

When are Description Logic Knowledge Bases Indistinguishable?

E. Botoeva,¹ R. Kontchakov,² V. Ryzhikov,¹ F. Wolter³ and M. Zakharyashev²

¹Free University of Bozen-Bolzano, Italy ²Birkbeck, University of London, UK ³University of Liverpool, UK

Versioning

\mathcal{K}_2 is a Σ -query version of \mathcal{K}_1 if $\mathcal{K}_2 \equiv_{\Sigma} \mathcal{K}_1$

Query Entailment and Inseparability

Let \mathcal{K}_1 and \mathcal{K}_2 be knowledge bases: $\mathcal{K}_i = (\mathcal{T}_i, \mathcal{A}_i)$

\mathcal{K}_1 Σ -query entails \mathcal{K}_2 if for each **CQ** $q(\vec{x})$ over Σ and each tuple \vec{a} in $\text{ind}(\mathcal{K}_2)$ $\mathcal{K}_2 \models q(\vec{a})$ implies $\mathcal{K}_1 \models q(\vec{a})$

\mathcal{K}_1 and \mathcal{K}_2 are Σ -query inseparable, $\mathcal{K}_1 \equiv_{\Sigma} \mathcal{K}_2$, if

\mathcal{K}_1 Σ -query entails \mathcal{K}_2 and \mathcal{K}_2 Σ -query entails \mathcal{K}_1

Knowledge Exchange

\mathcal{K}_2 is a **universal UCQ-solution** for \mathcal{K}_1 under \mathcal{T}_{12} if $\mathcal{K}_2 \equiv_{\Sigma_2} \mathcal{K}_1 \cup \mathcal{T}_{12}$

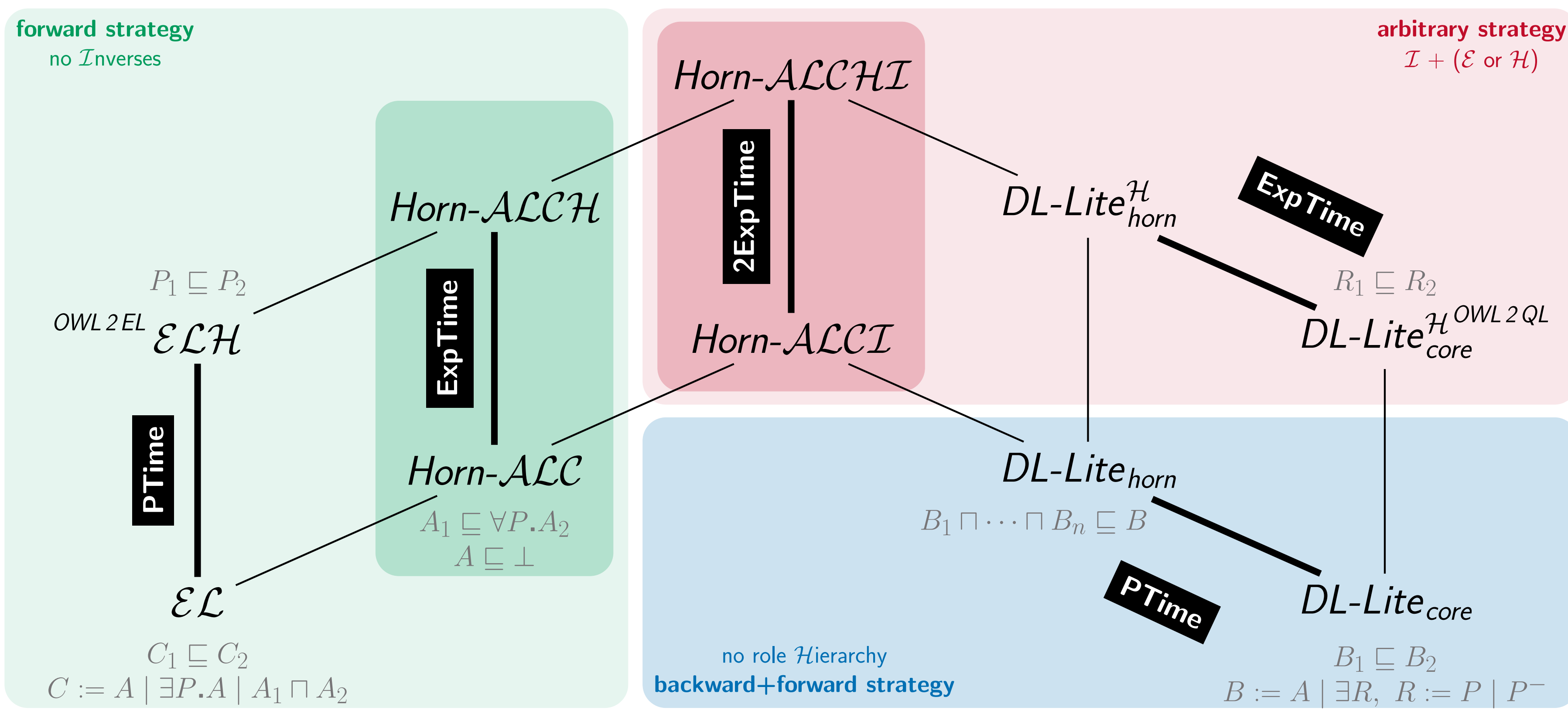
Forgetting

\mathcal{K}' is the result of **forgetting** Σ in \mathcal{K} if $\text{sig}(\mathcal{K}') \subseteq \text{sig}(\mathcal{K}) \setminus \Sigma$ and $\mathcal{K}' \equiv_{\text{sig}(\mathcal{K}) \setminus \Sigma} \mathcal{K}$

Modularisation

$\mathcal{K}' \subseteq \mathcal{K}$ is a Σ -module of \mathcal{K} if $\mathcal{K}' \equiv_{\Sigma} \mathcal{K}$

The Landscape of Logics



Horn-ALCHI normal form:

- $C \sqsubseteq A,$
- $C \sqsubseteq \forall R.A,$
- $A \sqsubseteq \exists R.C,$
- $\exists R.C \sqsubseteq A,$
- $A_1 \sqcap A_2 \sqsubseteq A,$
- $R_1 \sqsubseteq R_2,$

where

- A, A_1, A_2 are concept names
- C is a concept name or \top
- R, R_1, R_2 are role names or their inverses

Query Entailment

Let $\mathcal{K}_1 = (\mathcal{T}_1, \{A(a)\})$ and $\mathcal{K}_2 = (\mathcal{T}_2, \{A(a)\})$

$$\mathcal{T}_1 = \left\{ \begin{array}{l} A \sqsubseteq \exists S.(\exists R.A \sqcap \exists T.\exists W.\exists Q) \\ \exists Q^- \sqsubseteq \exists Q, S \sqsubseteq S_1, T \sqsubseteq T_1, W \sqsubseteq W_1 \end{array} \right\}$$

$$\mathcal{T}_2 = \left\{ \begin{array}{l} A \sqsubseteq \exists P.B, \exists Q^- \sqsubseteq \exists Q \\ B \sqsubseteq \exists R^-.(\exists S^-.B \sqcap \exists T) \\ \exists T^- \sqsubseteq \exists W.(\exists W_1^-.\exists T_1^-.\exists S_1^- \sqcap \exists Q) \end{array} \right\}$$

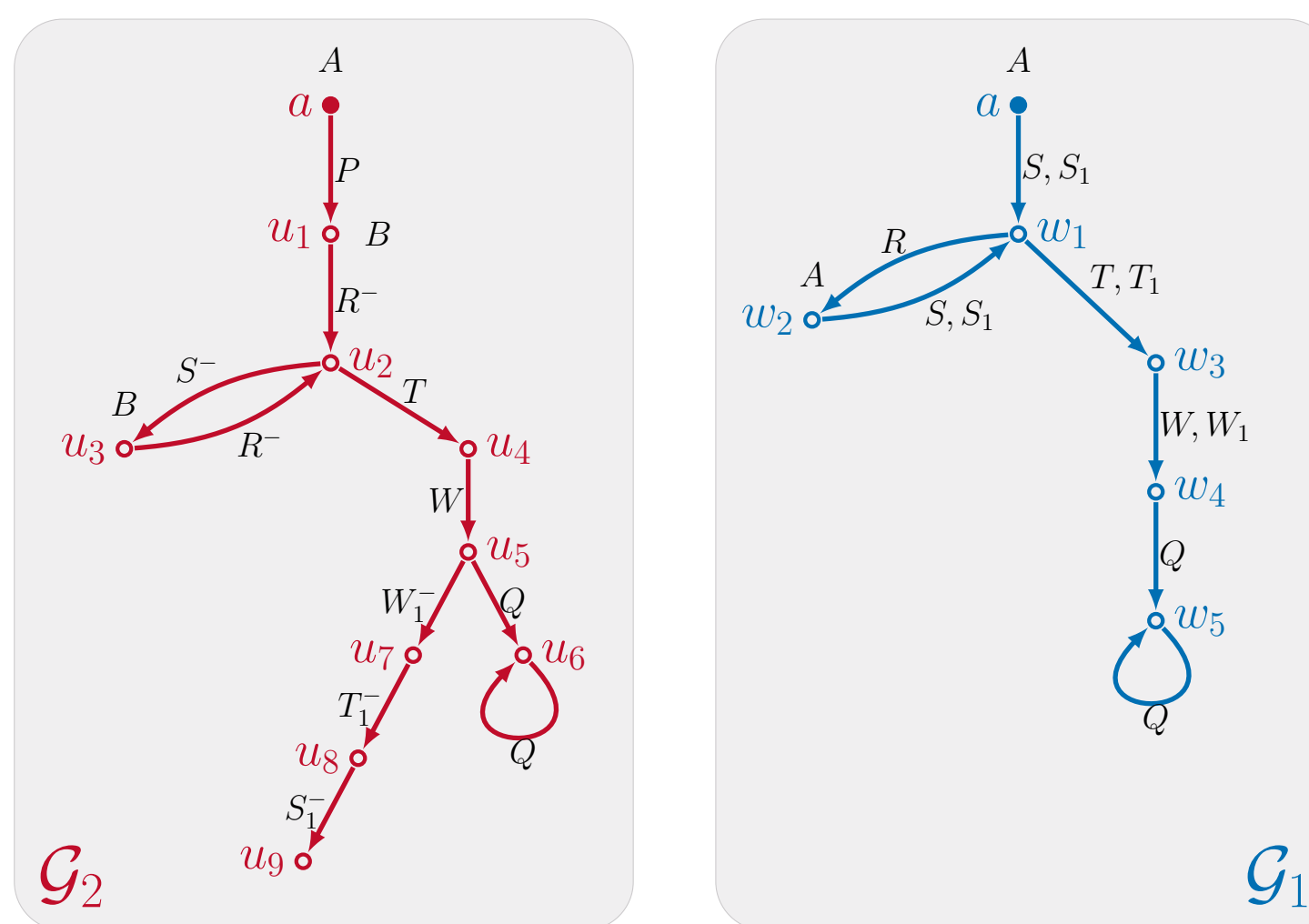
$$\Sigma = \{R, S, S_1, T, T_1, Q, W, W_1\}$$

\mathcal{K}_1 Σ -query entails \mathcal{K}_2

Games and Generating Structures

\mathcal{G}_i is a **generating structure** for \mathcal{M}_i if \mathcal{M}_i is the unravelling of \mathcal{G}_i

(used in the combined approach to query answering)



Forward strategies (no inverse roles): Player 1 can only move forward in \mathcal{M}_1

General strategies are combinations of a backward strategy

Player 1 can only move backwards in \mathcal{M}_1

and a number of **start-bounded** strategies

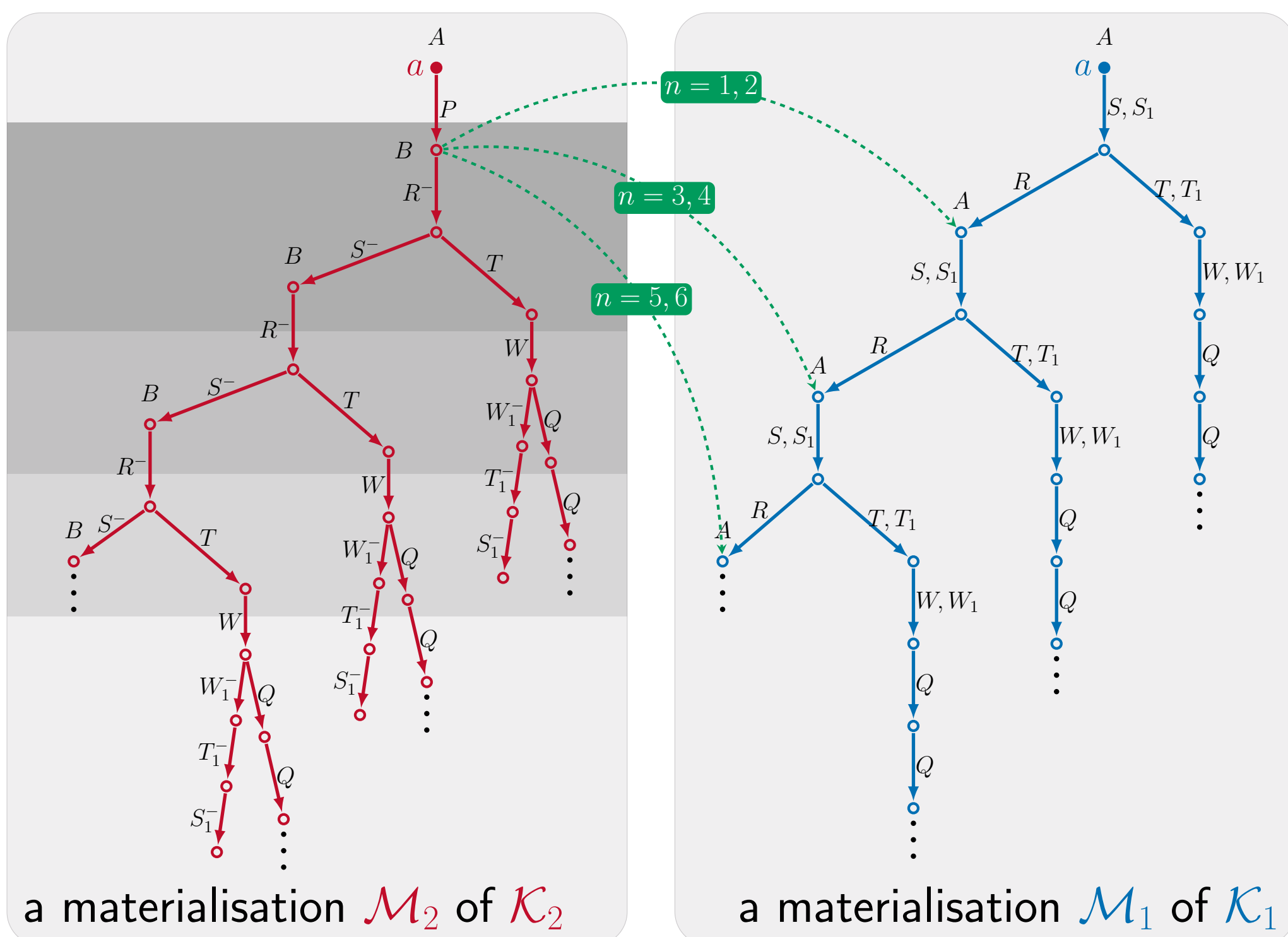
Player 1 cannot move in \mathcal{M}_1 closer to the ABox than the initial response

Player 2 has no winning strategy against Player 1 in an elaborate game on \mathcal{G}_2 and \mathcal{G}_1

Player 2 has no winning strategy against Player 1 in the naive game on \mathcal{M}_2 and \mathcal{M}_1

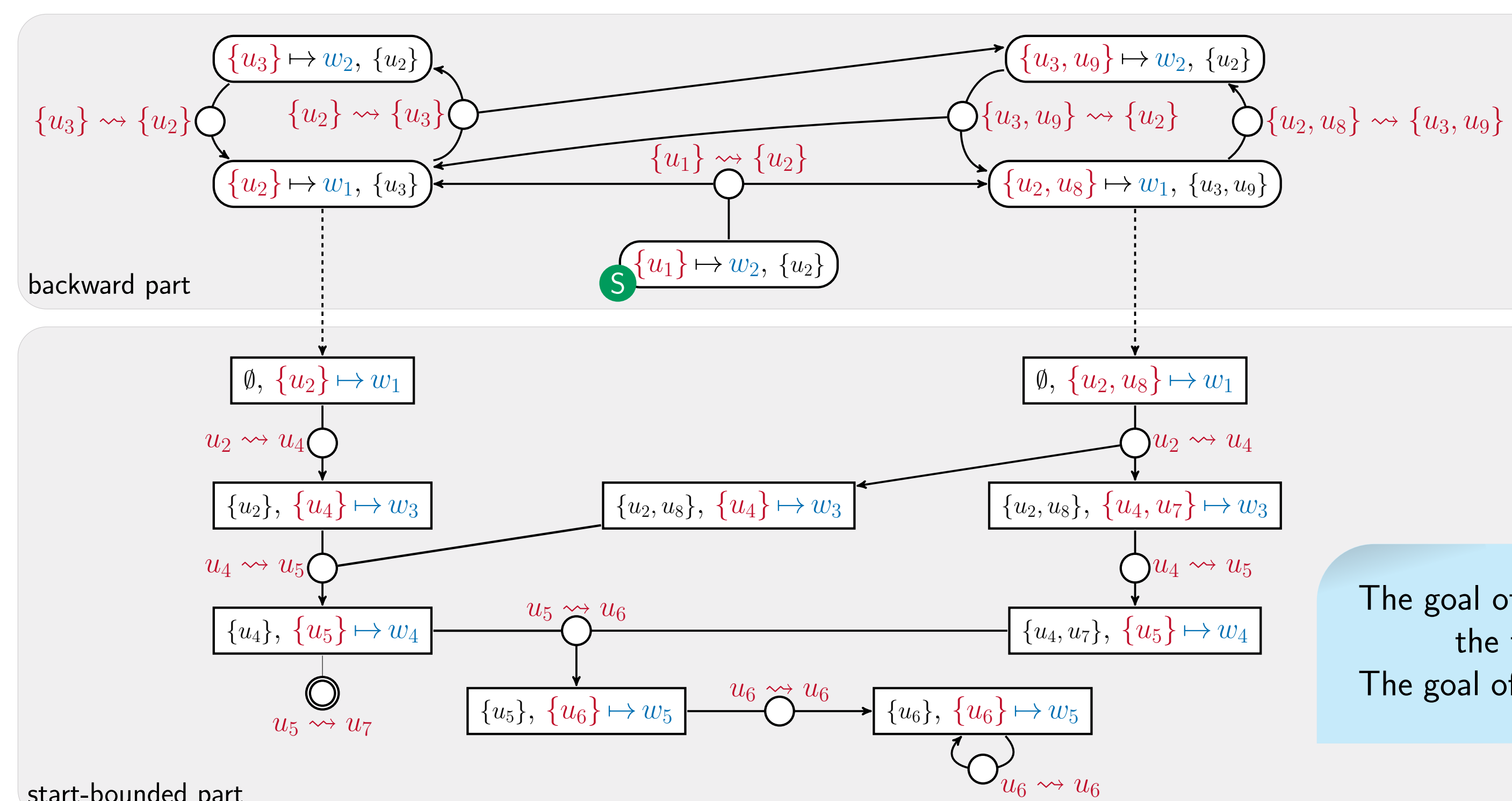
Σ -Homomorphisms

\mathcal{M}_2 is finitely Σ -homomorphically embeddable into \mathcal{M}_1 (each finite subinterpretation of \mathcal{M}_2 can be Σ -homomorphically mapped to \mathcal{M}_1)



\mathcal{M} is a **materialisation** of \mathcal{K} if for each $q(\vec{x})$ and each \vec{a} $\mathcal{K} \models q(\vec{a})$ iff $\mathcal{M} \models q(\vec{a})$

Let's play?



The goal of **Player 2** is to reach the final state (double circle)
The goal of **Player 1** is to avoid it