

**Foundations and Challenges of Change in Ontologies and Databases**

**Bozen-Bolzano, Italy**

**29-31 January 2014**

# **Temporal Dynamic Description Logic**

**Liang Chang**

**changl.guet@gmail.com**

**liang.chang@manchester.ac.uk**

# About Me

- **2005.09-2008.07**
  - PhD in Computer Science
  - Institute of Computing Technology, Chinese Academy of Sciences
- **2008.08-present**
  - Associate Professor, School of Computer Science, Guilin University of Electronic Technology
  - Deputy Director, Key Lab. of Trusted Software of Guangxi Province
- **2013.08-present**
  - Visiting Academic
  - School of Computer Science, University of Manchester

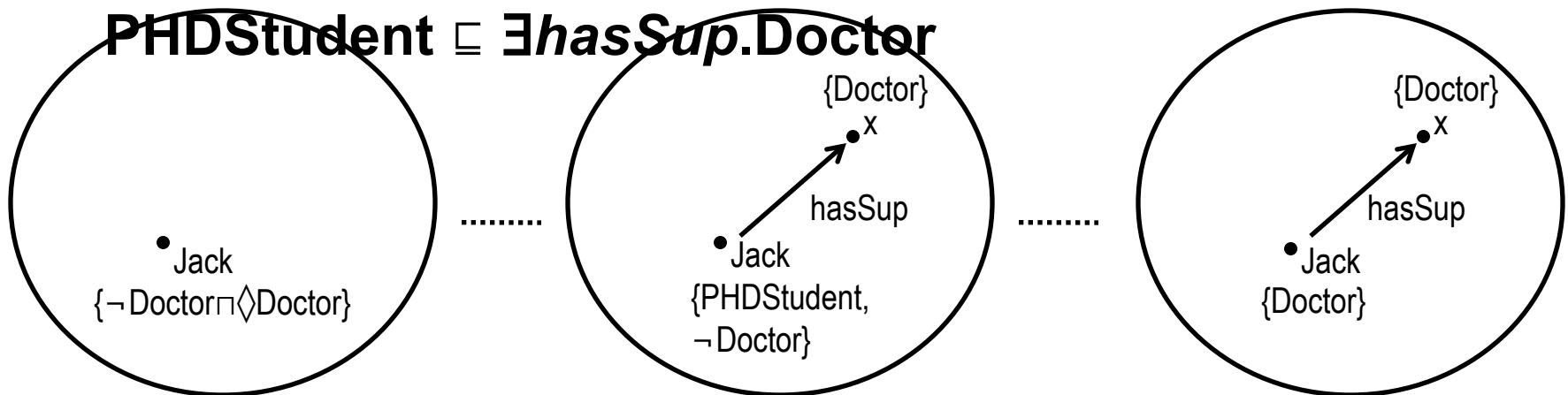
# Temporal description logic (TDL)

- For capturing temporal aspects of concepts in ontologies.

$\neg \text{Doctor} \sqcap \diamond \text{Doctor} \sqsubseteq$

$\diamond (\text{PHDStudent} \sqcap \neg \text{Doctor} \sqcap (\text{PHDStudent} \boxtimes \text{Doctor}))$

$\text{PHDStudent} \sqsubseteq \exists \text{hasSup}.\text{Doctor}$



- Two-dimensional logics [GrWZ03]
- Temporal description logic
- Dynamic description logic

joint?

# Different temporal extensions of DLs

- **Explicit** notion of time or **implicit** time
- **Interval-based** notion of time or **point-based** time
  - **External** representation of time or **internal** representation
- **Linear** time or **branching** time

# Different temporal extensions

- Varying DL component: **DL-Lite**, **EL**, **ALC**, **SHOIQ**, ...
- Different choice for applying temporal operators:  
**concepts**, **TBox axioms**, **ABox assertions**
  - $\neg\text{Doctor} \sqcap \diamond\text{Doctor} \sqsubseteq \diamond(\text{PHDStudent} \sqcap \text{Doctor})$
  - $\diamond\text{Apple}(\text{Citizen} \sqsubseteq \text{HASVote})$
  - $\text{PHDStudent}(\text{Jack}) \wedge \diamond(\text{PHDStudent}(\text{Jack}) \sqcap \text{Doctor}(\text{Jack}))$
- Additional constraints on concepts and roles:  
**rigid concepts**, **rigid roles**
- interpretation domains: **expanding**, **constant**
- .....

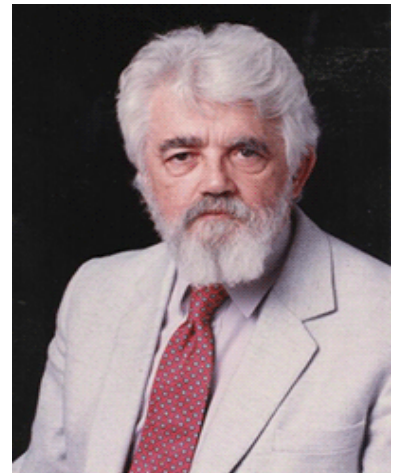
# Different temporal extensions

- **Explicit** notion of time or **implicit time**
- **Interval-based** notion of time or **point-based time**
  - **External** representation of time or **internal** representation
- **Linear** time or **branching time**
- Varying DL component: **DL-Lite**, **EL**, **ALC**, **SHOIQ**, ...
- Different choice for applying temporal operators:  
**concepts**, **TBox axioms**, **ABox assertions**
- Additional constraints on concepts and roles:  
**rigid concepts**, **rigid roles**
- interpretation domains: **expanding**, **constant**
- .....

**Dozens of combinations!**

# Reasoning about actions

- Representation and Reasoning about Actions
- Situation Calculus [Mcc63]
- John Mccarthy
  - father of AI, 1956
  - Winner of Turing Award, 1971



• *John Mccarthy (1927-2011)*

# Action Formalisms

Based on  
propositional  
logics

based on PDL  
[GL95]

based on LTL  
[CGV02]

Based on  
DL ?

Based on first-  
or higher-  
order logics

Situation Calculus

Fluent Calculus

.....

Gap ?

Action formalism based on  
DL [BLM<sup>+</sup>05]



# DL-Based Action Formalisms

- Background knowledge: RBox, TBox
- States: ABoxes
- Action:  $\alpha = (\text{pre}, \text{occ}, \text{post})$ 
  - **pre:** ABox assertions
  - **occ:** primitive literals
  - **post:** set of conditional post-conditions,  $\varphi/\psi$
- **Update ABox** after the execution of actions.

# Extension of the DL-based action formalism

$(M, w) \models \langle \pi \rangle \varphi$  iff some state  $w' \in W$  exists with  $(w, w') \in T(\pi)$  and  $(M, w') \models \varphi$ .

$(M, w) \models [\pi] \varphi$  iff for every state  $w' \in W$ : if  $(w, w') \in T(\pi)$  then  $(M, w') \models \varphi$ .

- Background knowledge: RBox, TBox
- Atomic actions: come from Baader et al.'s formalism

$\alpha \equiv (\text{pre}, \text{occ}, \text{post})$

- Complex actions:

$\pi, \pi' ::= \alpha \mid \phi? \mid \pi \cup \pi' \mid \pi; \pi' \mid \pi^*$

- Formulas:

$\phi, \psi ::= C(p) \mid R(p, q) \mid \langle \pi \rangle \phi \mid [\pi] \phi \mid \neg \phi \mid \phi \vee \psi \mid \phi \wedge \psi$

*Dynamic description logic  $DDL(X@)$*

*X: DLs ranging from ALCO to ALCHOIQ ,*

*X@: extension of X with the @ constructor.*

# Features of DDL(X@) (1/3)

## (1) Complex actions can be constructed

- TBox:

$Customer \equiv Person \sqcap \exists \text{holds.CreditCard}$

$VIPcustomer \equiv Customer \sqcap \geq 10 \text{ boughr.}(Book \sqcup CD)$

- Atomic Actions:

$buybook(a,b) \equiv ( \{Customer(a), Book(b)\}, \{ \};$

$\{Instore(b)/\neg Instore(b), Instore(b)/bought(a,b)\} )$

$order(b) \equiv ( \{(Book \sqcup CD)(b)\}, \{ \};$

$\{\neg Instore(b)/Instore(b)\} )$

- Complex Action:

$VIPbuybook(a,b) \equiv VIPcustomer(a)? ;$

$( (Instore(b)? ; buybook(a,b) ) \cup$

$(\neg Instore(b)? ; order(b); buybook(a,b)) )$

# Features of DDL(X@) (2/3)

## (2) Properties on (complex) actions can be described directly

- **necessary conditions** for the execution of (complex) actions

$\langle \text{VIPbuybook}(a,b) \rangle \text{true} \rightarrow (\text{VIPcustomer}(a) \wedge \text{Book}(b))$  ✓

$\langle \text{VIPbuybook}(a,b) \rangle \text{true} \rightarrow \text{Instore}(b)$  ✗

- **results** on the execution of actions

$[\text{VIPbuybook}(a,b)] \text{bought}(a,b)$  ✓

$[\text{buybook}(a,b)] \text{bought}(a,b)$  ✗

# Features of DDL( $X@$ ) (3/3)

## (3) Reasoning problems on actions be reduced to the satisfiability problem of formulas

- Executability of actions
- Projection problem
- Consistency/realizability of actions
  - whether a given action makes sense w.r.t. the knowledge base  
buybook(a1,b); buybook(a2,b)      ✗
- Satisfiability problem
  - a Tableau decision algorithm is provided.
  - the complexity upper-bound is
    - EXPSPACE if  $X \in \{\text{ALCO}, \text{ALCHO}, \text{ALCOQ}, \text{ALCHOQ}\}$ ,
    - N2EXPTIME if  $X \in \{\text{ALCOI}, \text{ALCHOI}, \text{ALCOIQ}, \text{ALCHOIQ}\}$ .

# Temporal extension of DDL( $X@$ )

To investigate temporal properties of actions.

## Approach:

- the ongoing of time is embodied as the execution of atomic actions (time units)
- two temporal assertions are introduced:

$\phi, \psi ::= C(p) \mid R(p,q) \mid \langle \pi \rangle \phi \mid [\pi] \phi \mid \neg \phi \mid \phi \vee \psi \mid \mathbf{E}(\phi \mathbf{U}^\pi \psi) \mid \mathbf{A}(\phi \mathbf{U}^\pi \psi)$

$\mathbf{E}(\phi \mathbf{U}^\pi \psi)$  : there exists some path of  $\pi$  such that “ $\phi$  until  $\psi$ ” holds

$\mathbf{A}(\phi \mathbf{U}^\pi \psi)$  : “ $\phi$  until  $\psi$ ” holds in any path of  $\pi$

$$\mathbf{EX} \phi =_{\text{def}} \vee_{\alpha \in N_A} \langle \alpha \rangle \phi$$

$$\mathbf{E}(\phi \mathbf{U} \psi) =_{\text{def}} \mathbf{E}(\phi \mathbf{U}^{(\alpha_1 \mathbf{U} \dots \mathbf{U} \alpha_n)^*} \psi)$$

$$\mathbf{A}(\phi \mathbf{U} \psi) =_{\text{def}} \mathbf{A}(\phi \mathbf{U}^{(\alpha_1 \mathbf{U} \dots \mathbf{U} \alpha_n)^*} \psi)$$

$$\mathbf{EF} \varphi =_{\text{def}} \mathbf{E}(\text{true} \mathbf{U} \varphi)$$

$$\mathbf{AF} \varphi =_{\text{def}} \mathbf{A}(\text{true} \mathbf{U} \varphi)$$

$$\mathbf{EG} \varphi =_{\text{def}} \neg \mathbf{AF}(\neg \varphi)$$

$$\mathbf{AG} \varphi =_{\text{def}} \neg \mathbf{EF}(\neg \varphi)$$

$$\mathbf{AX} \varphi =_{\text{def}} \neg \mathbf{EX}(\neg \varphi)$$

# Description example of TDDL( $X@$ )

- **liveness property**: good things will eventually happen.

**EF**(( $\exists$ bought<sup>-</sup>.Customer)(b))

**E**(Instore(b) **U**<sup>VIPbuybook(a,b)</sup>  $\neg$ Instore(b) )

- **safety property**: bad things will never happen.

**AG**  $\neg$ ( $\geq 2$  bought<sup>-</sup>.Customer)(b) )

**AG** ( Instore(b)  $\vee$  ( $\exists$ bought<sup>-</sup>.Customer)(b) )

- Reduced to satisfiability problem of formulas.
- A **Tableau decision algorithm** is provided.

# Limitation of DDL( $X@$ )/TDDL( $X@$ )

- TBox:
  - only concept definitions, no GCIs
  - acyclic
- RBox:
  - on transitive property
- Atomic action:
  - no defined concept name occurring in the effect set ***post***.

Result of: difficulty of ABox updating.



# Difficulty of ABox updating

Example.

- RBox & TBox:

Trans( $R$ ),

$$A \sqsubseteq \exists R.A, \quad A \sqcap B \sqsubseteq \perp, \quad B \sqsubseteq \forall R.B$$

- ABox:

$A(a)$

- Update or new information:

$(\exists R.B)(a)$

# Difficulty of ABox updating

Assumptions	DLs	Approach	References
Acyclic TBox; no def. concept names in U	ALC~AL CQIO	PMA semantics + distance on primitive concept names (NOT defined concept names).	LLMW06, LLMW11
	DL-Lite <sub>F</sub>	PMA semantics	GLPR06, GLPR07
	DL-Lite <sub>R</sub>	Both revision and update	KZ11, KZC13
	DL-Lite <sub>FR</sub>	Based on $cl_{\mathcal{T}}(A)$ (Coincide with PMA if no role names occurring in GCIs and concept assertions.)	CKNZ10

Thank you!

