i}\ \not b\ \not b \cdots\cdots \not b\ \not b$

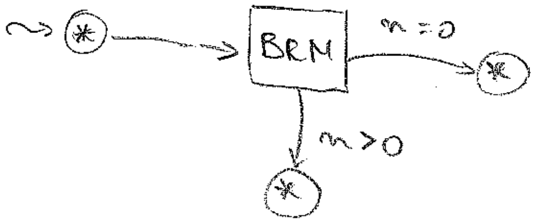
$\longrightarrow\ \underset{\uparrow}{\not b}\ \bar{n}_1\ \not b \cdots \not b\ \bar{n}_k\ \not b\ \bar{n}_{k+1}\ \not b \cdots \not b\ \bar{n}_{k+i}\ \not b\ \bar{n}_1\ \not b \cdots \not b\ \bar{n}_k\ \not b$.

for CPY_k and $CPY_{k,i}$ the blank portion following $\bar{n}_1 \not b \cdots \not b\, \bar{n}_k$ is assumed to be long enough to contain the copy.

T (translate) $\not b\ \not b \cdots \underset{\uparrow}{\not b}\ \bar{n}\ \not b$

$\longrightarrow\ \underset{\uparrow}{\not b}\ \bar{n}\ \not b \cdots \not b$

BRN (branch on zero)

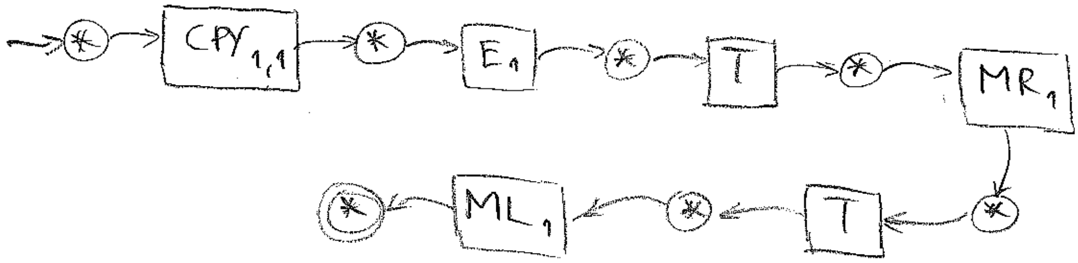


it does not alter the tape or change the head position

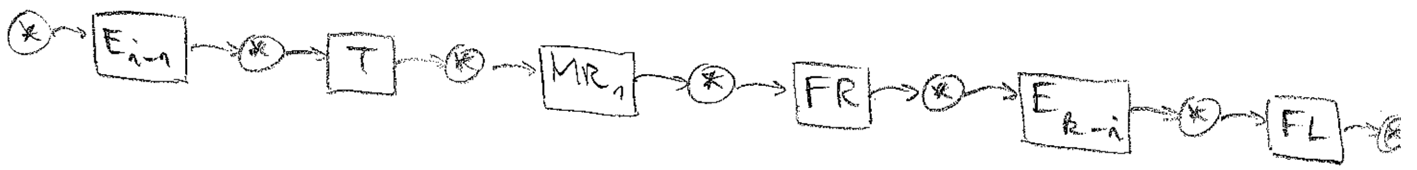
7) Using composition and the defined macros, construct a TM INT (interchange): $\not b\, \overline{n}\, \not b\, m\, \not b\, b^{n+1}\, \not b$

$\uparrow$

$\not b\, \overline{m}\, \not b\, n\, \not b\, \not b^{n+1}\, \not b$



8) Using composition and the defined macros, construct a TM for $P_i^{(k)}$.



Function composition:

Let $g_i : \mathbb{N}^k \to \mathbb{N}$ for $1 \le i \le n$

$h : \mathbb{N}^n \to \mathbb{N}$

The composition of h with $g_1, \dots, g_n$, written

$f = h \circ (g_1, \dots, g_n)$

is the function $f : \mathbb{N}^k \to \mathbb{N}$ defined by

$$f(x_1, \dots, x_k) = h(g_1(x_1, \dots, x_k), \dots, g_n(x_1, \dots, x_k))$$

The function $f(x_1, \dots, x_k)$ is undefined if either

1) $g_i(x_1, \dots, x_k)\uparrow$ for some $i \in \{1, \dots, n\}$

2) $g_i(x_1, \dots, x_k) = y_i$ for $i \in \{1, \dots, n\}$ and $h(y_1, \dots, y_n)\uparrow$

We show that the composition of Turing-computable functions is also Turing-computable.

Exercise: given $g_1 : \mathbb{N}^3 \to \mathbb{N}$, $g_2 : \mathbb{N}^3 \to \mathbb{N}$, $h : \mathbb{N}^2 \to \mathbb{N}$ and TMs G_1, G_2, H computing respectively g_1, g_2, h, construct a TM computing $h \circ (g_1, g_2)$.

Use macros and TM composition, and construct the configurations obtained after every macro application.

macro	configuration	
	$\underline{\not b}\,\overline{x_1}\,\not b\,\overline{x_2}\,\not b\,\overline{x_3}\,\not b$	
CPY_3	$\underline{\not b}\,\overline{x_1}\,\not b\,\overline{x_2}\,\not b\,\overline{x_3}\,\not b\,\overline{x_1}\,\not b\,\overline{x_2}\,\not b\,\overline{x_3}\,\not b$	
MR_3	$\not b$ —"— $\underline{\not b}$ —"— $\not b$	
G_1	$\not b$ —"— $\underline{\not b}\,\overline{y_1}\,\not b$	(where $y_1 = g_1(x_1, x_2, x_3)$)
ML_3	$\underline{\not b}\,\overline{x_1}\,\not b\,x_2\,\not b\,x_3\,\not b\,\overline{y}_1\,\not b$	
$CPY_{3,1}$	$\underline{\not b}$ —"— $\not b\,\overline{x_1}\,\not b\,\overline{x_2}\,\not b\,\overline{x_3}\,\not b$	
MR_4	$\not b$ —"— $\underline{\not b}$ —"— $\not b$	
G_2	$\not b$ —"— $\not b\,\overline{y_1}\,\underline{\not b}\,\overline{y_2}\,\not b$	(where $y_2 = g_2(x_1, x_2, x_3)$)
ML_1	$\not b$ —"— $\underline{\not b}\,\overline{y_1}\,\not b\,\overline{y_2}\,\not b$	
H	$\not b$ —"— $\underline{\not b}\,\overline{z}\,\not b$	(where $z = h(y_1, y_2)$)
ML_3	$\underline{\not b}$ —"— $\not b\,\overline{z}\,\not b$	
E_3	$\underline{\not b}\,\not b \dots \not b\,\overline{z}\,\not b$	
T	$\not b\,\overline{z}\,\not b$	