

Proving languages not to be regular

4.5

Consider: $L_{\text{alt}} = \{w \mid \text{has alternating 0's and 1's}\}$

$L_{\text{eq}} = \{w \mid \text{has an equal number of 0's and 1's}\}$

• Claim: L_{alt} is regular

Proof: easy $E_{\text{alt}} = (\epsilon + 0)(1 + 0)^*(\epsilon + 1)$ is such that $\mathcal{L}(E_{\text{alt}}) = L_{\text{alt}}$

• Claim: L_{eq} is not regular

How can we prove this?

Intuition: • DFA with n states can count up to n .

• to decide whether $w \in L_{\text{eq}}$ we need unbounded counting (since w may be arbitrarily long)

Pumping Lemma:

For all regular languages $L \subseteq \Sigma^*$

there exists n (which depends on L) such that

for all $w \in L$ with $|w| \geq n$

there exists a decomposition $w = xyz$ of w s.t.

1) $|y| \geq 1$ (i.e., $y \neq \epsilon$)

2) $|x \cdot y| \leq n$

3) for all $k \geq 0$, $xy^kz \in L$.

Intuitively, for every $w \in L$, we can find a substring y "near" the beginning of w that can be "pumped", while still obtaining words in L .

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↓
Proof:

4.6

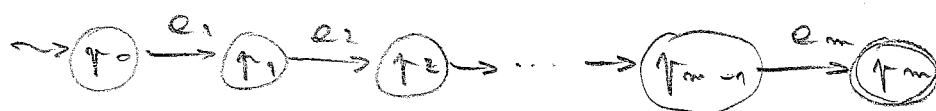
Given regular language L , let $A = (Q, \Sigma, \delta, q_0, F)$ be a DFA with $\mathcal{L}(A) = L$.

We take $n = |Q|$.

Consider any $w = e_1 e_2 \dots e_m \in L$ with $m = |w| \geq n$.

Since $w \in \mathcal{L}(A)$, we have that $\hat{\delta}(q_0, w) \in F$.

Define $q_i = \hat{\delta}(q_0, e_1 e_2 \dots e_i) \quad \forall i \in \{1, \dots, m\}$ and $q_0 = q_0$.



Since $m \geq n$,

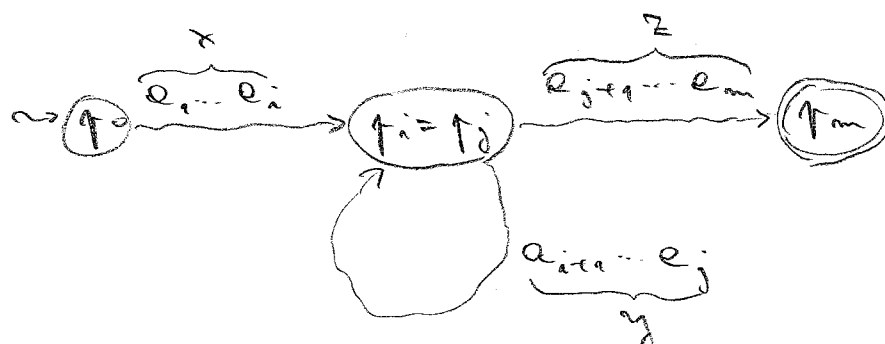
- each $q_i, 0 \leq i \leq m$ belongs to Q , and
- $|Q| = n$

by the pigeon-hole principle, $q_0, q_1, \dots, q_m$ are not all distinct

Let i, j with $0 \leq i < j \leq m$ be the best indices such that

$$q_i = q_j.$$

Hence, to accept w , the DFA goes through a cycle:



Observe: $|y| = j - i \geq 1$ (since $i < j$)

$$|xy| = j \leq m$$

$$\hat{\delta}(q_0, x^k z) = \hat{\delta}(\hat{\delta}(q_0, x), y^k z) = \hat{\delta}(q_i, y^k z) = \hat{\delta}(\hat{\delta}(q_i, y), y^{k-1} z)$$

$$= \hat{\delta}(q_i, y^{k-1} z) = \dots = \hat{\delta}(q_i, z) = q_m \in F \Rightarrow xy^k z \in L$$

q.e.d.

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END