

Propositions

- An interpretation is an assignment of values to all variables.
- A model is an interpretation that satisfies the constraints.
- Often we don't want to just find a model, but want to know what is true in all models.
- A proposition is statement that is true or false in each interpretation.

Why propositions?

- Specifying logical formulae is often more natural than filling in tables
- It is easier to check correctness and debug formulae than tables
- We can exploit the Boolean nature for efficient reasoning
- We need a language for asking queries (of what follows in all models) that may be more complicated than asking for the value of a variable
- It is easy to incrementally add formulae
- It can be extended to infinitely many variables with infinite domains (using logical quantification)

Human's view of semantics

Step 1 Begin with a task domain.

Step 2 Choose atoms in the computer to denote propositions. These atoms have meaning to the KB designer.

Step 3 Tell the system knowledge about the domain.

Step 4 Ask the system questions.

- the system can tell you whether the question is a logical consequence.
- You can interpret the answer with the meaning associated with the atoms.

Role of semantics

In computer:

$light1_broken \leftarrow sw_up$
 $\wedge power \wedge unlit_light1.$

$sw_up.$

$power \leftarrow lit_light2.$

$unlit_light1.$

$lit_light2.$

In user's mind:

- $light1_broken$: light #1 is broken
- sw_up : switch is up
- $power$: there is power in the building
- $unlit_light1$: light #1 isn't lit
- lit_light2 : light #2 is lit

Conclusion: $light1_broken$

- The computer doesn't know the meaning of the symbols
- The user can interpret the symbol using their meaning

Simple language: propositional definite clauses

- An **atom** is a symbol starting with a lower case letter
- A **body** is an atom or is of the form $b_1 \wedge b_2$ where b_1 and b_2 are bodies.
- A **definite clause** is an atom or is a rule of the form $h \leftarrow b$ where h is an atom and b is a body.
- A **knowledge base** is a set of definite clauses

- An **interpretation** I assigns a truth value to each atom.
- A body $b_1 \wedge b_2$ is true in I if b_1 is true in I and b_2 is true in I .
- A rule $h \leftarrow b$ is false in I if b is true in I and h is false in I . The rule is true otherwise.
- A knowledge base KB is true in I if and only if every clause in KB is true in I .

Models and Logical Consequence

- A **model** of a set of clauses is an interpretation in which all the clauses are *true*.
- If KB is a set of clauses and g is a conjunction of atoms, g is a **logical consequence** of KB , written $KB \models g$, if g is *true* in every model of KB .
- That is, $KB \models g$ if there is no interpretation in which KB is *true* and g is *false*.

Simple Example

$$KB = \begin{cases} p \leftarrow q. \\ q. \\ r \leftarrow s. \end{cases}$$

	<i>p</i>	<i>q</i>	<i>r</i>	<i>s</i>	model?
<i>l</i> ₁	<i>true</i>	<i>true</i>	<i>true</i>	<i>true</i>	
<i>l</i> ₂	<i>false</i>	<i>false</i>	<i>false</i>	<i>false</i>	
<i>l</i> ₃	<i>true</i>	<i>true</i>	<i>false</i>	<i>false</i>	
<i>l</i> ₄	<i>true</i>	<i>true</i>	<i>true</i>	<i>false</i>	
<i>l</i> ₅	<i>true</i>	<i>true</i>	<i>false</i>	<i>true</i>	

Simple Example

$$KB = \begin{cases} p \leftarrow q. \\ q. \\ r \leftarrow s. \end{cases}$$

	<i>p</i>	<i>q</i>	<i>r</i>	<i>s</i>	model?
<i>l</i> ₁	<i>true</i>	<i>true</i>	<i>true</i>	<i>true</i>	is a model of <i>KB</i>
<i>l</i> ₂	<i>false</i>	<i>false</i>	<i>false</i>	<i>false</i>	not a model of <i>KB</i>
<i>l</i> ₃	<i>true</i>	<i>true</i>	<i>false</i>	<i>false</i>	is a model of <i>KB</i>
<i>l</i> ₄	<i>true</i>	<i>true</i>	<i>true</i>	<i>false</i>	is a model of <i>KB</i>
<i>l</i> ₅	<i>true</i>	<i>true</i>	<i>false</i>	<i>true</i>	not a model of <i>KB</i>

Simple Example

$$KB = \begin{cases} p \leftarrow q. \\ q. \\ r \leftarrow s. \end{cases}$$

	<i>p</i>	<i>q</i>	<i>r</i>	<i>s</i>	model?
<i>l</i> ₁	<i>true</i>	<i>true</i>	<i>true</i>	<i>true</i>	is a model of <i>KB</i>
<i>l</i> ₂	<i>false</i>	<i>false</i>	<i>false</i>	<i>false</i>	not a model of <i>KB</i>
<i>l</i> ₃	<i>true</i>	<i>true</i>	<i>false</i>	<i>false</i>	is a model of <i>KB</i>
<i>l</i> ₄	<i>true</i>	<i>true</i>	<i>true</i>	<i>false</i>	is a model of <i>KB</i>
<i>l</i> ₅	<i>true</i>	<i>true</i>	<i>false</i>	<i>true</i>	not a model of <i>KB</i>

Which of *p*, *q*, *r*, *s* logically follow from *KB*?

Simple Example

$$KB = \begin{cases} p \leftarrow q. \\ q. \\ r \leftarrow s. \end{cases}$$

	<i>p</i>	<i>q</i>	<i>r</i>	<i>s</i>	model?
<i>l</i> ₁	<i>true</i>	<i>true</i>	<i>true</i>	<i>true</i>	is a model of <i>KB</i>
<i>l</i> ₂	<i>false</i>	<i>false</i>	<i>false</i>	<i>false</i>	not a model of <i>KB</i>
<i>l</i> ₃	<i>true</i>	<i>true</i>	<i>false</i>	<i>false</i>	is a model of <i>KB</i>
<i>l</i> ₄	<i>true</i>	<i>true</i>	<i>true</i>	<i>false</i>	is a model of <i>KB</i>
<i>l</i> ₅	<i>true</i>	<i>true</i>	<i>false</i>	<i>true</i>	not a model of <i>KB</i>

Which of *p*, *q*, *r*, *q* logically follow from *KB*?

$KB \models p$, $KB \models q$, $KB \not\models r$, $KB \not\models s$

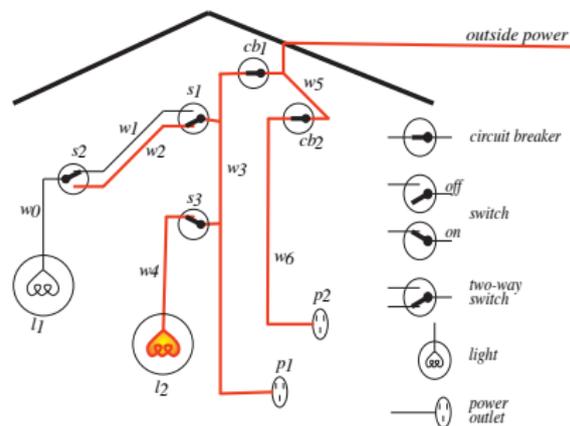
User's view of Semantics

- 1 Choose a task domain: **intended interpretation.**
- 2 Associate an atom with each proposition you want to represent.
- 3 Tell the system clauses that are true in the intended interpretation: **axiomatizing the domain.**
- 4 Ask questions about the intended interpretation.
- 5 If $KB \models g$, then g must be true in the intended interpretation.
- 6 Users can interpret the answer using their intended interpretation of the symbols.

Computer's view of semantics

- The computer doesn't have access to the intended interpretation.
- All it knows is the knowledge base.
- The computer can determine if a formula is a logical consequence of KB.
- If $KB \models g$ then g must be true in the intended interpretation.
- If $KB \not\models g$ then there is a model of KB in which g is false. This could be the intended interpretation.

Electrical Environment



$light_l_1.$ $lit_l_1 \leftarrow live_w_0 \wedge ok_l_1$
 $light_l_2.$ $live_w_0 \leftarrow live_w_1 \wedge up_s_2.$
 $down_s_1.$ $live_w_0 \leftarrow live_w_2 \wedge down_s_2.$
 $up_s_2.$ $live_w_1 \leftarrow live_w_3 \wedge up_s_1.$
 $up_s_3.$ $live_w_2 \leftarrow live_w_3 \wedge down_s_1.$
 $ok_l_1.$ $lit_l_2 \leftarrow live_w_4 \wedge ok_l_2.$
 $ok_l_2.$ $live_w_4 \leftarrow live_w_3 \wedge up_s_3.$
 $ok_cb_1.$ $live_p_1 \leftarrow live_w_3.$
 $ok_cb_2.$ $live_w_3 \leftarrow live_w_5 \wedge ok_cb_1.$
 $live_outside.$ $live_p_2 \leftarrow live_w_6.$
 $live_w_6 \leftarrow live_w_5 \wedge ok_cb_2.$
 $live_w_5 \leftarrow live_outside.$