

Belief Revision in Description Logic

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Bolzano, FCCOD, 2014

Motivation

- Study the dynamics of ontologies, specially “OWL-like” DL ontologies.
- AGM Belief Revision deals with the problem of adding/removing information in a consistent way.
- AGM is most commonly applied to propositional classical logic and cannot be directly used with DLs.
- How can we adapt AGM so that it can deal with interesting DLs?

In this work

- Show reasons why AGM fails to apply to DLs.
- Adapt Contraction (easy).
- Adapt Revision (less easy).

Outline of the Talk

- 1 Motivation
- 2 The AGM paradigm
- 3 Contraction and DLs
- 4 Revision and DLs
- 5 Conclusions and Future Work

AGM Belief Revision

Three operations defined to deal with knowledge base dynamics:

- **Expansion** - adding knowledge (possibly inconsistent)

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- **Expansion** - adding knowledge (possibly inconsistent)
- **Contraction** - removing knowledge
- **Revision** - adding knowledge consistently

Revision usually defined in terms of contraction:

$$K * \alpha = (K - \neg\alpha) + \alpha$$

AGM Theory

For contraction and revision:

- **Rationality Postulates**

AGM Theory

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AGM Assumptions: Tarskian, Compact, Deduction Theorem, Supraclassical.

AGM contraction

(closure) $K - \alpha = Cn(K - \alpha)$

(success) If $\alpha \notin Cn(\emptyset)$ then $\alpha \notin K - \alpha$

(inclusion) $K - \alpha \subseteq K$

(vacuity) If $\alpha \notin K$ then $K - \alpha = K$

(recovery) $K \subseteq K - \alpha + \alpha$

(extensionality) If $Cn(\alpha) = Cn(\beta)$ then
 $K - \alpha = K - \beta$

Applying to DL

- AGM cannot be applied to every logic. In particular it can not be applied to SHIF and SHOIN. [Flouris 2006]
- Solution: substitute recovery by relevance
(relevance) If $\beta \in K \setminus K - \alpha$, then there is K' s. t.
 $K - \alpha \subseteq K' \subseteq K$ and $\alpha \notin Cn(K')$, but $\alpha \in Cn(K' \cup \{\beta\})$.
- Good property: AGM assumptions + 5 postulates \Rightarrow recovery and relevance are equivalent.

Results - contraction

Representation Theorem [RW06]

If the underlying logic is tarskian and compact, partial meet contraction is equivalent to the AGM postulates with relevance instead of recovery.

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Can we do the same for revision???

AGM Revision

(closure) $K * \alpha = Cn(K * \alpha)$

(success) $\alpha \in K * \alpha$

(inclusion) $K * \alpha \subseteq K + \alpha$

(vacuity) If $K + \alpha$ is consistent then $K * \alpha = K + \alpha$

(consistency) If α is consistent then $K * \alpha$ is consistent.

(extensionality) If $Cn(\alpha) = Cn(\beta)$ then
 $K * \alpha = K * \beta$

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Definition (Revision without negation)

$$K *_{\gamma} \alpha = \bigcap \gamma(K \downarrow \alpha) + \alpha$$

where γ selects at least one element of $K \downarrow \alpha$.

Properties

- 1 Inconsistent explosion: Whenever K is inconsistent, then for all formulas α , $\alpha \in Cn(K)$
- 2 Distributivity: For all sets of formulas X, Y and W ,
$$Cn(X \cup (Cn(Y) \cap Cn(W))) = Cn(X \cup Y) \cap Cn(X \cup W)$$

Properties

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Representation Theorem [RW09]

If the logic is monotonic and compact and satisfies Inconsistent explosion and Distributivity, then $*$ is a revision without negation iff it satisfies closure, success, inclusion, consistency, relevance and uniformity.

(uniformity) If for all $K' \subseteq K$, $K' \cup \{\alpha\}$ is inconsistent iff $K' \cup \{\beta\}$ is inconsistent then $K \cap K * \alpha = K \cap K * \beta$

Which Logics Satisfy Distributivity?

- Classical logic does.
- But what about DLs?
 - \mathcal{ALC} does not.
 - \mathcal{ALC} with empty \mathcal{ABox} does.
 - not many more...

New characterisation

Representation Theorem [RW14]

If the logic is monotonic and compact and satisfies Inconsistent explosion and ~~Distributivity~~, then $*$ is a revision without negation iff it satisfies closure, success, strong inclusion, consistency, relevance and uniformity.

(strong inclusion) $K * \alpha \subseteq (K \cap K * \alpha) + \alpha$

In classical logics this postulate is equivalent to inclusion.

What was done

- Adapted AGM to DLs
 - Contraction - only 1 postulate changed
 - Revision - Contraction and postulates
- Provided representation results.

What we want to do

- Study other forms of revision for DLs avoiding negation.
- Apply the solutions to other fragments
 - Horn
 - ???