

# DETERMINISTIC FINITE AUTOMATA

30/10/2009  
E2.1

## EXERCISE 1

give a DFA accepting the following language over the alphabet  $\{a, b\}$ : the set of all strings such that the second last symbol is  $b$ .

## EXERCISE 2

give a DFA accepting the following language over the alphabet  $\{x, y\}$ : the set of strings that either begin or end (or both) with  $yx$ .

## EXERCISE 3

give a DFA accepting the following language over the alphabet  $\{0, 1\}$ : the set of strings such that the number of 0's is divisible by five and the number of 1's is divisible by three.

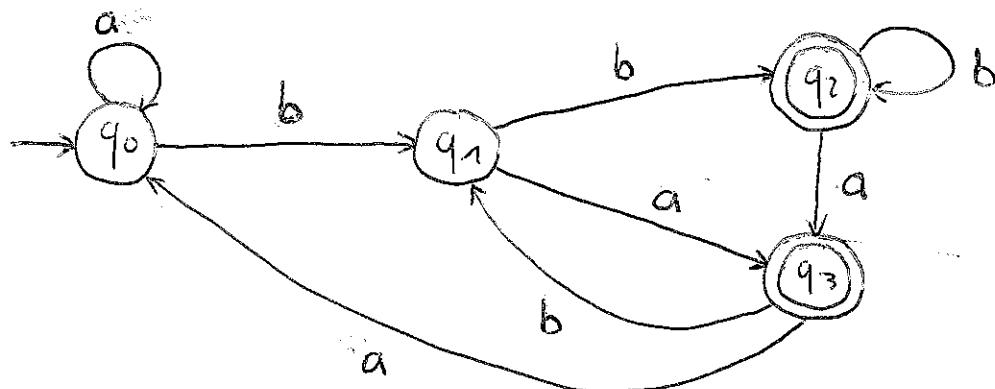
## EXERCISE 4

give a DFA accepting the following language over the alphabet  $\{a, b, c, d\}$ : the set of strings consisting of zero or more a's followed by zero or more b's, followed by zero or more c's.

# SOLUTIONS

E2.2

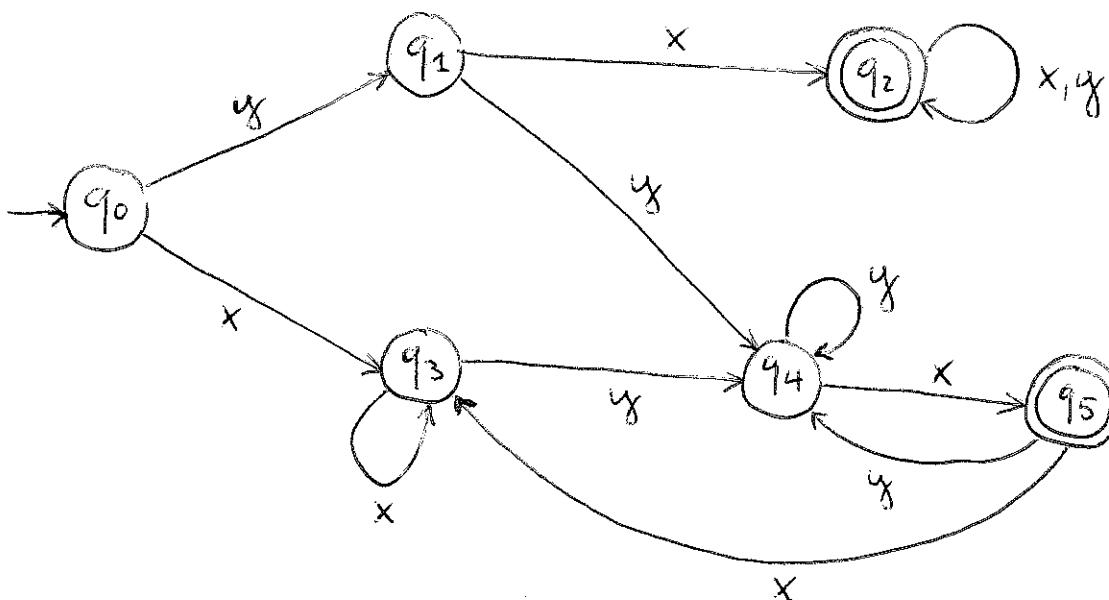
1) The DFA looks as follows:



We show that it accepts bbaaba but not abbbbaaa

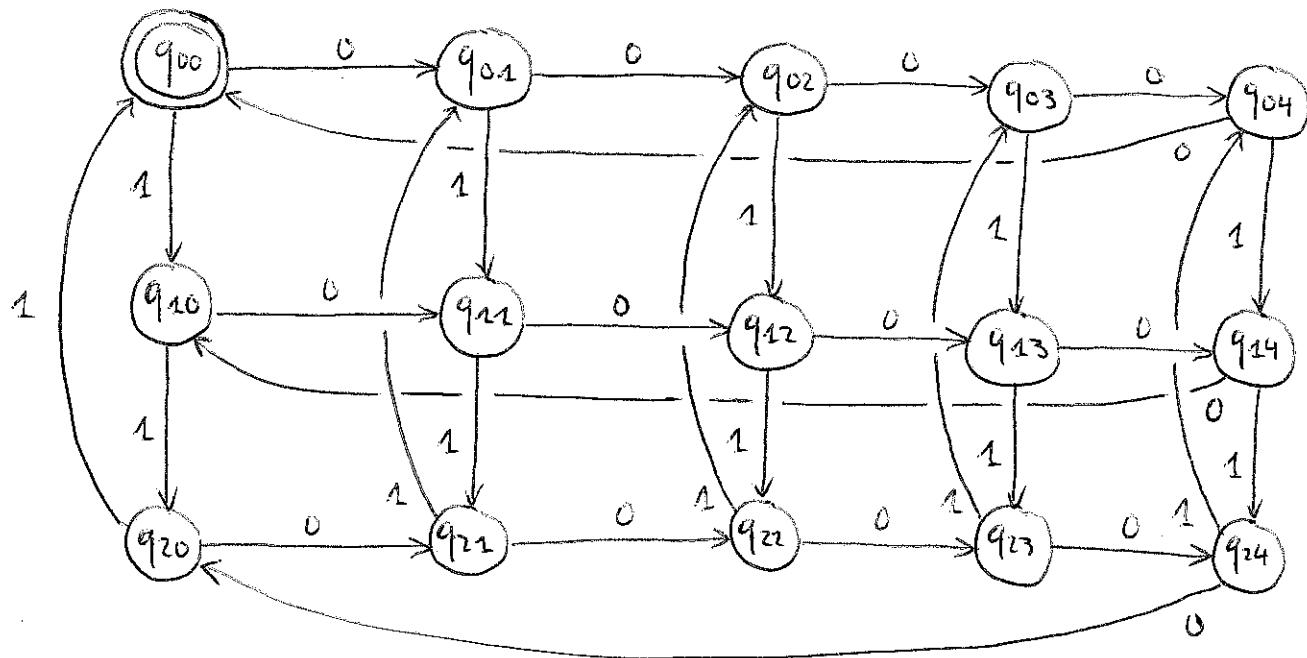
- $(q_0, b) \rightarrow q_1, (q_1, b) \rightarrow q_2, (q_2, a) \rightarrow q_3, (q_3, a) \rightarrow q_0$   
 $(q_0, b) \rightarrow q_1, (q_1, a) \rightarrow q_2$        $q_2$  is a final state
- $(q_0, a) \rightarrow q_0, (q_0, b) \rightarrow q_1, (q_1, b) \rightarrow q_2, (q_2, b) \rightarrow q_2$   
 $(q_2, a) \rightarrow q_3, (q_3, a) \rightarrow q_0$        $q_0$  is not a final state

2) The DFA looks as follows:

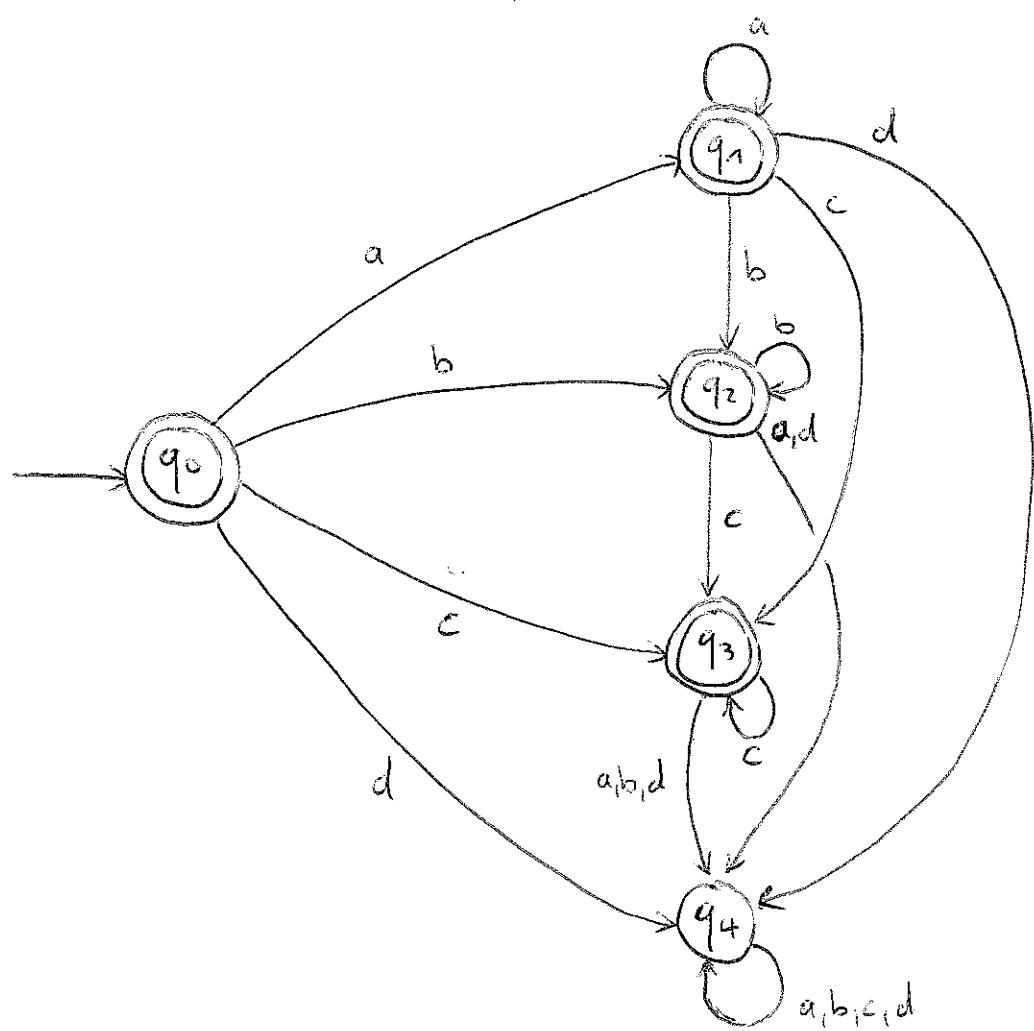


3) The DFA looks as follows:

EZ.3



4) The DFA looks as follows:



# NON-DETERMINISTIC FINITE AUTOMATA

## EXERCISE 1

Pick out one of the DFA's from exercise E2 (16/10/2008) and two strings of length at least five over the corresponding alphabet. Show whether the strings are accepted or not by using the extended transition function.

## EXERCISE 2

Give a NFA accepting the following language over the alphabet  $\{a, b\}$ : the set of strings that end with ba, bb, or baa. Then show that the string baab is not accepted by the NFA.

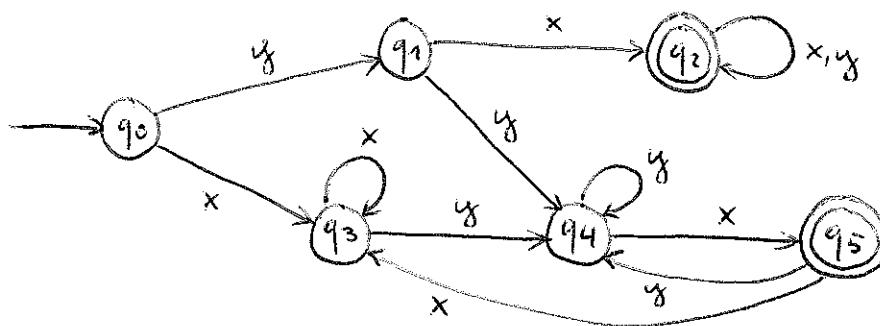
## EXERCISE 3

Give NFA's accepting the following languages:

- the set of strings over  $\{0, 1, \dots, 9\}$  such that the final digit has appeared before;
- the set of strings over  $\{0, 1, \dots, 9\}$  such that the final digit has not appeared before;
- the set of strings over  $\{0, 1\}$  such that there are two 0's separated by a number of positions that is a multiple of four.

SOLUTIONS

1) We choose the DFA from exercise 2



and show that:  $yxyxxyx$  is accepted;  $xyxxxy$  is not accepted.  
In other words, we show that:  $\hat{\delta}(q_0, yxyxxyx) = q_5 \in F$ ;  
 $\hat{\delta}(q_0, xyxxxy) = q_4 \notin F$ .

$$\hat{\delta}(q_0, \epsilon) = q_0$$

$$\hat{\delta}(q_0, y) = \delta(\hat{\delta}(q_0, \epsilon), y) = \delta(q_0, y) = q_1$$

$$\hat{\delta}(q_0, yy) = \delta(\hat{\delta}(q_0, y), y) = \delta(q_1, y) = q_4$$

$$\hat{\delta}(q_0, yyx) = \delta(\hat{\delta}(q_0, yy), x) = \delta(q_4, x) = q_5$$

$$\hat{\delta}(q_0, yyxy) = \delta(\hat{\delta}(q_0, yyx), y) = \delta(q_5, y) = q_4$$

$$\hat{\delta}(q_0, yyxxy) = \delta(\hat{\delta}(q_0, yyxy), x) = \delta(q_4, x) = q_5$$

$$\hat{\delta}(q_0, \epsilon) = q_0$$

$$\hat{\delta}(q_0, x) = \delta(\hat{\delta}(q_0, \epsilon), x) = \delta(q_0, x) = q_3$$

$$\hat{\delta}(q_0, xy) = \delta(\hat{\delta}(q_0, x), y) = \delta(q_3, y) = q_4$$

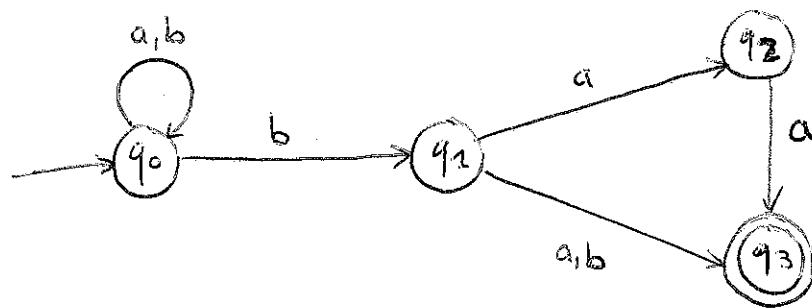
$$\hat{\delta}(q_0, xyy) = \delta(\hat{\delta}(q_0, xy), y) = \delta(q_4, y) = q_5$$

$$\hat{\delta}(q_0, xyxx) = \delta(\hat{\delta}(q_0, xyy), x) = \delta(q_5, x) = q_3$$

$$\hat{\delta}(q_0, xyxxx) = \delta(\hat{\delta}(q_0, xyxx), x) = \delta(q_3, x) = q_4$$

Note that we have underlined the recursive calls  
of the extended transition function  $\hat{\delta}$  in our calculations.

2) The NFA looks as follows:



We show that  $baab$  is not accepted, i.e. that  $\hat{S}(q_0, baab) = \{q_0, q_1\}$  and thus  $q_3 \notin \hat{S}(q_0, baab)$ .

$$\hat{S}(q_0, \epsilon) = \{q_0\}$$

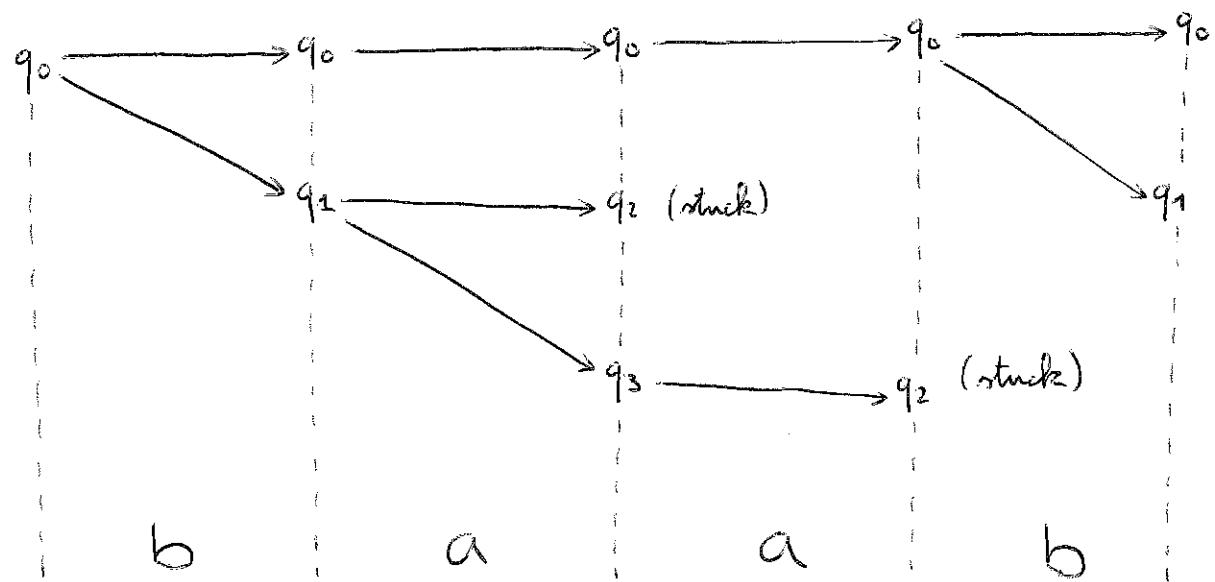
$$\hat{S}(q_0, b) = S(q_0, b) = \{q_0, q_1\}$$

$$\hat{S}(q_0, ba) = S(q_0, a) \cup S(q_1, a) = \{q_0\} \cup \{q_2, q_3\} = \{q_0, q_2, q_3\}$$

$$\hat{S}(q_0, baa) = S(q_0, a) \cup S(q_2, a) \cup S(q_3, a) = \{q_0\} \cup \{q_3\} \cup \emptyset = \{q_0, q_3\}$$

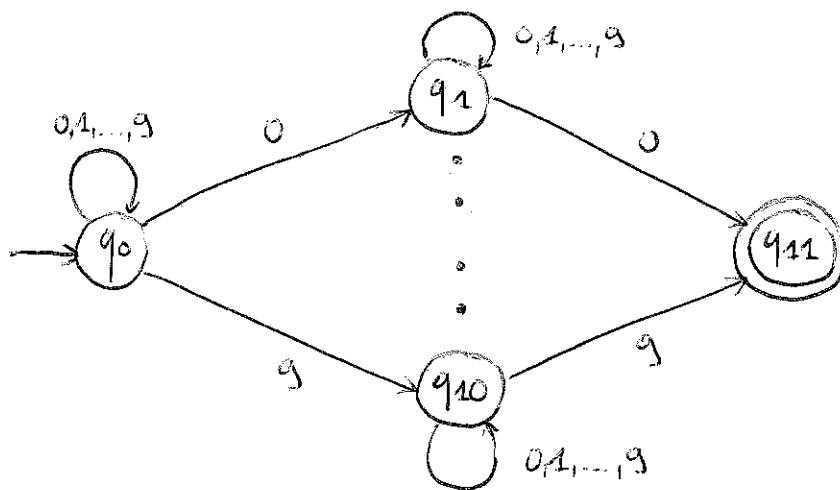
$$\hat{S}(q_0, baab) = S(q_0, b) \cup S(q_3, b) = \{q_0, q_1\} \cup \emptyset = \{q_0, q_1\}$$

The following graph might help to get a more intuitive understanding of what is going on.

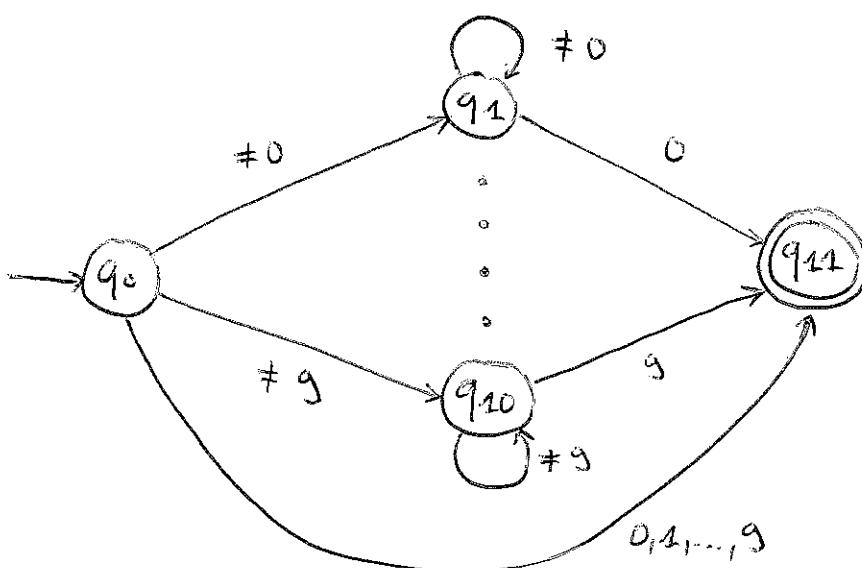


3) The NFA's look as follows:

a)



b)



c)

