Description Logics

Structural Description Logics

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Description Logics

- A logical reconstruction and unifying formalism for the representation tools
 - Frame-based systems
 - Semantic Networks
 - Object-Oriented representations
 - Semantic data models
 - Ontology languages
 - •
- A structured fragment of predicate logic
- Provide theories and systems for expressing structured information and for accessing and reasoning with it in a principled way.

Applications

Description logics based systems are currently in use in many applications.

- Configuration
- Conceptual Modeling
- Query Optimization and View Maintenance
- Natural Language Semantics
- I3 (Intelligent Integration of Information)
- Information Access and Intelligent Interfaces
- Terminologies and Ontologies
- Software Management
- Planning

A formalism

- Description Logics formalize many Object-Oriented representation approaches.
- As such, their purpose is to disambiguate many imprecise representations.

Frames or Objects

- Identifier
- Class
- Instance
- Slot (attribute)
 - Value
 - Identifier
 - Default
 - Value restriction
 - Type
 - Concrete Domain
 - Cardinality
 - Encapsulated method

Ambiguities: classes and instances

```
Person: AGE: Number,
```

SEX: M, F,

HEIGHT : Number,

WIFE: Person.

john : AGE : 29,

SEX: M,

HEIGHT: 76,

WIFE: mary.

Ambiguities: incomplete information

```
\begin{aligned} \textit{john}: \text{AGE}: 29, \\ \text{SEX}: M, \\ \text{HEIGHT}: \text{Number}, \\ \text{WIFE}: \text{Person}. \end{aligned}
```

Ambiguities: is-a

Sub-class:

```
Person: AGE: Number,
```

SEX: M, F,

HEIGHT: Number,

WIFE: Person.

 \uparrow

Male: AGE: Number,

SEX : M,

HEIGHT : Number,

WIFE: Female.

Ambiguities: is-a

Instance-of:

```
Male: AGE: Number,
```

SEX : M,

HEIGHT : Number,

WIFE: Female.



john : AGE : 35,

SEX: M,

HEIGHT: 76,

WIFE: mary.

Ambiguities: is-a

Instance-of:

```
29'er: AGE: 29, {\tt SEX}: M, \\ {\tt HEIGHT}: {\tt Number},
```

 \uparrow

WIFE: Person.

```
\begin{aligned} \textit{john}: \text{AGE}: 29, \\ \text{SEX}: M, \\ \text{HEIGHT}: \text{Number}, \\ \text{WIFE}: \text{Person}. \end{aligned}
```

Implicit relation:

```
john : AGE : 35,
```

SEX: M,

HEIGHT: 76,

 $\underline{\mathtt{WIFE}}: mary.$

mary : AGE : 32,

SEX : F,

HEIGHT: 59,

 $\underline{\mathtt{HUSBAND}}: john.$

Explicit relation:

john : AGE : 35,

SEX : M,

HEIGHT: 76.

mary : AGE : 32,

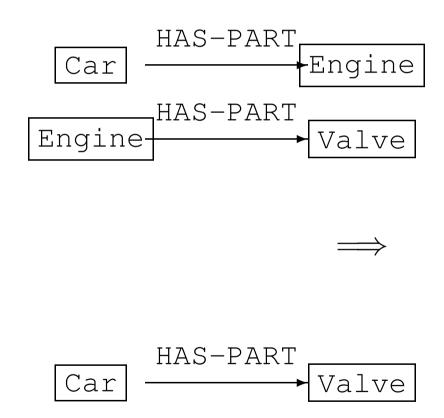
SEX : F,

HEIGHT : 59.

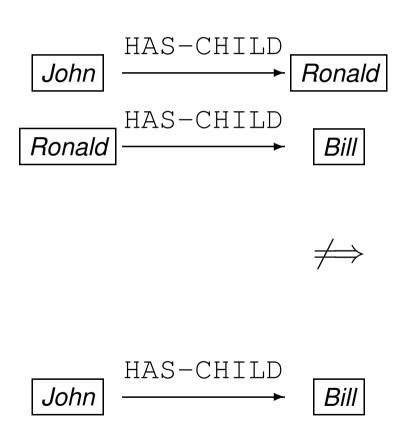
 $\textit{m-j-family}: \underline{\mathtt{WIFE}}: mary,$

 $\underline{\mathtt{HUSBAND}}: john.$

Special relation:

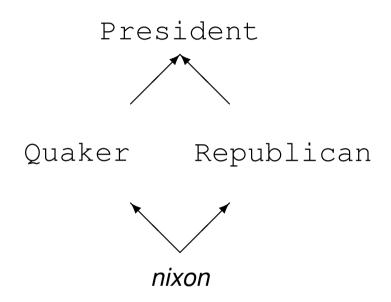


Normal relation:



Ambiguities: default

The *Nixon* diamond:



Quakers are pacifist, Republicans are not pacifist.

⇒ Is Nixon pacifist or not pacifist?

What is the exact meaning of:



What is the exact meaning of:



Every frog is just green

What is the exact meaning of:



- Every frog is just green
- Every frog is also green

What is the exact meaning of:



- Every frog is just green
- Every frog is also green
- Every frog is of some green

What is the exact meaning of:



- Every frog is just green
- Every frog is also green
- Every frog is of some green
- There is a frog, which is just green

•

What is the exact meaning of:



- Every frog is just green
- Every frog is also green
- Every frog is of some green
- There is a frog, which is just green
- •
- Frogs are typically green, but there may be exceptions

False friends

- The meaning of object-oriented representations is logically very ambiguous.
- The appeal of the graphical nature of object-oriented representation tools has led to forms of reasoning that do not fall into standard logical categories, and are not yet very well understood.
- It is unfortunately much easier to develop some algorithm that appears to reason over structures of a certain kind, than to *justify* its reasoning by explaining what the structures are saying about the domain.

A structured logic

- Any (basic) Description Logic is a fragment of FOL.
- The representation is at the predicate level: no variables are present in the formalism.
- A Description Logic theory is divided in two parts:
 - the definition of predicates (*TBox*)
 - the assertion over constants (*ABox*)
- Any (basic) Description Logic is a subset of \mathcal{L}_3 , i.e. the function-free FOL using only at most *three* variable names.

Why not FOL

If FOL is directly used without additional restrictions then

- the structure of the knowledge is destroyed, and it can not be exploited for driving the inference;
- the expressive power is too high for obtaining decidable and efficient inference problems;
- the inference power may be too low for expressing interesting, but still decidable theories.

Structured Inheritance Networks: KL-ONE

- Structured Descriptions
 - corresponding to the complex relational structure of objects,
 - built using a restricted set of epistemologically adequate constructs
- distinction between conceptual (terminological) and instance (assertional) knowledge;
- central role of automatic classification for determining the subsumption i.e., universal implication – lattice;
- strict reasoning, no defaults.

Types of the TBox Language

Concepts – denote entities

 (unary predicates, classes)

Example: Student, Married $\{x \mid \mathtt{Student}(x)\},\$ $\{x \mid \mathtt{Married}(x)\}$

Roles
 – denote properties
 (binary predicates, relations)

Example: FRIEND, LOVES

$$\{\langle x, y \rangle \mid \text{FRIEND}(x, y)\},\$$

 $\{\langle x, y \rangle \mid \text{LOVES}(x, y)\}$

Concept Expressions

Description Logics organize the information in classes – *concepts* – gathering homogeneous data, according to the relevant common properties among a collection of instances.

Example:

Student □ ∃FRIEND.Married

$$\{x \mid \mathtt{Student}(x) \land \\ \exists y \mathtt{.FRIEND}(x,y) \land \mathtt{Married}(y) \}$$

A note on λ 's

In general, λ is an explicit way of forming *names* of functions:

 λx . f(x) is the function that, given input x, returns the value f(x)

The λ -conversion rule says that:

$$(\lambda x. f(x))(a) = f(a)$$

Thus, λx . $(x^2 + 3x - 1)$ is the function that applied to 2 gives 9:

$$(\lambda x. (x^2 + 3x - 1))(2) = 9$$

We can give a name to this function, so that:

$$f_{231} \doteq \lambda x$$
. $(x^2 + 3x - 1)$
 $f_{231}(2) = 9$

λ to define predicates

Predicates are special case of functions: they are *truth* functions. So, if we think of a formula P(x) as denoting a truth value which may vary as the value of x varies, we have:

 λx . P(x) denotes a function from domain individuals to truth values.

In this way, as we have learned from FOL, P denotes exactly the set of individuals for which it is true. So, P(a) means that the individual a makes the predicate P true, or, in other words, that a is in the extension of P.

For example, we can write for the *unary* predicate Person:

Person
$$\doteq \lambda x$$
. Person (x)

which is equivalent to say that Person denotes the set of persons:

$$\begin{aligned} & \text{Person} \leadsto \{x \mid \text{Person}(x)\} \\ & \text{Person}^{\mathcal{I}} = \{x \mid \text{Person}(x)\} \\ & \text{Person}(john) \text{ IFF } john^{\mathcal{I}} \in \text{Person}^{\mathcal{I}} \end{aligned}$$

In the same way for the *binary* predicate FRIEND:

FRIEND
$$\doteq \lambda x, y$$
. FRIEND (x, y)
FRIEND $^{\mathcal{I}} = \{ \langle x, y \rangle \mid \text{FRIEND}(x, y) \}$

The functions we are defining with the λ operator may be parametric:

Student
$$\sqcap$$
 Worker = λx . (Student $(x) \land \text{Worker}(x)$)

(Student \sqcap Worker) $^{\mathcal{I}} = \{x \mid (\text{Student}(x) \land \text{Worker}(x))\}$

(Student \sqcap Worker) $^{\mathcal{I}} = \text{Student}^{\mathcal{I}} \cap \text{Worker}^{\mathcal{I}}$

(Verify as exercise)

Concept Expressions

```
(\mathtt{Student} \sqcap \exists \mathtt{FRIEND}.\mathtt{Married})^{\mathcal{I}}
(\mathtt{Student})^{\mathcal{I}} \cap (\exists \mathtt{FRIEND}.\mathtt{Married})^{\mathcal{I}}
\{x \mid \mathtt{Student}(x)\} \cap
\{x \mid \exists y \text{.FRIEND}(x, y) \land \text{Married}(y)\}
\{x \mid \mathtt{Student}(x) \land
         \exists y. \mathtt{FRIEND}(x, y) \land \mathtt{Married}(y) \}
```

Objects: classes

Student

Person

name: [String]

address: [String]

enrolled: [Course]

```
\{x \mid \mathtt{Student}(x)\} = \{x \mid \mathtt{Person}(x) \land \\ (\exists y. \, \mathtt{NAME}(x,y) \land \mathtt{String}(y)) \land \\ (\exists z. \, \mathtt{ADDRESS}(x,z) \land \mathtt{String}(z)) \land \\ (\exists w. \, \mathtt{ENROLLED}(x,w) \land \mathtt{Course}(w)) \, \} \mathtt{Student} \doteq \mathtt{Person} \, \sqcap \\ \exists \mathtt{NAME.String} \, \sqcap \\ \exists \mathtt{ADDRESS.String} \, \sqcap \\ \exists \mathtt{ENROLLED.Course}
```

Objects: instances

s1: Student

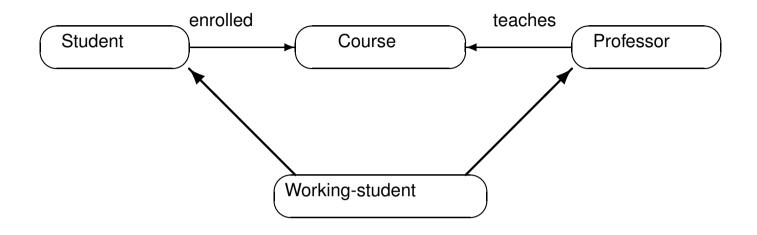
name: "John"

address: "Abbey Road..."

enrolled: cs415

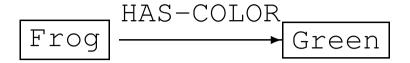
$$\label{eq:student} \begin{split} &\text{Student}(\texttt{s1}) \land \\ &\text{NAME}(\texttt{s1}, \texttt{"john"}) \land \texttt{String}(\texttt{"john"}) \land \\ &\text{ADDRESS}(\texttt{s1}, \texttt{"abbey-road"}) \land \texttt{String}(\texttt{"abbey-road"}) \land \\ &\text{ENROLLED}(\texttt{s1}, \texttt{cs415}) \land \texttt{Course}(\texttt{cs415}) \end{split}$$

Semantic Networks



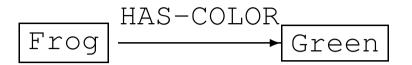
- $\forall x. \ \, \text{Working-student}(x) \rightarrow \\ \text{Student}(x) \land \text{Professor}(x)$

Quantification



- Frog
 □ ∃HAS—COLOR.Green:
 Every frog is also green
- Frog
 □ ∀HAS—COLOR.Green:
 Every frog is just green
- Frog \sqsubseteq \forall HAS-COLOR.Green Frog(x), HAS-COLOR(x,y): There is a frog, which is just green

Quantification: existential



Every frog is also green

Frog
$$\sqsubseteq \exists HAS-COLOR.Green$$

$$\forall x. \ \mathrm{Frog}(x) \rightarrow \\ \exists y. \ (\mathrm{HAS-COLOR}(x,y) \land \mathrm{Green}(y))$$

Exercise: is this a model?

Frog(oscar), Green(green),

HAS-COLOR(oscar, green),

Red(red),

HAS-COLOR(oscar,red).

Quantification: universal

Every frog is only green

Frog
$$\sqsubseteq \forall \texttt{HAS-COLOR}.\texttt{Green}$$

$$\forall x. \operatorname{Frog}(x) \rightarrow \\ \forall y. (\operatorname{HAS-COLOR}(x,y) \rightarrow \operatorname{Green}(y))$$

Exercise: is this a model? and this?

Frog(oscar), Green(green), Frog(sing),

HAS-COLOR(oscar, green), AGENT(sing, oscar).

Red(red),

HAS-COLOR(oscar,red).

Analytic reasoning (intuition)

Person
subsumes
(Person with every male friend is a doctor)
subsumes
(Person with every friend is a
(Doctor with a specialty is surgery))

Analytic reasoning (intuition)

```
Person
subsumes
(Person with every male friend is a doctor)
subsumes
(Person with every friend is a
   (Doctor with a specialty is surgery))
(Person with \geq 2 children)
subsumes
(Person with \geq 3 male children)
```

Analytic reasoning (intuition)

```
Person
subsumes
(Person with every male friend is a doctor)
subsumes
(Person with every friend is a
   (Doctor with a specialty is surgery))
(Person with \geq 2 children)
subsumes
(Person with \geq 3 male children)
(Person with \geq 3 young children)
disjoint
(Person with \leq 2 children)
```