

Formal Languages and Compilers

Lab V: Semantics

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Formal Languages and Compilers — BSc course

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Board

PRODUCTION	SEMANTIC RULES
$Prog \rightarrow S$	$S.next := newlabel; Prog.code := S.code \parallel gen(S.next ':')$
$S \rightarrow S_1 ; S_2$	$S_1.next := newlabel; S_2.next := S.next;$ $S.code := S_1.code \parallel gen(S_1.next ':') \parallel S_2.code$
$S \rightarrow \text{while } Test \text{ do } \{S_1\}$	$Test.begin := newlabel; Test.true := newlabel;$ $Test.false := S.next; S_1.next := Test.begin;$ $S.code := gen(Test.begin ':') \parallel Test.code \parallel gen(Test.true ':') \parallel S_1.code \parallel gen('goto' Test.begin)$
$S \rightarrow id := E$	$S.code := E.code \parallel gen(id.place ' :=' E.place)$
$Test \rightarrow id_1 \leq id_2$	$Test.code := gen('if' id_1.place ' \leq ' id_2.place ' goto' Test.true) \parallel gen('goto' Test.false)$
$E \rightarrow E_1 + id$	$E.place := newtemp;$ $E.code := E_1.code \parallel gen(E.place ' :=' E_1.place ' +' id.place)$
$E \rightarrow id$	$E.place := id.place; E.code := '$

Board

Given the input:

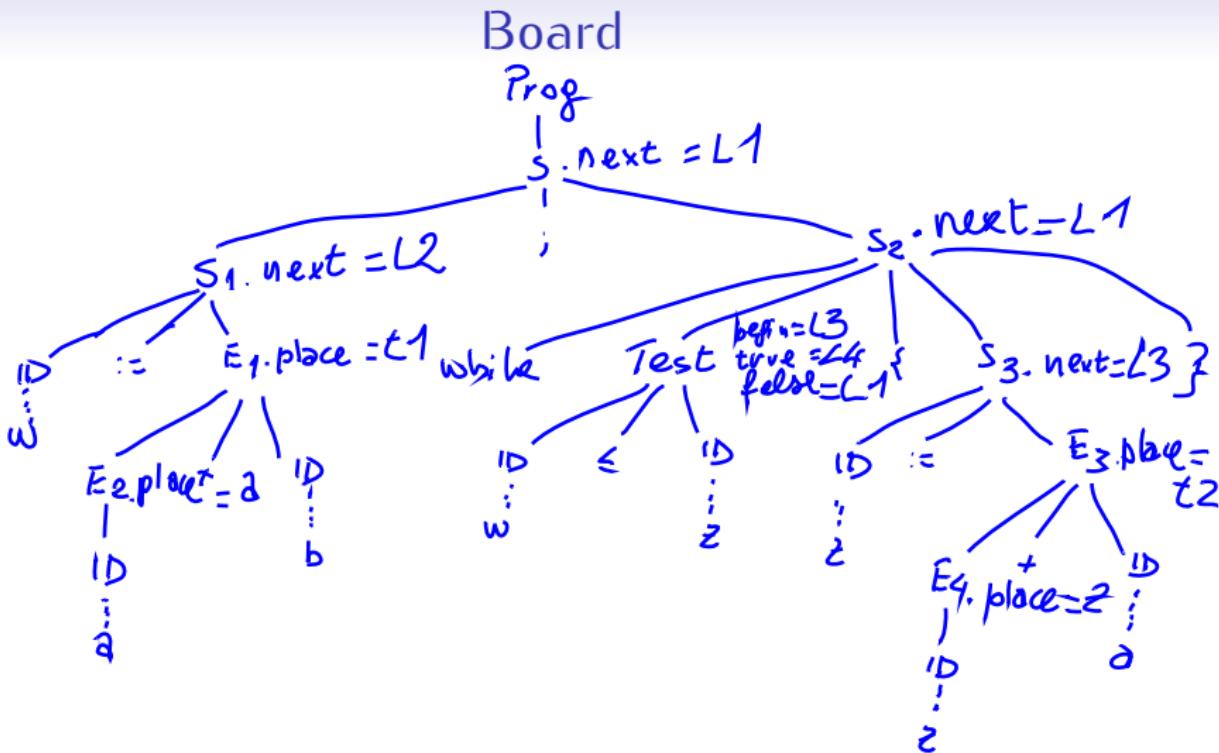
```
w := a + b;  
while w ≤ z do {  
    z := z + a}
```

Show the following:

- ① The annotated parse tree for the input together with the values of the attributes (without the *code* attribute).
- ② The three-address code produced computed for the given input.
- ③ Considering the production

$$S \rightarrow \text{while } Test \text{ do } \{S_1\}$$

show what are the synthesized attributes and what are the inherited attributes (mark with **s** the synthesized and with **h** the inherited attributes in the above semantic rules), and show the corresponding Translation Scheme.



Board

$E_2.\text{code} = ''$

$E_1.\text{code} = t1 := z + b$

$S_1.\text{code} = t1 := z + b$
 $w := t1$

$\text{Test.code} = \text{if } w \leq 2 \text{ goto L4}$
 goto L1

$E_4.\text{code} = ''$

$E_3.\text{code} = t2 := z + a$

$S_3.\text{code} = t2 := z + a$
 $z := t2$

$S_2.\text{code} =$

$L3 : \text{if } w \leq 2 \text{ goto L4}$
 goto L1

$L4 : t2 := z + a$
 $z := t2$
 goto L3

$\text{Prog.code} =$

$t1 := a + b$
 $w := t1$
 $L2:L3 : \text{if } w \leq 2 \text{ goto L4}$

$\text{L4 : } t2 := z + a$
 $z := t2$
 goto L3

$L1 :$