

# Advanced Algorithms

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Academic Year 2013-2014

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## Lab 9 – Exercises on network optimization and others

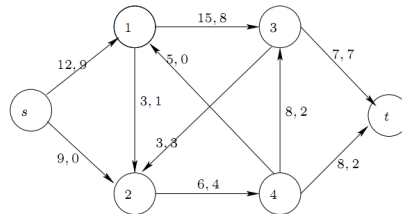
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## Exercise 1

### Network Flow

In the network reported in the following figure the pair of numbers on each edge indicate the capacity and the current flow on that edge, respectively

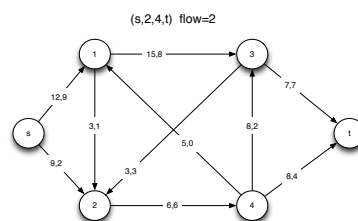
- Find the maximum flow from  $s$  to  $t$  using the Edmond-Karp algorithm



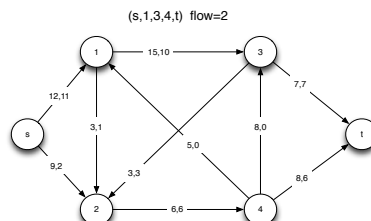
## Exercise 1

### Solution

- Iteration 1



- Iteration 2



Maximum flow =  $11+2 = 7+6 = 13$

## Exercise 2

The knapsack problem is a problem in combinatorial optimization that can be summarized as follows.

- A thief breaks into a house carrying a knapsack
- The capacity of the knapsack is limited to  $m$  Kgs, since the thief can carry at most  $m$  Kgs
- In the house there are  $n$  objects among which the thief can choose
- Each object  $i$  has some weight  $w_i$  and some value  $v_i$

**The thief has to choose which objects to steal in order to maximize his revenue but not exceeding the weight he can carry**

This is called 0-1 knapsack problem because each object can either be stolen (1) or not stolen (0)

- A. Provide an LP formulation of the 0-1 Knapsack problem
- B. How a brute-force approach will work? Which is the complexity of this brute force algorithm?

## Exercise 2

### Solution

- A. This is an optimization problem for solving which we can use  $n$  binary variables  $x_1, \dots, x_n$ , one for each object in the house
  - A linear programming formulation for the 0-1 knapsack is the following:
    - maximize  $\sum_{i=1}^n v_i * x_i$
    - Subject to
    - $\sum_{i=1}^n w_i * x_i \leq m$
    - $x_i \in \{0,1\}, 1 \leq i \leq n$

## Exercise 2

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### Solution

#### B. Brute-force approach

- Since there are  $n$  objects, there are  $2^n$  possible combinations of these objects
- For each combination we have to compute the value and the weight with cost  $O(n)$ , in order to find the one with the maximum value and total weight less than or equals to  $m$
- Thus, the total cost is  $O(n*2^n)$

## Assignment 08

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### Exercise 7.18 page 226 DPV

There are many common variations of the maximum flow problem. Here are three of them.

- A. There are many sources and many sinks, and we wish to maximize the total flow from all sources to all sinks
- B. Each vertex also has a capacity on the maximum flow that can enter it
- C. Each edge has not only a capacity, but also a lower bound on the flow it can carry

Each of these variations can be solved efficiently. Show this by reducing (A) and (B) to the original max-flow problem, and reducing (C) and to linear programming

## Assignment 08

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### Exercise

- Consider the 0-1 knapsack problem of Exercise 2
- A. Propose some greedy strategy to solve the problem and find a counterexample for each of them
- B. Consider the greedy strategy of selecting at each step the object with the maximum value per unit of weight and provide a counterexample showing that this strategy does not return the optimal solution