

Review of basic definitions:

- Alphabet: finite, nonempty set of symbols Σ

e.g. $\Sigma = \{0, 1\}$

$\Sigma = \{e, k, \dots, z\}$

$\Sigma =$ set of Unicode characters

- String: finite sequence of symbols from Σ

$w = a_1 a_2 \dots a_n$, with $a_i \in \Sigma$ for $i \in \{1, \dots, n\}$

e.g. $\cdot 01101$

\cdot ~~ciociocio~~

• empty string: denoted ϵ : string with no symbols

• length of a string = number of (positions for) symbols in the string

denoted $|w|$ \exists if $w = a_1 \dots a_n$, then $|w| = n$

e.g. $|\epsilon| = 0$

ϵ is the only string of length 0

$|0| = 1$

~~$|ciociocio| = 8$~~

Notice: strictly speaking, the number of symbols in ~~ciociocio~~ is 4

- Powers of an alphabet:

$\Sigma^k = \underbrace{\Sigma \times \Sigma \times \dots \times \Sigma}_{k \text{ times}}$... set of all strings over Σ of length k

e.g. $\Sigma^0 = \{\epsilon\}$

$\{0, 1\}^1 = \{0, 1\}$

← what is the difference between lhs and rhs?

$\{0, 1\}^2 = \{00, 01, 10, 11\}$

Closure of an alphabet Σ : Σ^* is the set of all finite strings over Σ

i.e. $\Sigma^* = \Sigma^0 \cup \Sigma^1 \cup \Sigma^2 \cup \dots$

also $\Sigma^+ = \Sigma^1 \cup \Sigma^2 \cup \dots$ hence $\Sigma^* = \Sigma^0 \cup \Sigma^+$

Note: all strings in Σ^* are finite
 Σ^* is an infinite set

e.g. $\Sigma = \{0, 1\}$

$\Sigma^* = \{0, 1\}^* = \{\epsilon, 0, 1, 00, 01, 10, 11, 000, 001, \dots\}$

Concatenation of two strings:

$x = a_1 a_2 \dots a_m \in \Sigma^*$

$y = b_1 b_2 \dots b_n \in \Sigma^*$

$\Rightarrow xy = a_1 \dots a_m b_1 \dots b_n$ (we may omit the \cdot)

Note: $\epsilon \cdot x = x \cdot \epsilon = x$, i.e. ϵ is the identity for conc.

$|xy| = |x| + |y|$

Language L over Σ : is any subset of Σ^* (i.e. $L \subseteq \Sigma^*$)

Note: L contains only finite strings, but it may be infinite

Examples:

$\left\{ \begin{array}{l} \Sigma = \{e, b, \dots, z\} \\ L = \text{set of all English words} \end{array} \right.$

$\left\{ \begin{array}{l} \Sigma = \text{Unicode characters} \\ L = \text{compilable Java programs} \end{array} \right.$

$\left\{ \begin{array}{l} \Sigma = \{0, 1\} \\ L = \{\epsilon, 01, 0011, 000111, \dots\} \end{array} \right.$ all strings with equal # of 0 and 1, with all 0's preceding the 1's

\emptyset the empty language ($\neq \{\epsilon\}$)