

Automata with  $\epsilon$ -Transitions

Exercise 1

Design  $\epsilon$ -NFA's for the following languages:

- a) The set of strings consisting of zero or more a's followed by zero or more b's, followed by zero or more c's.
- b) The set of all strings consist of either of repeated one or more times, or oio repeated one or more times.

Exercise 2

Consider the following  $\epsilon$ -NFA:

|                 | $\epsilon$  | a           | b           | c           |                              |
|-----------------|-------------|-------------|-------------|-------------|------------------------------|
| $\rightarrow p$ | {q, r}      | $\emptyset$ | {q}         | {r}         | $\epsilon$ -NFA <sub>1</sub> |
| q               | $\emptyset$ | {p}         | {r}         | {p, q}      |                              |
| *r              | $\emptyset$ | $\emptyset$ | $\emptyset$ | $\emptyset$ |                              |

- a) Compute the  $\epsilon$ -closure of each state.
- b) which of following strings are accepted by the automata:

bb      bba      a       $\epsilon$       aba

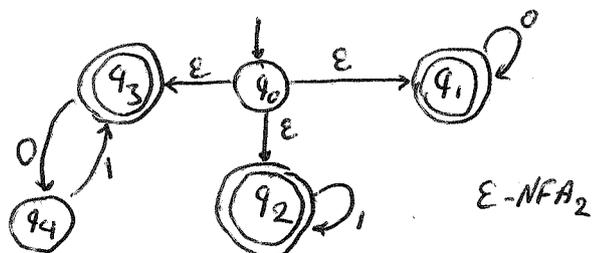
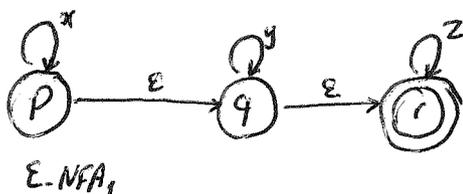
- c) Convert the automata to a DFA.

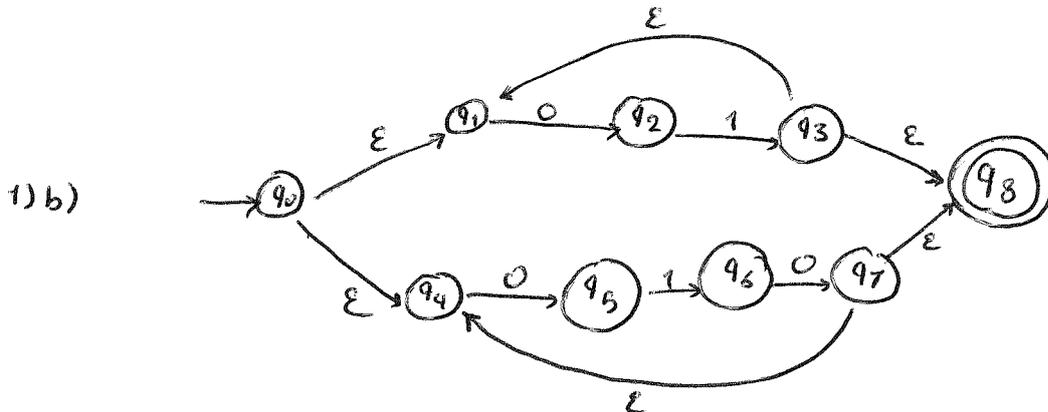
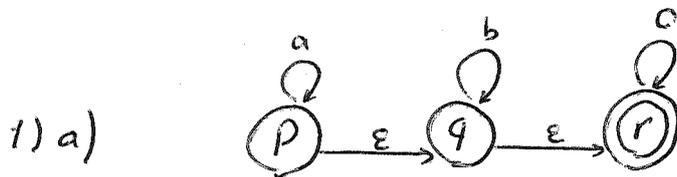
- d) Repeat steps a), b), c) for the following  $\epsilon$ -NFA:

|                 | $\epsilon$  | a   | b           | c           |                              |
|-----------------|-------------|-----|-------------|-------------|------------------------------|
| $\rightarrow p$ | $\emptyset$ | {p} | {q}         | {r}         | $\epsilon$ -NFA <sub>2</sub> |
| q               | {p}         | {q} | {r}         | $\emptyset$ |                              |
| *r              | {q}         | {r} | $\emptyset$ | {p}         |                              |

Exercise 3.

Convert the following NFA's to a DFA's:





2) a) Remember that:  $\begin{cases} q \in ECLOSE(q) \\ p \in ECLOSE(q) \ \& \ r \in \delta(p, \epsilon) \Rightarrow r \in ECLOSE(q) \end{cases}$

$\epsilon$ -NFA<sub>1</sub>:  $ECLOSE(p) = \{p, q, r\}$ ,  $ECLOSE(q) = \{q\}$ ,  $ECLOSE(r) = \{r\}$

2) b)  $bb \checkmark$     $bbax \checkmark$     $a \checkmark$     $\epsilon \checkmark$     $aba \checkmark$

2) c) From  $\epsilon$ -NFA to NFA  
Remember that  $\delta_N(q, a) = \hat{\delta}_\epsilon(q, a) = ECLOSE\left(\bigcup_{p \in ECLOSE(q)} \delta(p, a)\right)$

The calculation of  $\epsilon$ -NFA go a follow:

$$\hat{\delta}_\epsilon(p, a) = ECLOSE(\delta(p, a) \cup \delta(q, a) \cup \delta(r, a)) = ECLOSE(\emptyset \cup \{p\} \cup \emptyset) = \{p, q, r\}$$

$$\hat{\delta}_\epsilon(p, b) = ECLOSE(\delta(p, b) \cup \delta(q, b) \cup \delta(r, b)) = ECLOSE(\{q\} \cup \{r\} \cup \emptyset) = \{q, r\}$$

$$\hat{\delta}_\epsilon(p, c) = ECLOSE(\delta(p, c) \cup \delta(q, c) \cup \delta(r, c)) = ECLOSE(\{r\} \cup \{p, q\} \cup \emptyset) = \{p, q, r\}$$

$$\hat{\delta}_\epsilon(q, a) = ECLOSE(\delta(q, a)) = ECLOSE(\{p\}) = \{p, q, r\}$$

$$\hat{\delta}_\epsilon(q, b) = ECLOSE(\delta(q, b)) = ECLOSE(\{r\}) = \{r\}$$

$$\hat{\delta}_\epsilon(q, c) = ECLOSE(\delta(q, c)) = ECLOSE(\{p, q\}) = \{p, q, r\}$$

$$\hat{\delta}_\epsilon(r, a) = ECLOSE(\delta(r, a)) = ECLOSE(\emptyset) = \emptyset$$

$$\hat{\delta}_\epsilon(r, b) = ECLOSE(\delta(r, b)) = ECLOSE(\emptyset) = \emptyset$$

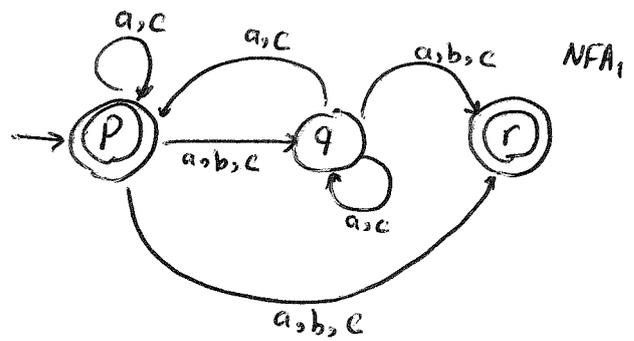
$$\hat{\delta}_\epsilon(r, c) = ECLOSE(\delta(r, c)) = ECLOSE(\emptyset) = \emptyset$$

2) c) Cont.

we get the following NFA:

| $\delta_N$      | a             | b           | c             |
|-----------------|---------------|-------------|---------------|
| $\rightarrow P$ | $\{P, q, r\}$ | $\{q, r\}$  | $\{P, q, r\}$ |
| q               | $\{P, q, r\}$ | $\{r\}$     | $\{P, q, r\}$ |
| $* r$           | $\emptyset$   | $\emptyset$ | $\emptyset$   |

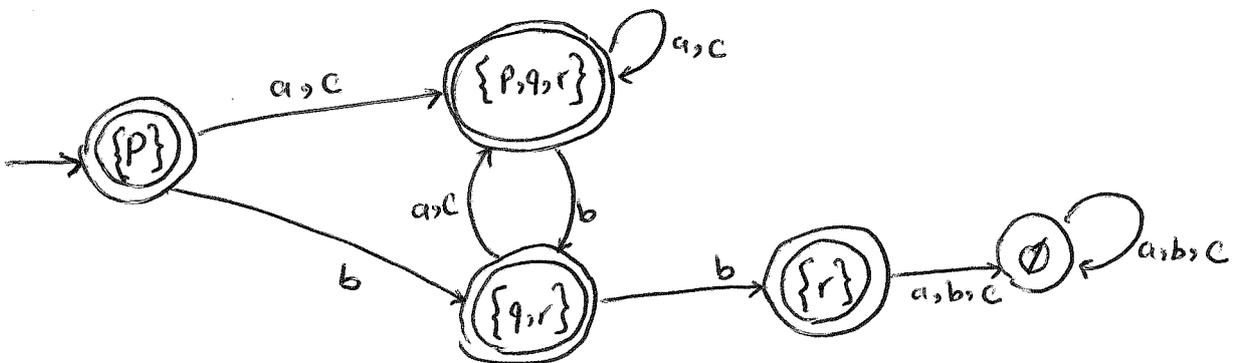
Note that  $\epsilon$ -NFA<sub>1</sub> accepts the empty string  $\epsilon$ .



From NFA to DFA

The subset construction yields:

|                     | a             | b           | c             |
|---------------------|---------------|-------------|---------------|
| $\emptyset$         | $\emptyset$   | $\emptyset$ | $\emptyset$   |
| $\rightarrow \{P\}$ | $\{P, q, r\}$ | $\{q, r\}$  | $\{P, q, r\}$ |
| $\{q\}$             | $\{P, q, r\}$ | $\{r\}$     | $\{P, q, r\}$ |
| $* \{r\}$           | $\emptyset$   | $\emptyset$ | $\emptyset$   |
| $* \{P, q\}$        | $\{P, q, r\}$ | $\{q, r\}$  | $\{P, q, r\}$ |
| $* \{P, r\}$        | $\{P, q, r\}$ | $\{q, r\}$  | $\{P, q, r\}$ |
| $\{q, r\}$          | $\{P, q, r\}$ | $\{r\}$     | $\{P, q, r\}$ |
| $* \{P, q, r\}$     | $\{P, q, r\}$ | $\{q, r\}$  | $\{P, q, r\}$ |



Attention: The states:  $\{q\}$ ,  $\{P, q\}$ ,  $\{P, r\}$  are not reachable from the initial state, so they are excluded from the NFA diagram.

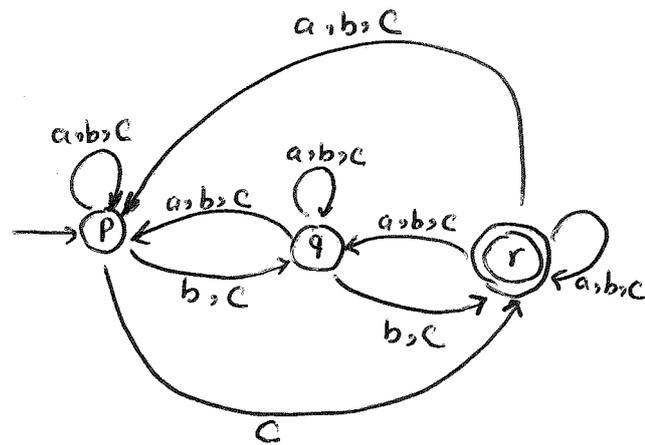
2)d)

a) E-NFA2:  $ECLOSE(P) = \{P\}$ ,  $ECLOSE(Q) = \{P, Q\}$ ,  $ECLOSE(r) = \{P, Q, r\}$

b)  $bb \checkmark$     $bba \checkmark$     $a \times$     $E \times$     $aba \checkmark$

c) E-NFA  $\rightarrow$  NFA:

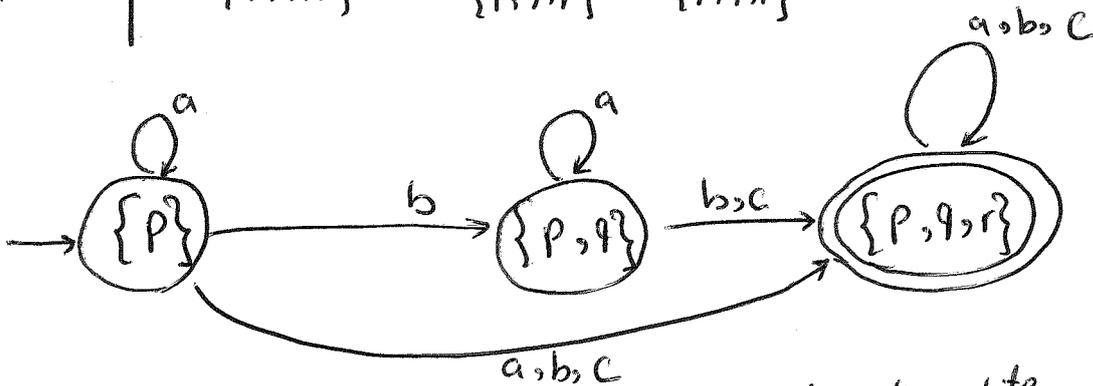
| $S_N$           | a             | b             | c             |
|-----------------|---------------|---------------|---------------|
| $\rightarrow P$ | $\{P\}$       | $\{P, Q\}$    | $\{P, Q, r\}$ |
| Q               | $\{P, Q\}$    | $\{P, Q, r\}$ | $\{P, Q, r\}$ |
| *r              | $\{P, Q, r\}$ | $\{P, Q, r\}$ | $\{P, Q, r\}$ |



NFA  $\rightarrow$  DFA:

The subset construction yields:

|                     | a             | b             | c             |
|---------------------|---------------|---------------|---------------|
| $\emptyset$         | $\emptyset$   | $\emptyset$   | $\emptyset$   |
| $\rightarrow \{P\}$ | $\{P\}$       | $\{P, Q\}$    | $\{P, Q, r\}$ |
| $\{Q\}$             | $\{P, Q\}$    | $\{P, Q, r\}$ | $\{P, Q, r\}$ |
| $\{r\}$             | $\{P, Q, r\}$ | $\{P, Q, r\}$ | $\{P, Q, r\}$ |
| $\{P, Q\}$          | $\{P, Q\}$    | $\{P, Q, r\}$ | $\{P, Q, r\}$ |
| $\{P, r\}$          | $\{P, Q, r\}$ | $\{P, Q, r\}$ | $\{P, Q, r\}$ |
| $\{Q, r\}$          | $\{P, Q, r\}$ | $\{P, Q, r\}$ | $\{P, Q, r\}$ |
| * $\{P, Q, r\}$     | $\{P, Q, r\}$ | $\{P, Q, r\}$ | $\{P, Q, r\}$ |



• other states are not reachable from the initial state.

3)  $\epsilon$ -NFA<sub>1</sub>From  $\epsilon$ -NFA<sub>1</sub> to NFA<sub>1</sub> (as above):

|                 | x             | y           | z       |
|-----------------|---------------|-------------|---------|
| $\rightarrow P$ | $\{P, q, r\}$ | $\{q, r\}$  | $\{r\}$ |
| q               | $\emptyset$   | $\{q, r\}$  | $\{r\}$ |
| * r             | $\emptyset$   | $\emptyset$ | $\{r\}$ |

From NFA<sub>1</sub> to DFA<sub>1</sub> (subset construction):

| $S_D$               | x             | y           | z           |
|---------------------|---------------|-------------|-------------|
| $\emptyset$         | $\emptyset$   | $\emptyset$ | $\emptyset$ |
| $\rightarrow \{P\}$ | $\{P, q, r\}$ | $\{q, r\}$  | $\{r\}$     |
| $\{q\}$             | $\emptyset$   | $\{q, r\}$  | $\{r\}$     |
| * $\{r\}$           | $\emptyset$   | $\emptyset$ | $\{r\}$     |
| $\{P, q\}$          | $\{P, q, r\}$ | $\{q, r\}$  | $\{r\}$     |
| * $\{P, r\}$        | $\{P, q, r\}$ | $\{q, r\}$  | $\{r\}$     |
| * $\{q, r\}$        | $\emptyset$   | $\{q, r\}$  | $\{r\}$     |
| * $\{P, q, r\}$     | $\{P, q, r\}$ | $\{q, r\}$  | $\{r\}$     |

 $\epsilon$ -NFA<sub>2</sub>

we only provide the final DFA:

