

# Logic

## The Meaning of Entity-Relationship Diagrams

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# What is a Conceptual Schema

- A conceptual schema is a formal conceptualisation of the world.

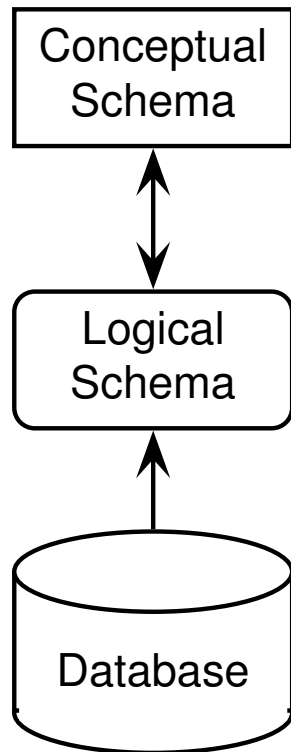
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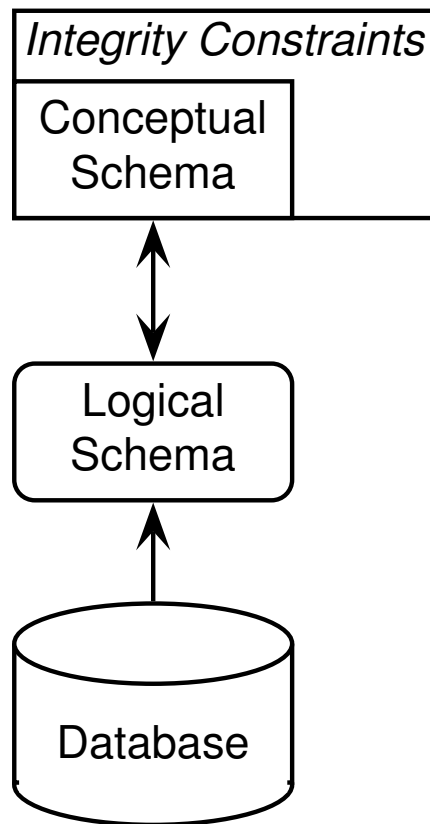
# What is a Conceptual Schema

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- A conceptual schema specifies a set of *constraints*, which declare what should necessarily hold in any possible database.
- Given a conceptual schema, a *legal database* is a database satisfying the constraints.

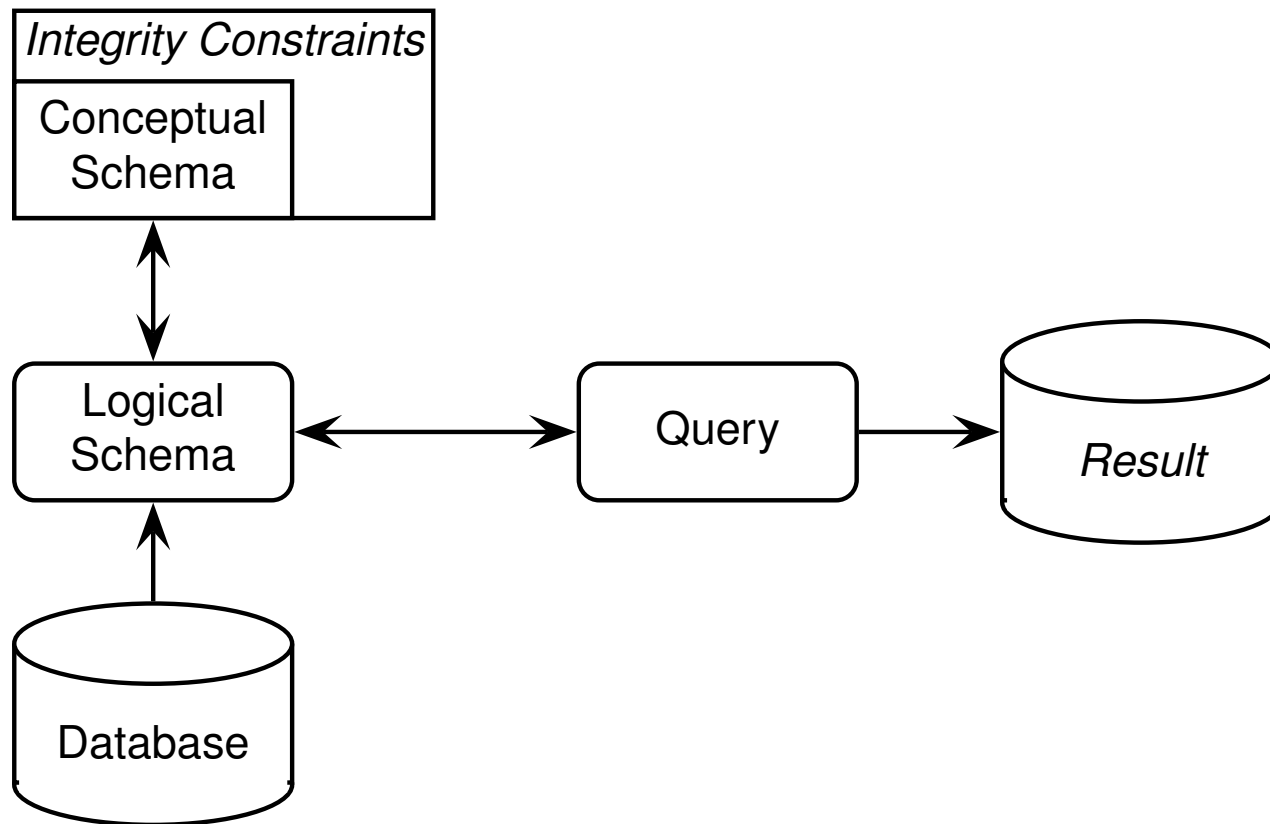
# The Architecture of a Database



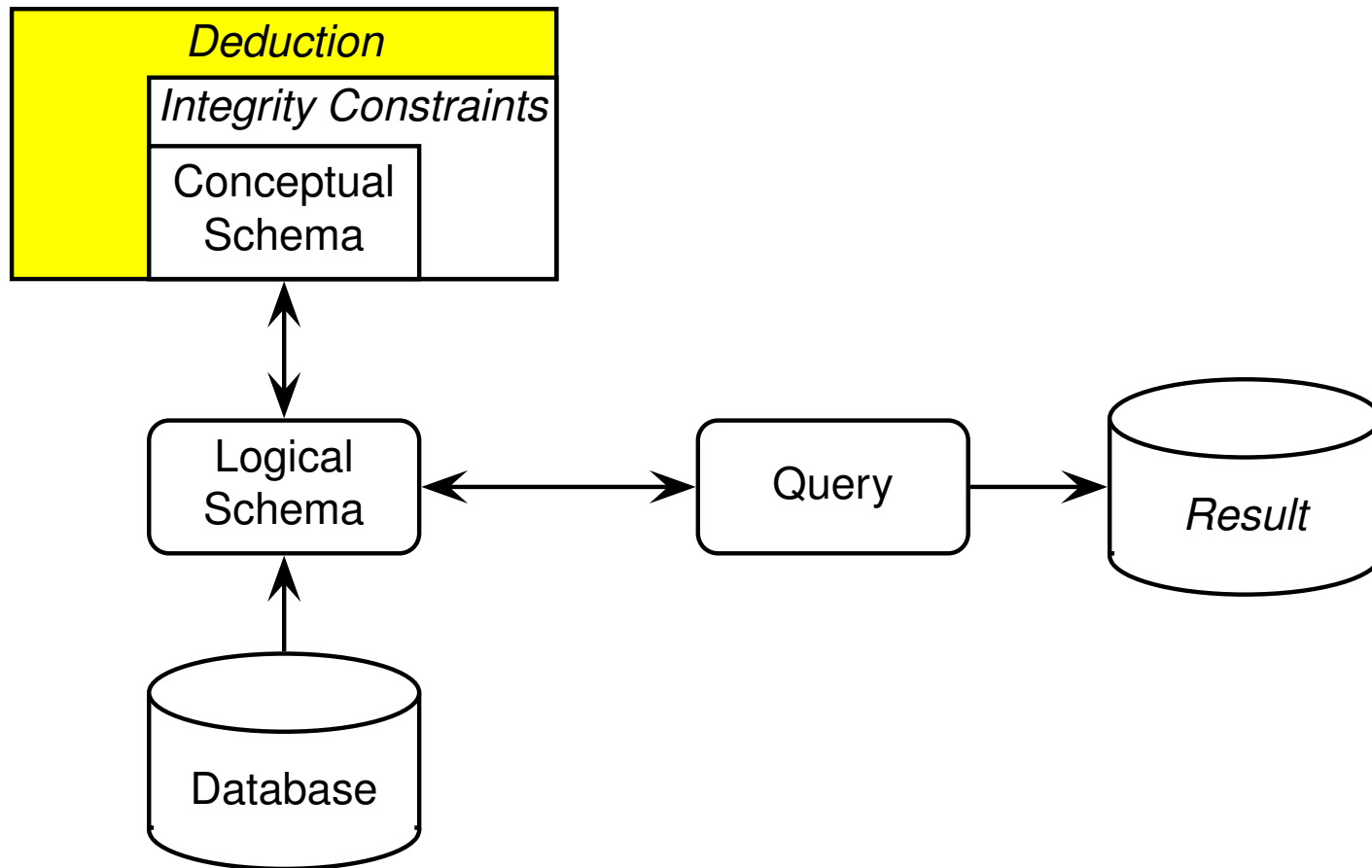
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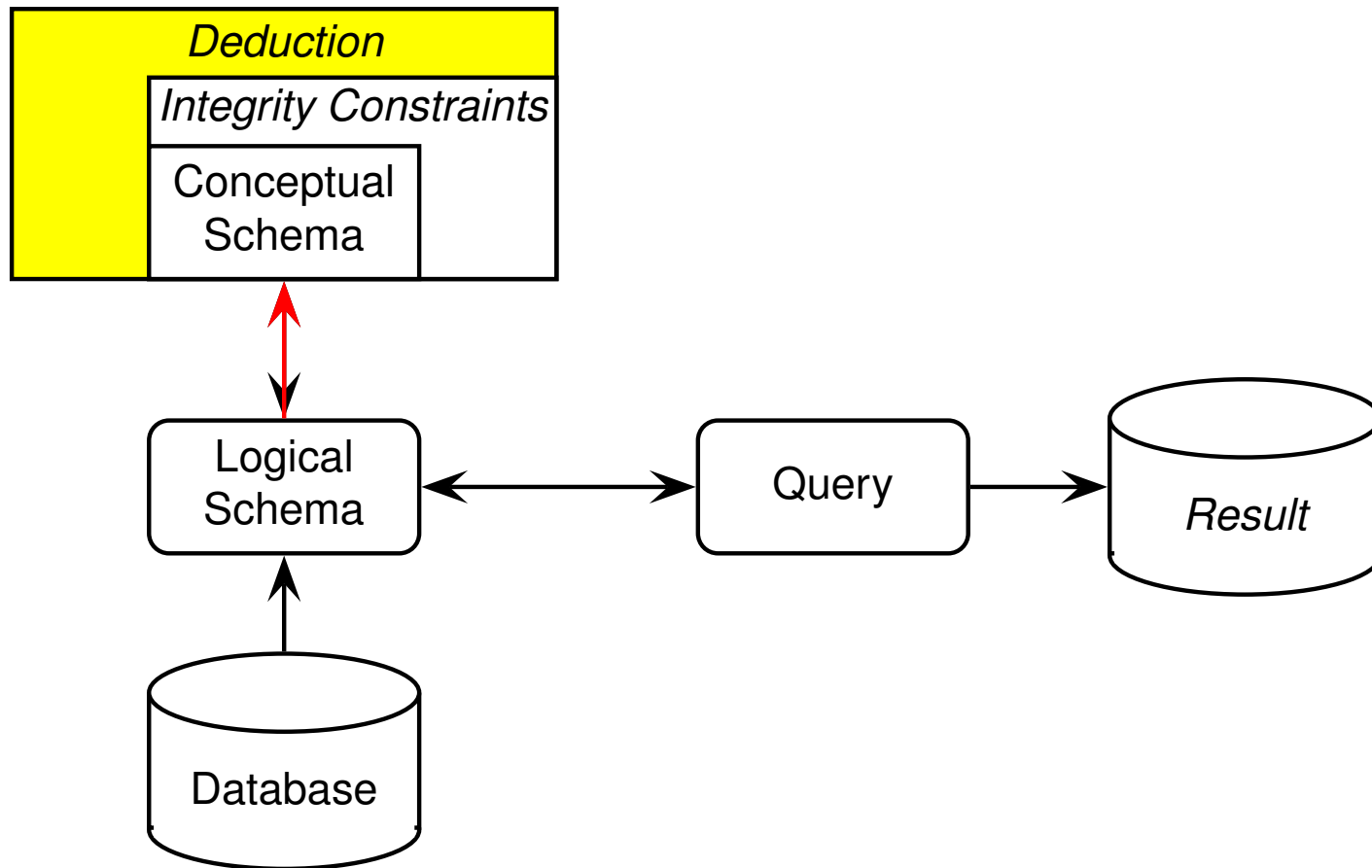


# The Architecture of a Database





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# Ontology languages and Conceptual Data Models

- An ontology language usually introduces *concepts* (aka classes, entities), *properties* of concepts (aka slots, attributes, roles), *relationships* between concepts (aka associations), and additional *constraints*.

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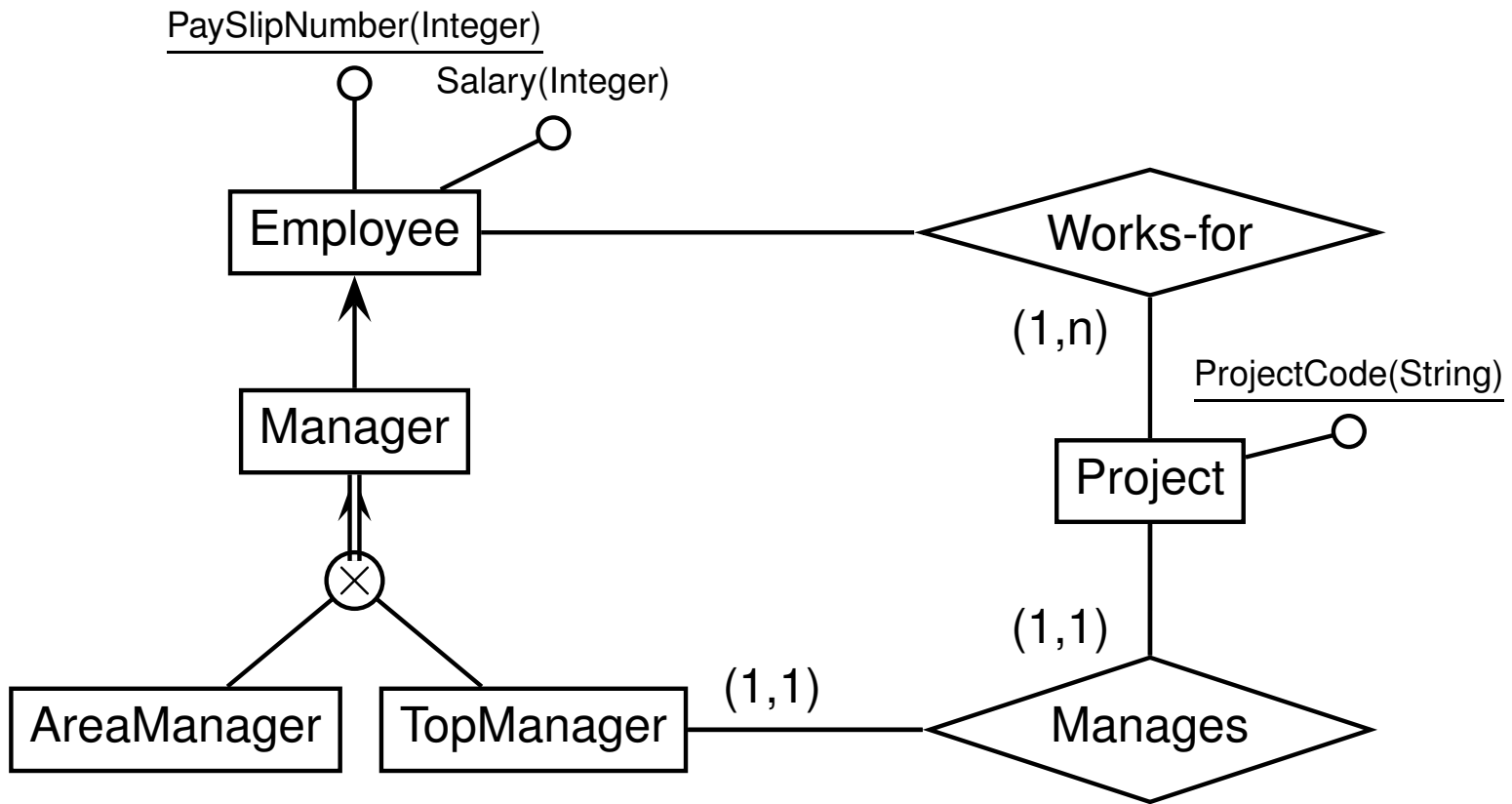
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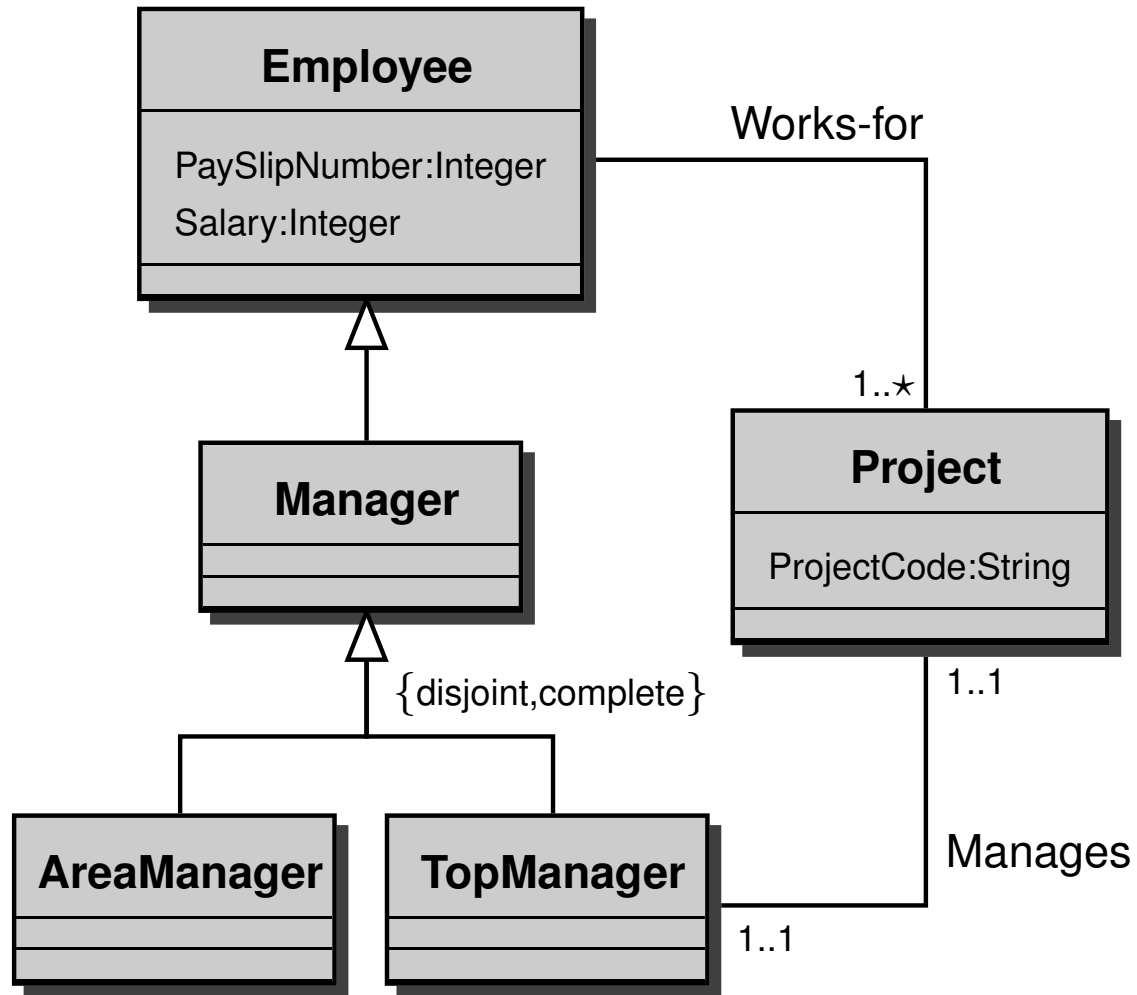
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- The Entity-Relationship conceptual data model and UML Class Diagrams can be considered as ontology languages.

# Entity-Relationship Diagram



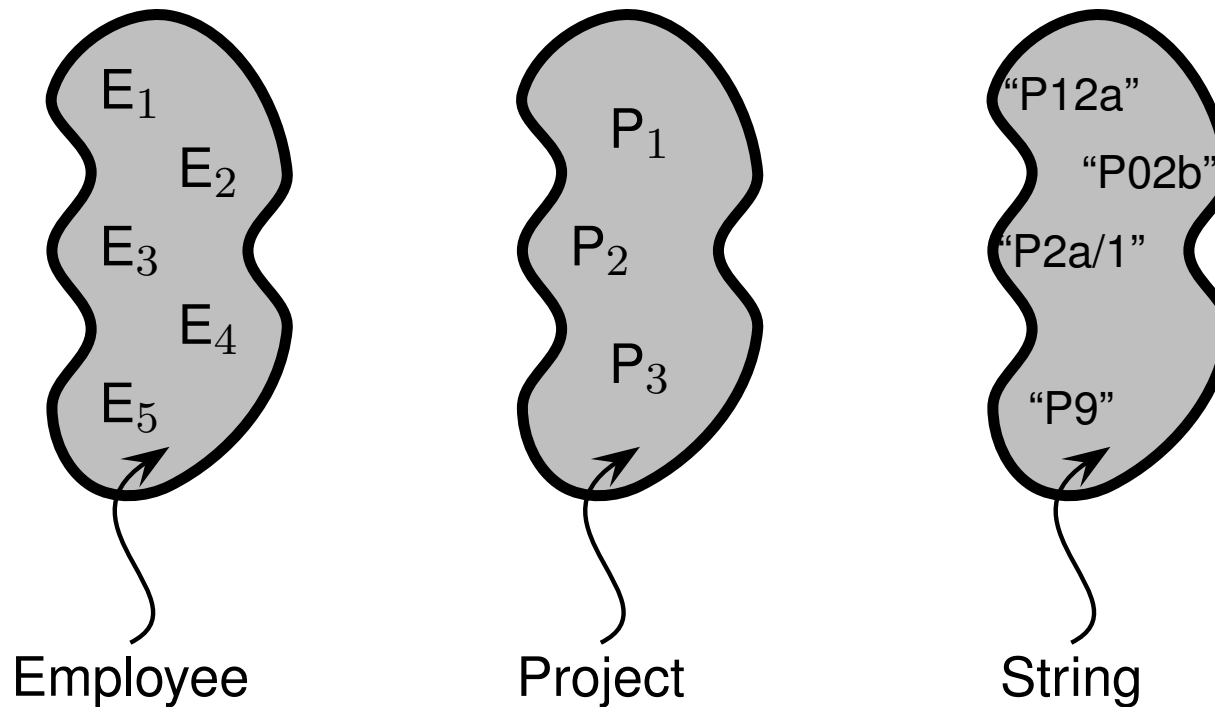
# UML Class Diagram



# Meaning of basic constructs

In a specific legal database:

- An entity is a **set of instances**;
- a n-ary relationship is a **set of n-tuples of instances**;
- an attribute is a **set of pairs of an instance and a domain element**.

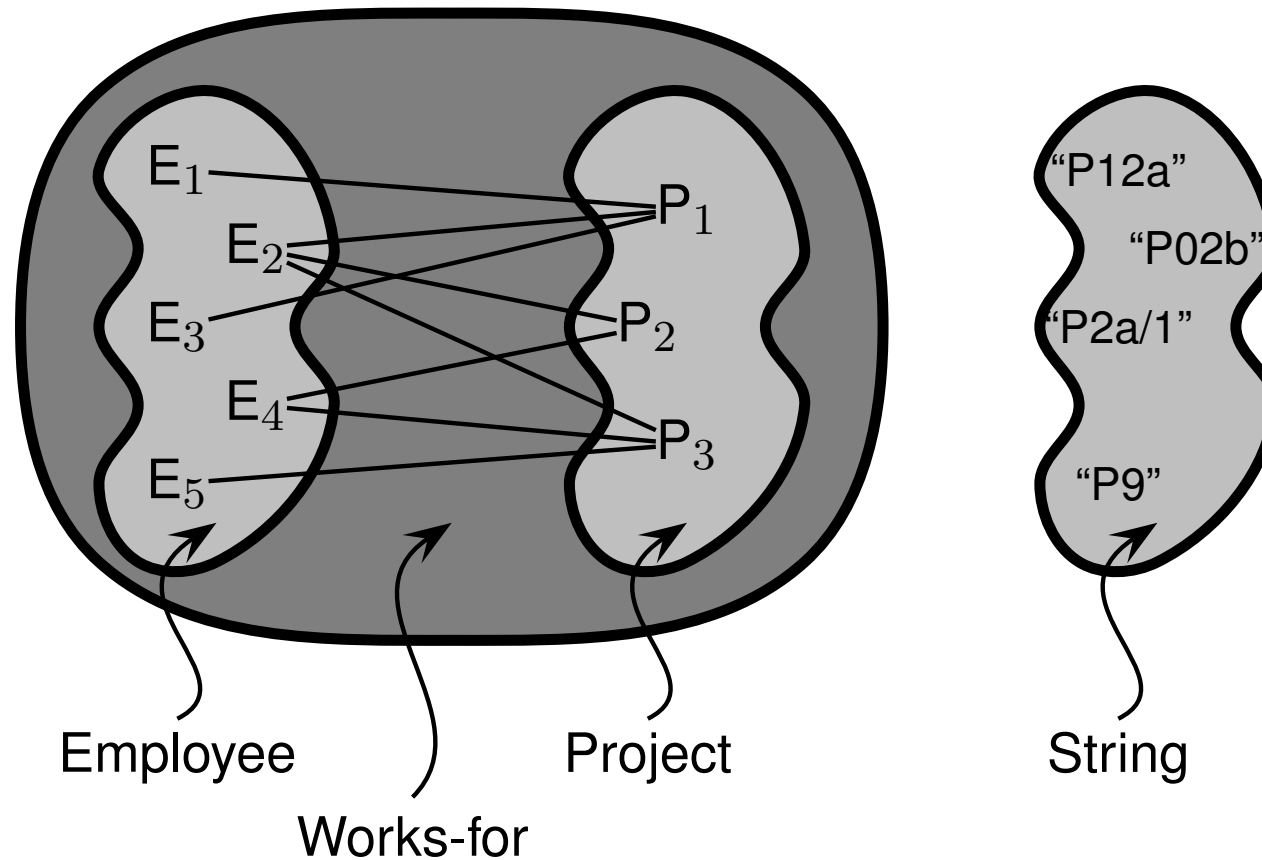




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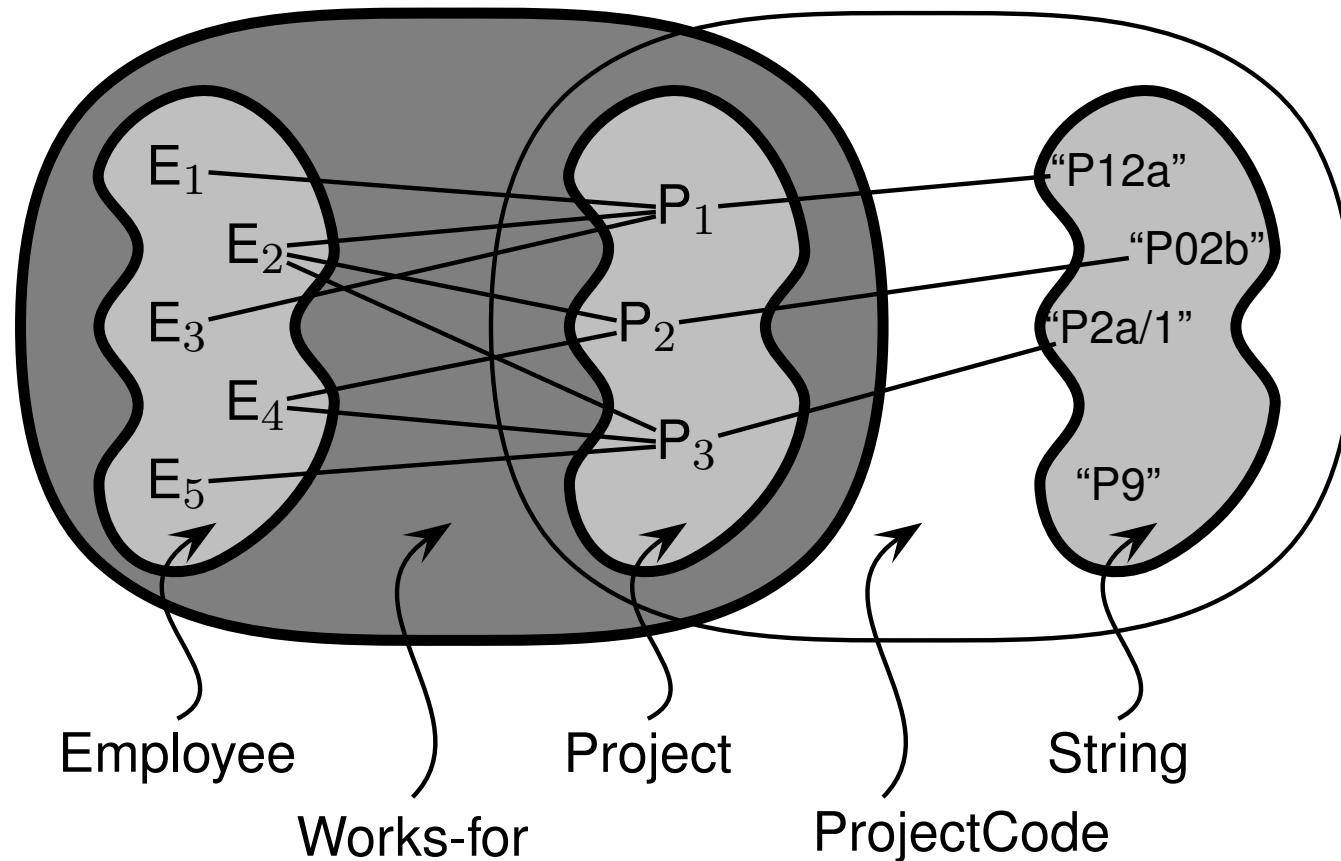
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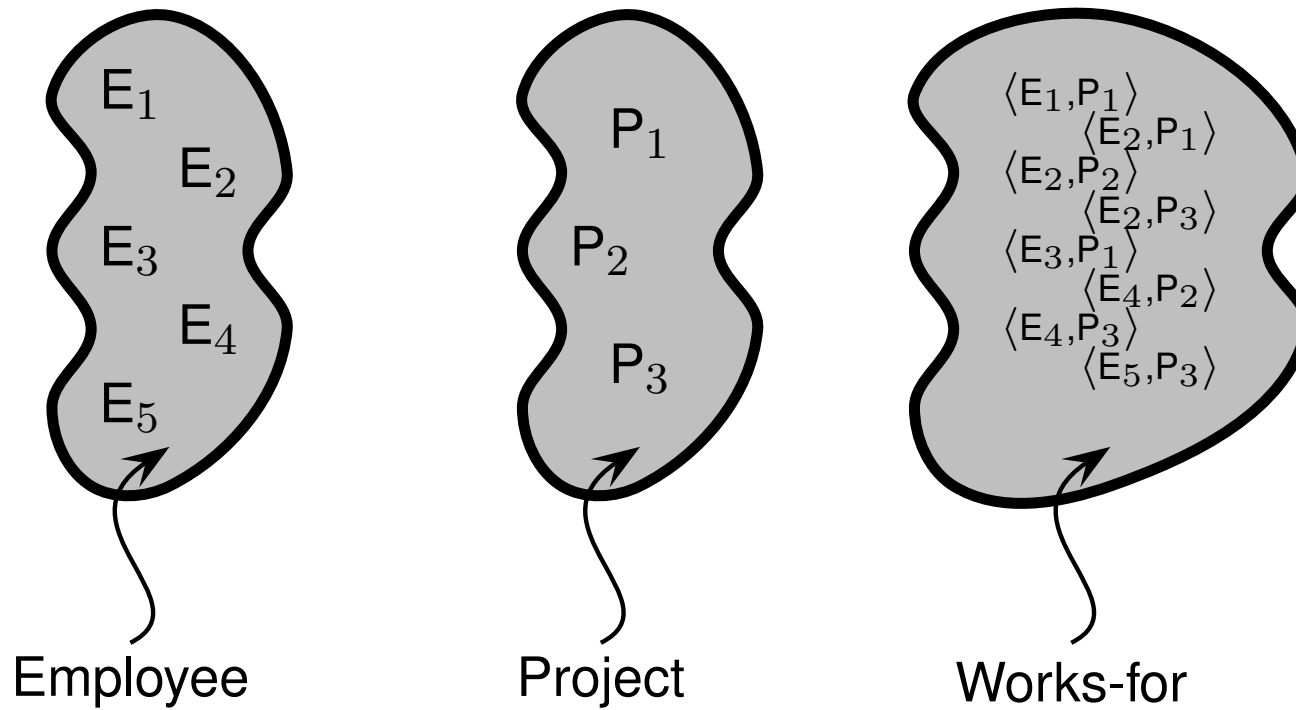
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# Relations as sets of tuples



# The relational representation

Employee

<i>employeeId</i>
E <sub>1</sub>
E <sub>2</sub>
E <sub>3</sub>
E <sub>4</sub>
E <sub>5</sub>

Project

<i>projectId</i>
P <sub>1</sub>
P <sub>2</sub>
P <sub>3</sub>

String

<i>anystring</i>
"P12a"
"P02b"
"P2a/1"
"P9"
...

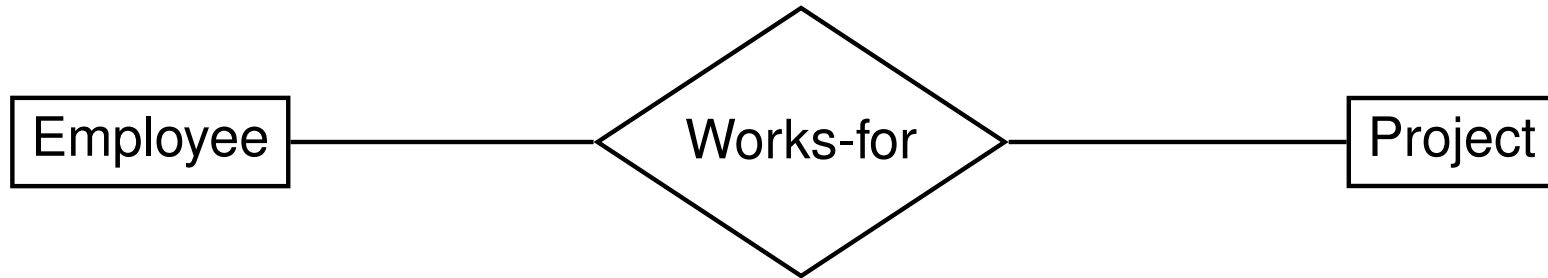
Works-for

<i>employeeId</i>	<i>projectId</i>
E <sub>1</sub>	P <sub>1</sub>
E <sub>2</sub>	P <sub>1</sub>
E <sub>2</sub>	P <sub>2</sub>
E <sub>2</sub>	P <sub>3</sub>
E <sub>3</sub>	P <sub>1</sub>
E <sub>4</sub>	P <sub>2</sub>
E <sub>4</sub>	P <sub>3</sub>
E <sub>5</sub>	P <sub>3</sub>

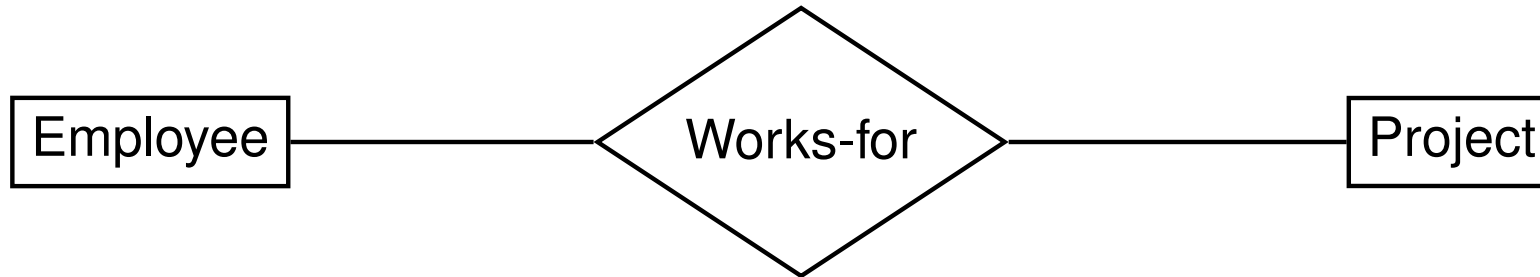
ProjectCode

<i>projectId</i>	<i>pcode</i>
P <sub>1</sub>	"P12a"
P <sub>2</sub>	"P02b"
P <sub>3</sub>	"P2a/1"

# Meaning of Relationships

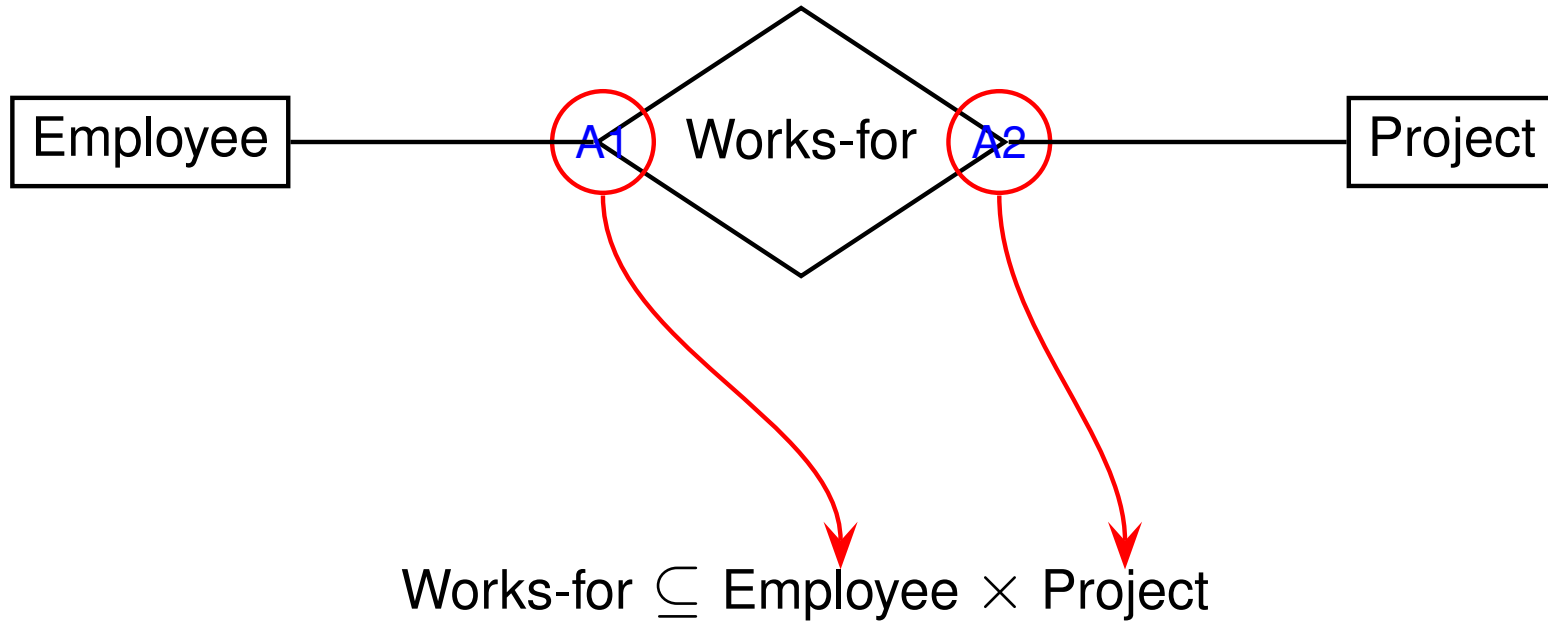


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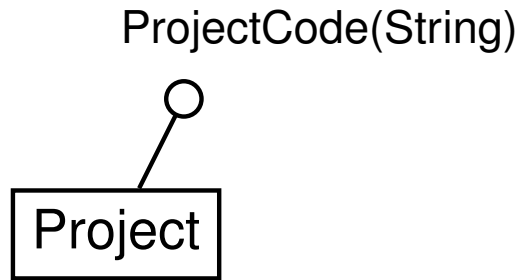


$\text{Works-for} \subseteq \text{Employee} \times \text{Project}$

# Meaning of Relationships

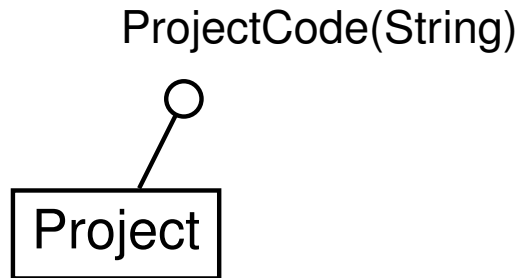


# Meaning of Attributes



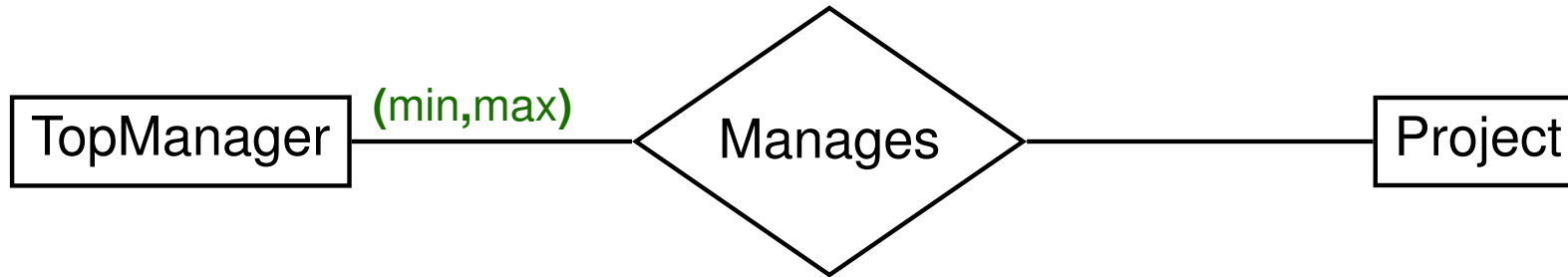


# Meaning of Attributes

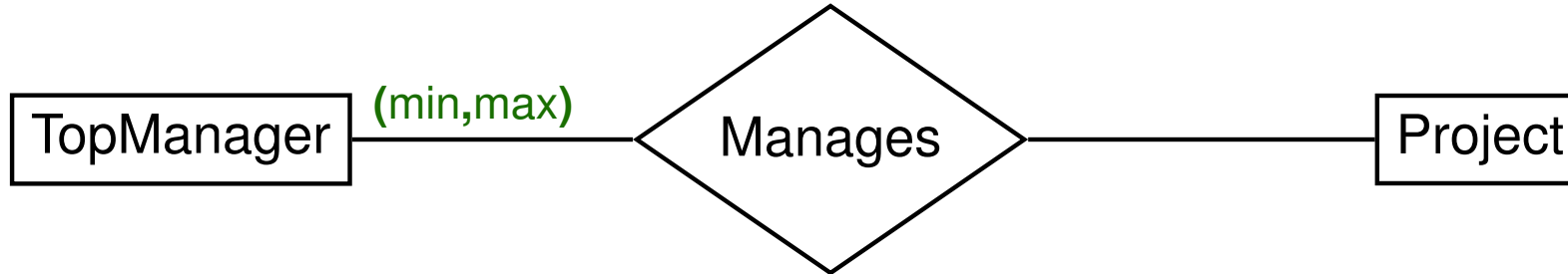


$$\text{Project} \subseteq \{p \mid \#(\text{ProjectCode} \cap (\{p\} \times \text{String})) \geq 1\}$$

# Meaning of Cardinality Constraints



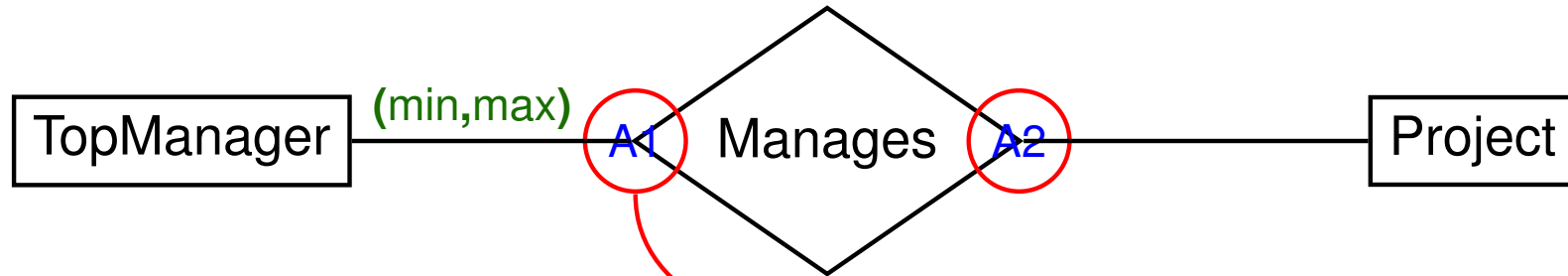
# Meaning of Cardinality Constraints



$$\text{TopManager} \subseteq \{m \mid \text{max} \geq \#(\text{Manages} \cap (\{m\} \times \Omega)) \geq \text{min}\}$$

(where  $\Omega$  is the set of all instances)

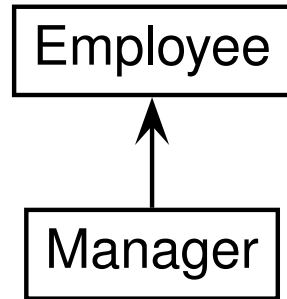
# Meaning of Cardinality Constraints



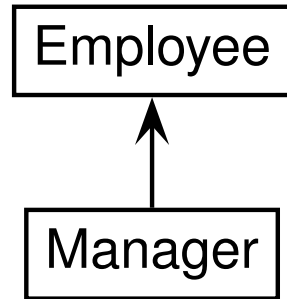
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# Meaning of ISA

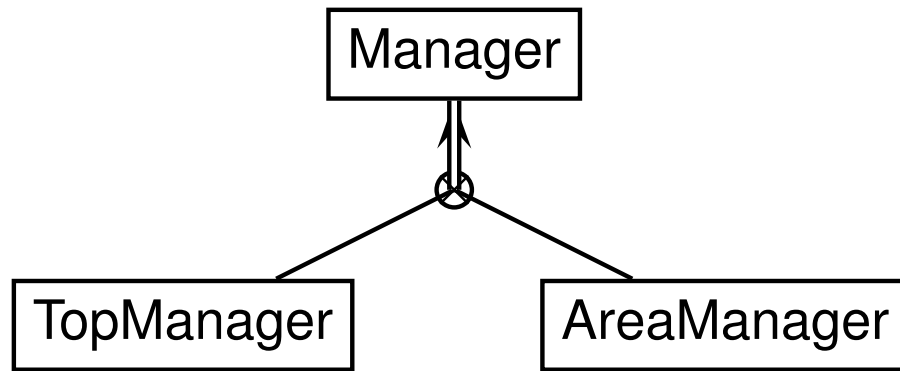


# Meaning of ISA

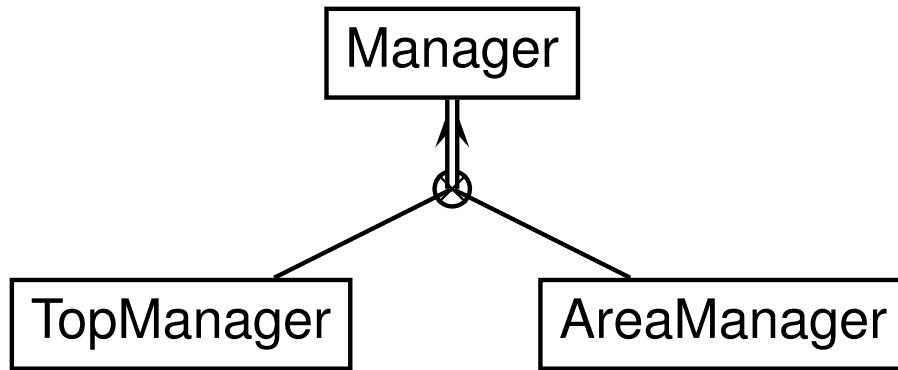


Manager  $\subseteq$  Employee

# Meaning of *disjoint* and *total* constraints



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- *ISA*:  $\text{AreaManager} \subseteq \text{Manager}$
- *ISA*:  $\text{TopManager} \subseteq \text{Manager}$
- *disjoint*:  $\text{AreaManager} \cap \text{TopManager} = \emptyset$
- *total*  $\text{Manager} \subseteq \text{AreaManager} \cup \text{TopManager}$



# Meaning of the initial diagram

Works-for  $\subseteq$  Employee  $\times$  Project

Manages  $\subseteq$  TopManager  $\times$  Project

Employee  $\subseteq \{e \mid \#(\text{PaySlipNumber} \cap (\{e\} \times \text{Integer})) \geq 1\}$

Employee  $\subseteq \{e \mid \#(\text{Salary} \cap (\{e\} \times \text{Integer})) \geq 1\}$

Project  $\subseteq \{p \mid \#(\text{ProjectCode} \cap (\{p\} \times \text{String})) \geq 1\}$

TopManager  $\subseteq \{m \mid 1 \geq \#(\text{Manages} \cap (\{m\} \times \Omega)) \geq 1\}$

Project  $\subseteq \{p \mid 1 \geq \#(\text{Manages} \cap (\Omega \times \{p\})) \geq 1\}$

Project  $\subseteq \{p \mid \#(\text{Works-for} \cap (\Omega \times \{p\})) \geq 1\}$

Manager  $\subseteq$  Employee

AreaManager  $\subseteq$  Manager

TopManager  $\subseteq$  Manager

AreaManager  $\cap$  TopManager =  $\emptyset$

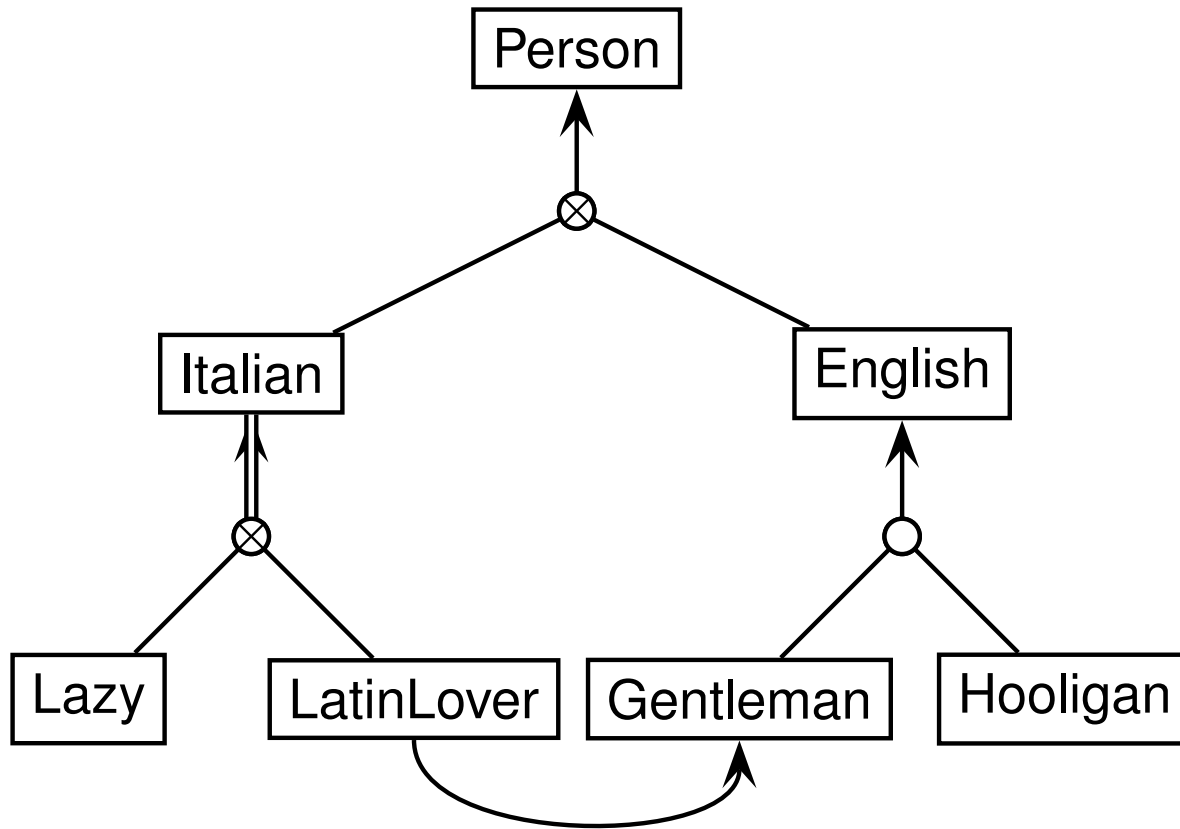
Manager  $\subseteq$  AreaManager  $\cup$  TopManager

# Inference

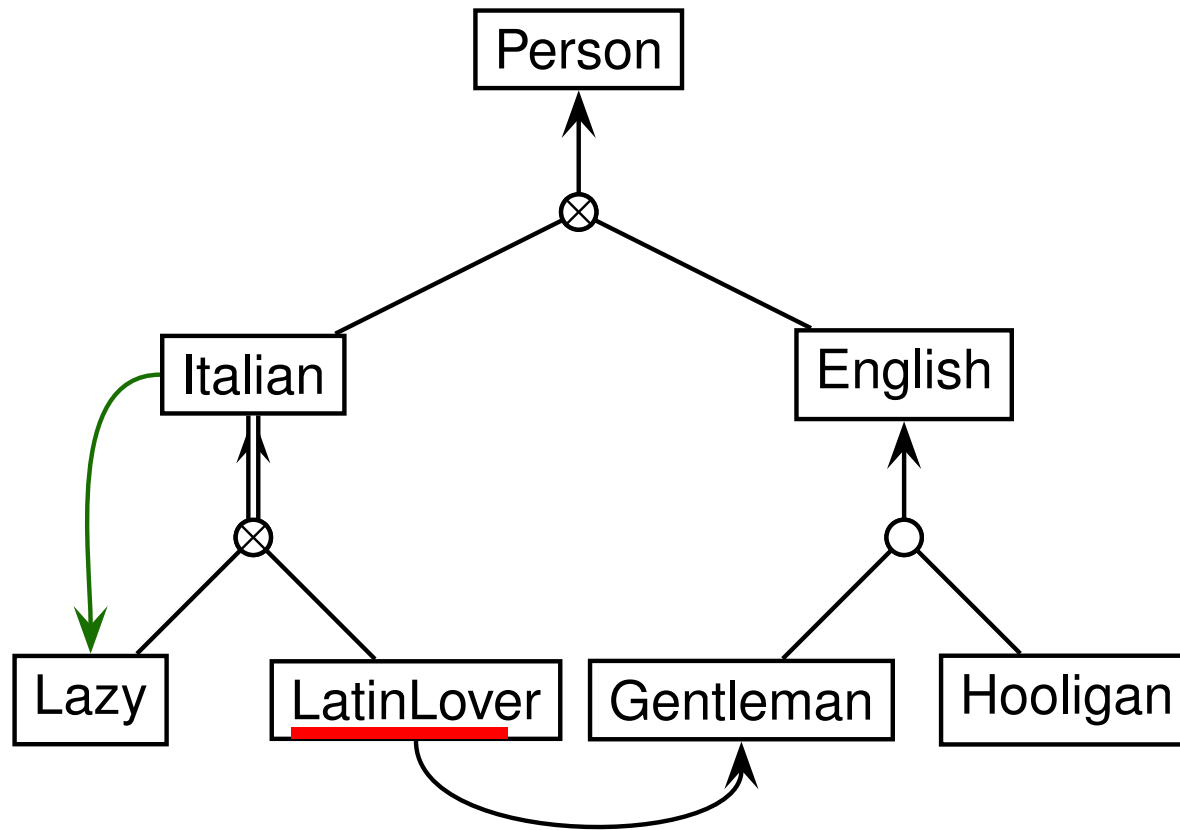
Given a collection of constraints, such as an Entity-Relationship diagram, it is possible that additional constraints can be inferred.

- An entity is **inconsistent** if it denotes the empty set in any legal database.
- An entity is a **subentity** of another entity if the former denotes a subset of the set denoted by the latter in any legal database.
- Two entities are **equivalent** if they denote the same set in any legal database.
- A **stricter** constraint is inferred – e.g., a **cardinality** constraint – if it holds in in any legal database.
- . . .

# Inference



# Inference



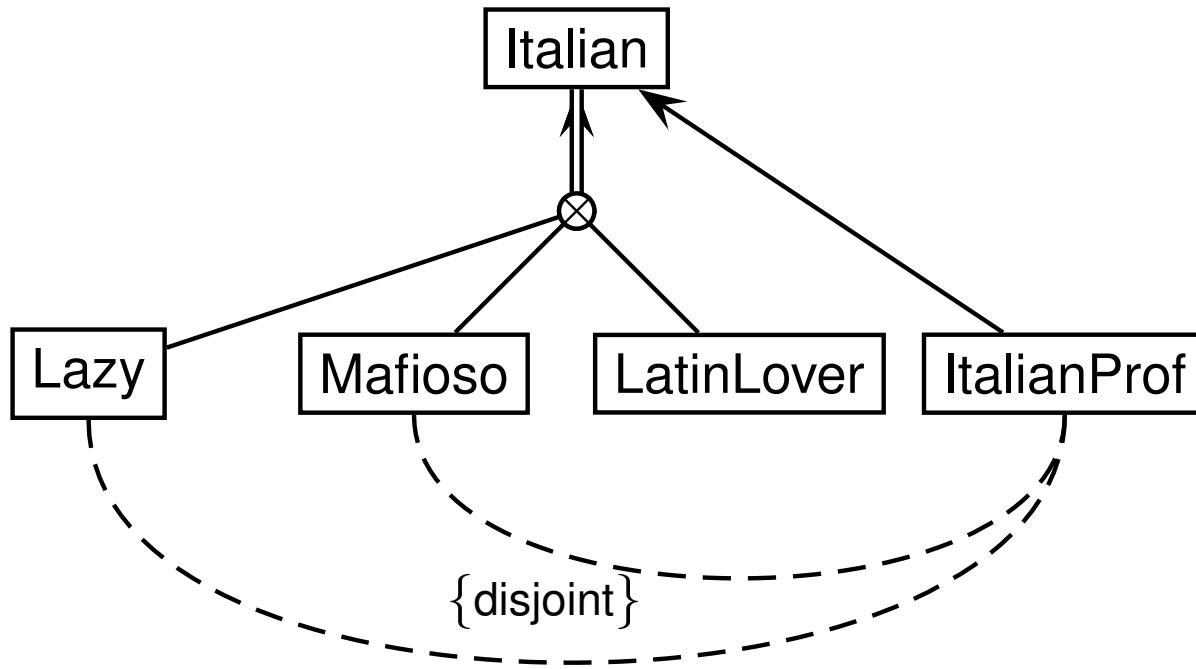
*implies*

LatinLover =  $\emptyset$

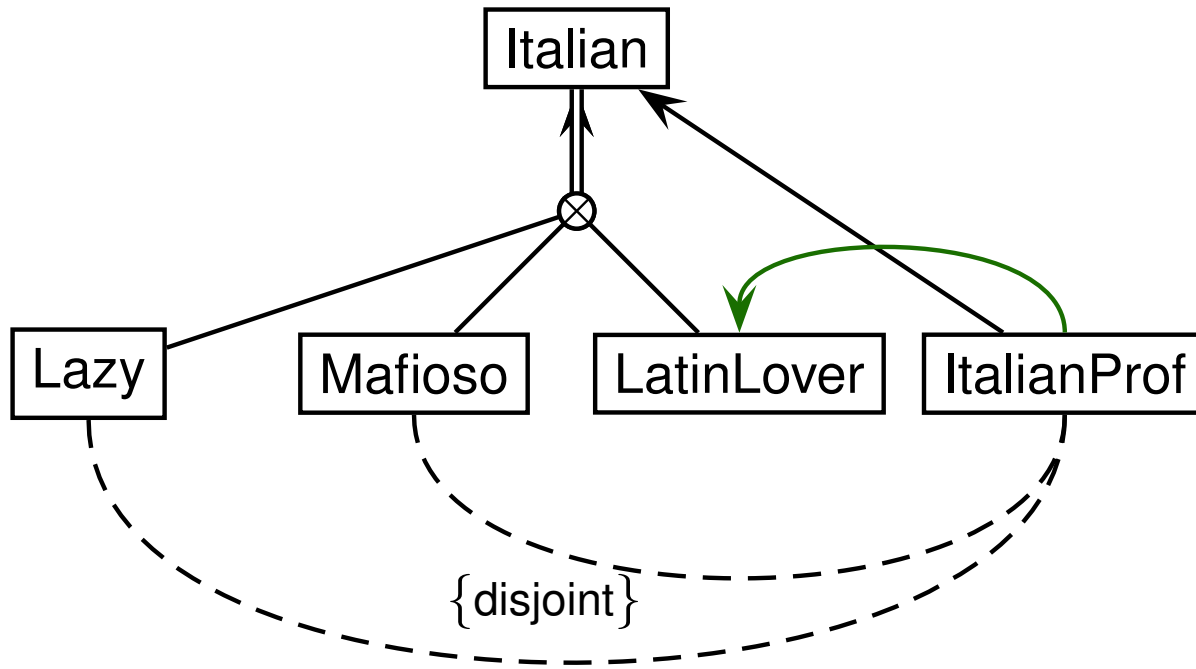
Italian  $\subseteq$  Lazy

Italian  $\equiv$  Lazy

# Reasoning by cases



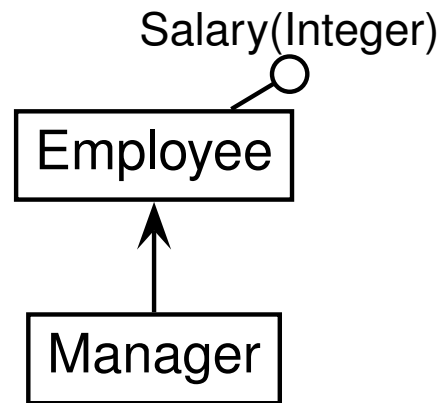
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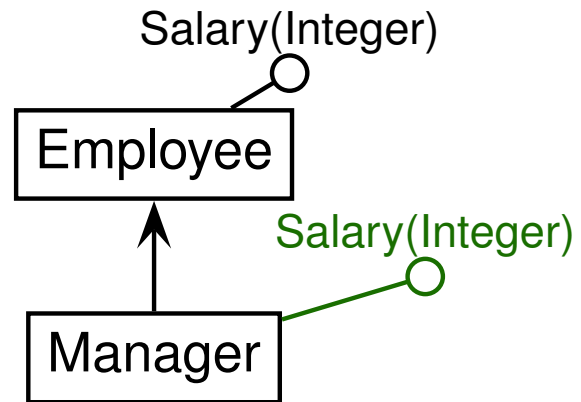
*implies*

ItalianProf  $\subseteq$  LatinLover

# ISA and Inheritance



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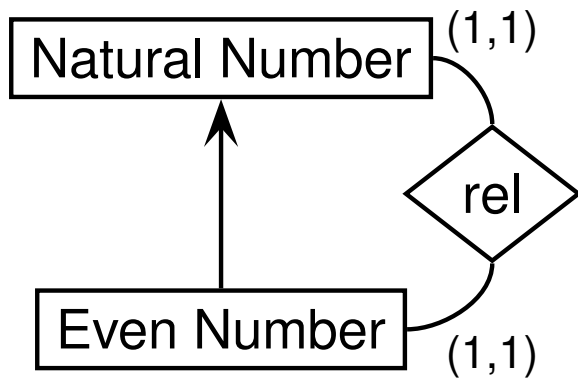


*implies*

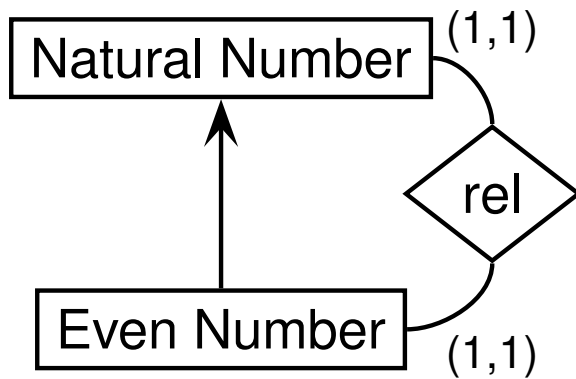
$$\text{Manager} \subseteq \{m \mid \#(\text{Salary} \cap (\{m\} \times \text{Integer})) \geq 1\}$$



# Bijection bewteen Entities



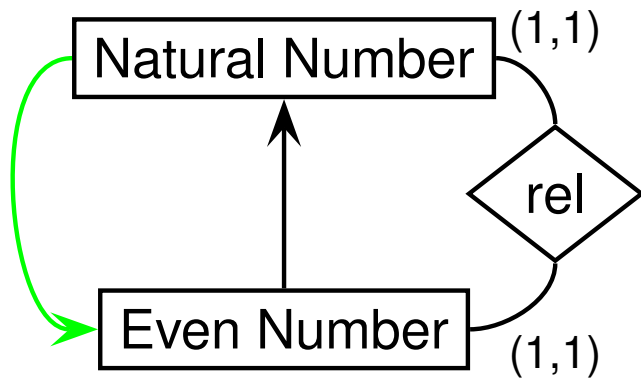
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*implies*

“the entities '*Natural Number*' and '*Even Number*' contain the same number of instances”.

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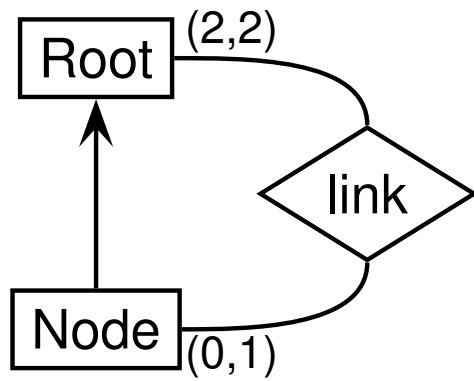


*implies*

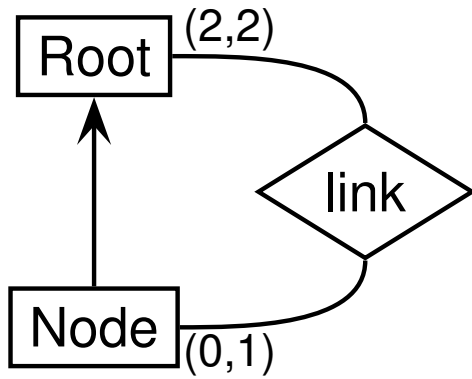
“the entities '*Natural Number*' and '*Even Number*' contain the same number of instances”.

If the domain is finite:       $\text{Natural Number} \equiv \text{Even Number}$

# Infinite Databases



# Infinite Databases



*implies*

“the classes Root and Node contain an infinite number of instances”.