

Foundations and Challenges of Change in Ontologies and Databases

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Temporal Dynamic Description Logic

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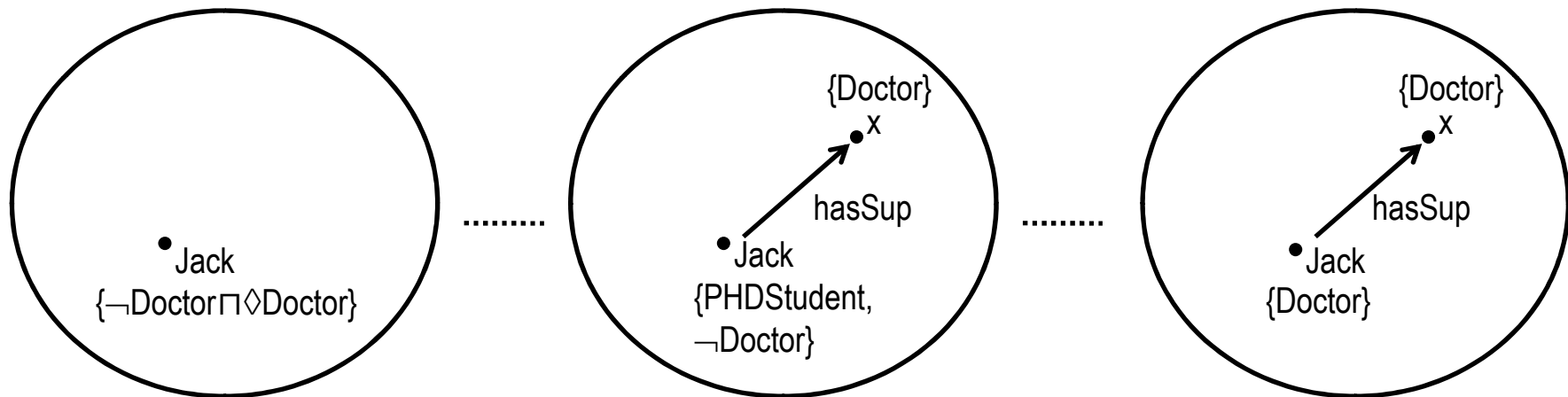
Temporal description logic (TDL)

- For capturing temporal aspects of concepts in ontologies.

$\neg\text{Doctor} \sqcap \diamond\text{Doctor} \sqsubseteq$

$\diamond(\text{PHDStudent} \sqcap \neg\text{Doctor} \sqcap (\text{PHDStudent} \sqcup \text{Doctor}))$

$\text{PHDStudent} \sqsubseteq \exists \text{hasSup}.\text{Doctor}$



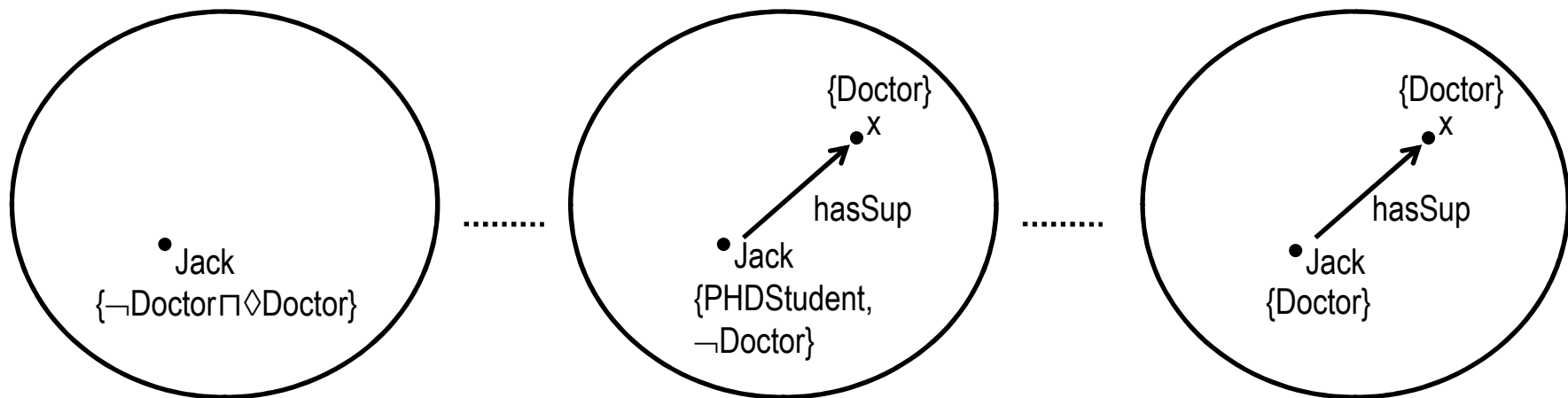
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- Two-dimensional logics [GKWZ03]
 - Temporal description logic
 - Dynamic description logic
 -

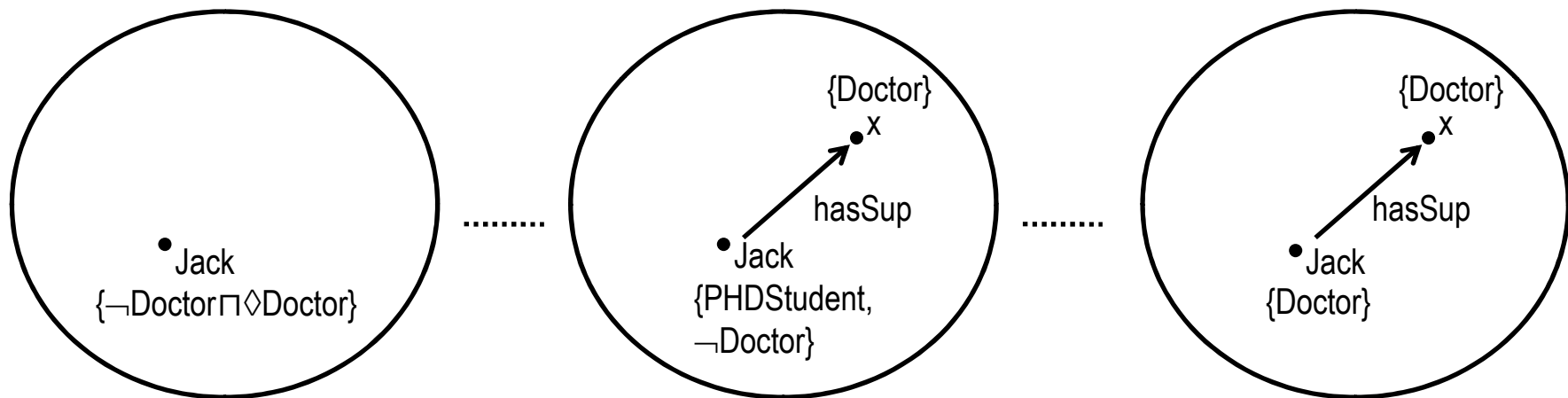
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} joint?

–

Different temporal extensions of DLs

- **Explicit** notion of time or **implicit** time
- **Interval-based** notion of time or **point-based** time
 - **External** representation of time or **internal** representation
- **Linear** time or **branching** time

Different temporal extensions

- Varying DL component: **DL-Lite**, **EL**, **ALC**, **SHOIQ**, ...
- Different choice for applying temporal operators:
concepts, **TBox axioms**, **ABox assertions**
 - $\neg\text{Doctor} \sqcap \diamond\text{Doctor} \sqsubseteq \diamond(\text{PHDStudent} \sqcup \text{Doctor})$
 - $\diamond\Box(\text{Citizen} \sqsubseteq \text{HASVote})$
 - $\text{PHDStudent}(\text{Jack}) \wedge \diamond(\text{PHDStudent}(\text{Jack}) \sqcup \text{Doctor}(\text{Jack}))$
- Additional constraints on concepts and roles:
rigid concepts, **rigid roles**
- interpretation domains: **expanding**, **constant**
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Dozens of combinations!

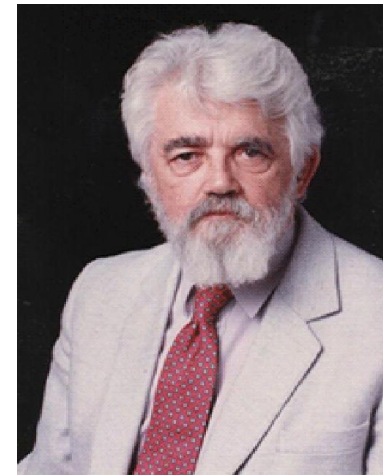
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Reasoning about actions

- Representation and Reasoning about Actions
- Situation Calculus [Mcc63]
- John Mccarthy
 - father of AI, 1956
 - Winner of Turing Award, 1971



• *John Mccarthy (1927-2011)*

Action Formalisms

Based on
propositional
logics

based on PDL
[GL95]

based on LTL
[CGV02]

Based on DL ?

Gap ?

Based on
first- or
higher-order
logics

Situation Calculus

Fluent Calculus

.....

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Action formalism based on DL
[BLM'05]

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DL-Based Action Formalisms

- Background knowledge: RBox, TBox
- States: ABoxes
- Action: $\alpha = (\text{pre}, \text{occ}, \text{post})$
 - **pre**: ABox assertions
 - **occ**: primitive literals
 - **post**: set of conditional post-conditions, φ/ψ
- **Update ABox** after the execution of actions.

Extension of the DL-based action formalism

Basic idea: construct more powerful formalism,
action theory + description logic + dynamic logic

- Background knowledge: RBox, TBox
- Atomic actions: come from Baader et al.'s formalism

$\alpha \equiv (\text{pre}, \text{occ}, \text{post})$

- Complex actions:

$\pi, \pi' ::= \alpha \mid \phi? \mid \pi \cup \pi' \mid \pi; \pi' \mid \pi^*$

- Formulas:

$\phi, \psi ::= C(p) \mid R(p, q) \mid \langle \pi \rangle \phi \mid [\pi] \phi \mid \neg \phi \mid \phi \vee \psi \mid \phi \wedge \psi$

- *Dynamic description logic* $DDL(X@)$

X : DLs ranging from ALCO to ALCHOIQ ,

$X@$: extension of X with the @ constructor.

Features of DDL(X@) (1/3)

(1) Complex actions can be constructed

- TBox:

$Customer \equiv Person \sqcap \exists \text{holds.CreditCard}$

$VIPcustomer \equiv Customer \sqcap \geq 10 \text{ boughr.}(Book \sqcup CD)$

- Atomic Actions:

$buybook(a,b) \equiv (\{Customer(a), Book(b)\}, \{ \};$
 $\{Instore(b)/\neg Instore(b), Instore(b)/bought(a,b)\})$

$order(b) \equiv (\{(Book \sqcup CD)(b)\}, \{ \};$
 $\{ \neg Instore(b)/Instore(b)\})$

- Complex Action:

$VIPbuybook(a,b) \equiv VIPcustomer(a)? ;$
 $((Instore(b)? ; buybook(a,b)) \cup$
 $(\neg Instore(b)? ; order(b); buybook(a,b)))$

Features of DDL($X@$) (2/3)

(2) Properties on (complex) actions can be described directly

- **necessary conditions** for the execution of (complex) actions
 - $\langle \text{VIPbuybook}(a,b) \rangle \text{true} \rightarrow (\text{VIPcustomer}(a) \wedge \text{Book}(b))$
 - $\langle \text{VIPbuybook}(a,b) \rangle \text{true} \rightarrow \text{Instore}(b)$
- **results** on the execution of actions
 - $[\text{VIPbuybook}(a,b)]\text{bought}(a,b)$
 - $[\text{buybook}(a,b)]\text{bought}(a,b)$

Features of DDL($X@$) (3/3)

(3) Reasoning problems on actions be reduced to the satisfiability problem of formulas

- Executability of actions
- Projection problem
- Consistency/realizability of actions
 - whether a given action makes sense w.r.t. the knowledge base
buybook(a1,b); buybook(a2,b)
- Satisfiability problem
 - a Tableau decision algorithm is provided.
 - the complexity upper-bound is
 - EXPSpace if $X \in \{ALCO, ALCHO, ALCOQ, ALCHOQ\}$,
 - N2EXPTIME if $X \in \{ALCOI, ALCHOI, ALCOIQ, ALCHOIQ\}$.

Temporal extension of DDL($X@$)

To investigate temporal properties of actions.

Approach:

- the ongoing of time is embodied as the execution of atomic actions (time units)
- two temporal assertions are introduced:

$\phi, \psi ::= C(p) \mid R(p,q) \mid \langle \pi \rangle \phi \mid [\pi] \phi \mid \neg \phi \mid \phi \vee \psi \mid \mathbf{E}(\phi \mathbf{U}^\pi \psi) \mid \mathbf{A}(\phi \mathbf{U}^\pi \psi)$

$\mathbf{E}(\phi \mathbf{U}^\pi \psi)$: there exists some path of π such that “ ϕ until ψ ” holds.

$\mathbf{A}(\phi \mathbf{U}^\pi \psi)$: “ ϕ until ψ ” holds in any path of π .

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$$\mathbf{EX} \phi =_{\text{def}} \bigvee_{\alpha \in N_A} \langle \alpha \rangle \phi$$

$$\mathbf{E}(\phi \mathbf{U} \psi) =_{\text{def}} \mathbf{E}(\phi \mathbf{U}^{(\alpha_1 \cup \dots \cup \alpha_n)^*} \psi)$$

$$\mathbf{A}(\phi \mathbf{U} \psi) =_{\text{def}} \mathbf{A}(\phi \mathbf{U}^{(\alpha_1 \cup \dots \cup \alpha_n)^*} \psi)$$

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$$\mathbf{EF} \phi =_{\text{def}} \mathbf{E}(\text{true} \mathbf{U} \phi)$$

$$\mathbf{AF} \phi =_{\text{def}} \mathbf{A}(\text{true} \mathbf{U} \phi)$$

$$\mathbf{EG} \phi =_{\text{def}} \neg \mathbf{AF}(\neg \phi)$$

$$\mathbf{AG} \phi =_{\text{def}} \neg \mathbf{EF}(\neg \phi)$$

$$\mathbf{AX} \phi =_{\text{def}} \neg \mathbf{EX}(\neg \phi)$$

Description example of TDDL($X@$)

- **liveness property**: good things will eventually happen.

EF((\exists bought $^-$.Customer)(b))

E(Instore(b) **U**^{VIPbuybook(a,b)} \neg Instore(b))

- **safety property**: bad things will never happen.

AG \neg (≥ 2 bought $^-$.Customer)(b))

AG (Instore(b) \vee (\exists bought $^-$.Customer)(b))

- Reduced to satisfiability problem of formulas.
- A **Tableau decision algorithm** is provided.

Limitation of DDL($X@$)/TDDL($X@$)

- TBox:
 - only concept definitions, no GCIs
 - acyclic
- RBox:
 - on transitive property
- Atomic action:
 - no defined concept name occurring in the effect set *post*.

Why?

- difficulty of ABox updating.

Difficulty of ABox updating

Example.

- TBox:

$\text{Trans}(R), \quad A \sqsubseteq \exists R.A, \quad A \sqcap B \sqsubseteq \perp, \quad B \sqsubseteq \forall R.B$

- ABox:

$A(a)$

- Update or new information:

$(\exists R.B)(a)$

Some results on ABox update

Assumptions	DLs	Approach	References
Acyclic TBox; no defined concept names occurring in U	ALC~ALC QIO	PMA semantics & only primitive concept names are counted when measuring distance.	LLMW06, LLMW11
	DL-Lite _F	PMA semantics.	GLPR06, GLPR07
	DL-Lite _R ^{pr}	Both revision and update. Based on $fcl_{\mathcal{T}}(A)$.	KZ11, KZC13
	DL-Lite _{FR}	Based on $cl_{\mathcal{T}}(A)$.	CKNZ10

Thank you!

