

Single agent or multiple agents

- Many domains are characterized by multiple agents rather than a single agent.
- **Game theory** studies what agents should do in a multi-agent setting.
- Agents can be cooperative, competitive or somewhere in between.
- Agents that are strategic can't be modeled as nature.

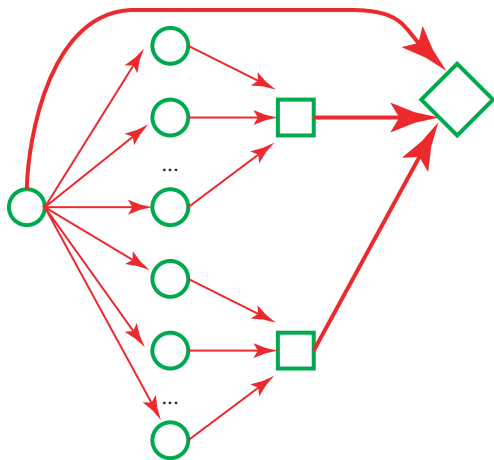
Multi-agent framework

- Each agent can have its own values.
- Agents select actions autonomously.
- Agents can have different information.
- The outcome can depend on the actions of all of the agents.
- Each agent's value depends on the outcome.

Fully Observable + Multiple Agents

- If agents act sequentially and can observe the state before acting: **Perfect Information Games.**
- Can do dynamic programming or search:
Each agent maximizes for itself.
- Two person, competitive (zero sum) \implies minimax.

Multiple Agents, shared value



Complexity of Multi-agent decision theory

- It can be exponentially harder to find optimal multi-agent policy even with a shared values.
- **Why?** Because dynamic programming doesn't work:
 - ▶ If a decision node has n binary parents, dynamic programming lets us solve 2^n decision problems.
 - ▶ This is much better than d^{2^n} policies (where d is the number of decision alternatives).
- Multiple agents with shared values is equivalent to having a single forgetful agent.

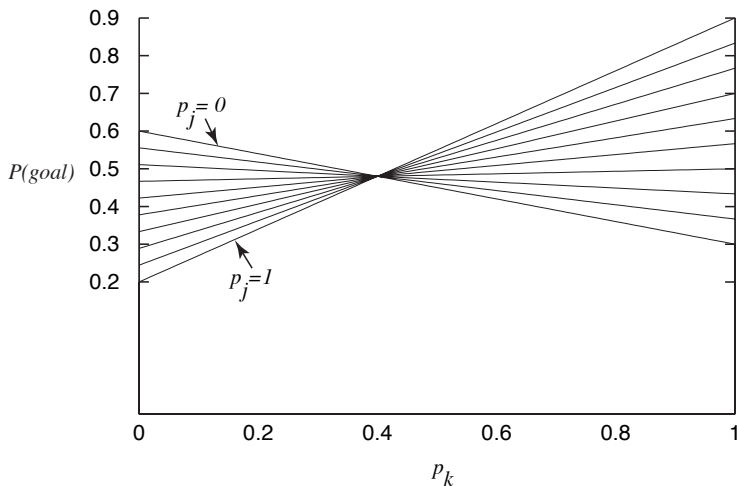
Partial Observability and Competition



		goalie	
		left	right
kicker	left	0.6	0.2
	right	0.3	0.9

Probability of a goal.

Stochastic Policies



Strategy Profiles

- Assume a general n -player game,
- A **strategy** for an agent is a probability distribution over the actions for this agent.
- A **strategy profile** is an assignment of a strategy to each agent.
- If σ is a strategy profile:
 - σ_i is the strategy of agent i in σ ,
 - σ_{-i} is the set of strategies of the other agents.
 - Thus σ is $\sigma_i\sigma_{-i}$
- A strategy profile σ has a utility for each agent. Let $utility(\sigma, i)$ be the utility of strategy profile σ for agent i .

Nash Equilibria

- σ_i is a **best response** to σ_{-i} if for all other strategies σ'_i for agent i ,

$$utility(\sigma_i \sigma_{-i}, i) \geq utility(\sigma'_i \sigma_{-i}, i).$$

- A strategy profile σ is a **Nash equilibrium** if for each agent i , strategy σ_i is a best response to σ_{-i} . That is, a Nash equilibrium is a strategy profile such that no agent can be better by unilaterally deviating from that profile.
- Theorem [Nash, 1950] Every finite game has at least one Nash equilibrium.

Multiple Equilibria

Hawk-Dove Game:
Agent 2

		dove	hawk
Agent 1	dove	$R/2, R/2$	$0, R$
	hawk	$R, 0$	$-D, -D$

D and R are both positive with $D \gg R$.

Coordination

Just because you know the Nash equilibria doesn't mean you know what to do:

		Agent 2	
		shopping	football
Agent 1	shopping	2,1	0,0
	football	0,0	1,2

Prisoner's Dilemma

Two strangers are in a game show. They each have the choice:

- Take \$100 for yourself
- Give \$1000 to the other player

This can be depicted as the payoff matrix:

		Player 2	
		take	give
Player 1	take	100,100	1100,0
	give	0,1100	1000,1000

Computing Nash Equilibria

Suppose you are given a game that specifies the expected value for each agent for each strategy profile. (This is the **strategic form of a game**).

To compute a Nash equilibria for a game in strategic form, there are three steps:

- Eliminate dominated strategies
- Determine which actions will have non-zero probabilities; this is called the **support set**
- Determine the probability for the actions in the support set

Eliminating Dominated Strategies

		Agent 2		
		d_2	e_2	f_2
Agent 1	a_1	3,5	5,1	1,2
	b_1	1,1	2,9	6,4
	c_1	2,6	4,7	0,8

Computing probabilities in randomized strategies

Given a support set:

- The only reason that an agent will randomize between actions $a_1 \dots a_k$ is if actions $a_1 \dots a_k$ have the same value for that agent given the randomized strategies actions for the other agents
- This forms a set of simultaneous equations that can be solved where the free variable are the probabilities of the actions
- If there is a solution with all the probabilities in range $(0,1)$ this is a Nash equilibrium.
- You may have to search over support sets to find a Nash equilibrium