

Free University of Bozen-Bolzano – Faculty of Computer Science  
Bachelor in Computer Science

Formal Languages and Compilers – A.Y. 2018/19 Mid-Term Exam – 04/08/2019

Prof. Alessandro Artale – *Time: 120 minutes*

This is a closed book exam: the only resources allowed are blank paper, pens, and your head. Explain your reasoning. Write clearly, in the sense of logic, language and legibility. The clarity of your explanations affects your grade.

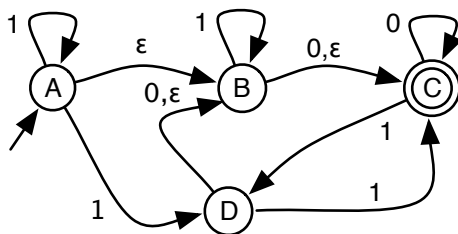
**Problem 1** [9 points] Decide which of the following statements is TRUE and which is FALSE. You must give a **clear explanation** of your answer to receive full credit.

- (a) Let  $L_1, L_2$  be any two regular languages over the same alphabet  $\Sigma$ , then the language  $L = \{w \in \Sigma^* \mid w \in L_1 \text{ or } w \notin L_2\}$  is regular.
- (b) If a language  $L$  is constituted by a **finite** set of strings, then  $L$  is a regular language.
- (c) The language  $L = \{a^n b^n \mid n \geq 1\}$  is regular (use the Pumping Lemma).

**Problem 2** [12 points]

- (a) Construct the automaton (either NFA or DFA) that recognises the language over the alphabet  $\{a, b, c\}$  constituted by all strings starting with the letter  $a$ , ending with the letter  $c$ , and never containing any of the following as a substring:  $ba, bb, ca, cc$ . E.g.,  $abcbc \in \mathcal{L}(A)$  and  $ac \in \mathcal{L}(A)$ , while  $ababc \notin \mathcal{L}(A)$  since it contains  $ba$ , and  $abbc \notin \mathcal{L}(A)$  since it contains  $bb$ . [3 POINTS]
- (b) Construct the RE for the language over the alphabet  $V = \{a, b, c\}$  where if a string contains two or more letters  $a$ , they are not adjacent. E.g.,  $bcbbbcc \in L(RE)$  since there are no  $as$ ,  $cbcacbb \in L(RE)$  since it contains a single  $a$ ,  $bcabcaba \in L(RE)$  since the  $as$  are not adjacent, while  $baac \notin L(RE)$  since the  $as$  are adjacent. [3 POINTS]
- (c) Given two DFA's over the same alphabet  $V$ , say  $A = (Q_A, V, q_0^A, \delta_A, F_A)$  and  $B = (Q_B, V, q_0^B, \delta_B, F_B)$ , formally describe the *Product Automaton*, say  $A \times B = (Q_{A \times B}, V, q_0^{A \times B}, \delta_{A \times B}, F_{A \times B})$ . In particular, say how the set of states  $Q_{A \times B}$ , the initial state  $q_0^{A \times B}$ , the transition function  $\delta_{A \times B}$ , and the set of final states  $F_{A \times B}$  can be constructed to recognise the intersection  $L(A) \cap L(B)$ . [3 POINTS]
- (d) Give the CFG for the language over the alphabet  $V = \{a, b\}$ ,  $L = \{a^n b^k \mid k > n\}$ . [3 POINTS]

**Problem 3** [5 points] Consider the following  $\epsilon$ -NFA  $A_\epsilon$  over  $\{0, 1\}$ :

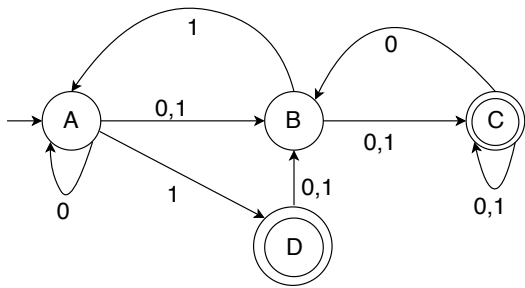


(a) Show the  $\epsilon$ -closure of each state. [2 POINTS]

(b) Construct an NFA  $A_N$  such that  $\mathcal{L}(A_N) = \mathcal{L}(A_\epsilon)$ . Just show in a table format the transition function of the resulting NFA indicating the initial and the final states. [3 POINTS]

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**Problem 4** [4 points] Consider the following NFA,  $A_N$ , over  $\{0, 1\}$ :



- (a) Construct a DFA,  $A_D$ , such that  $\mathcal{L}(A_D) = \mathcal{L}(A_N)$ . Just show in a table format the transition function of the resulting DFA indicating the initial and the final states. [3 POINTS]
- (b) Show **all** possible sequences of states of  $A_N$  that are traversed for the string 1001. [1 POINT]

**Problem 5** [5 points]

- (a) Apply the sequence of steps that are necessary to simplify a context free grammar and convert it into a Clean-up Form to the context free grammar  $G = (\{S, A, B, C, D, E\}, \{a, b, c\}, P, S)$ , where  $P$  consists of the following productions:

$$\begin{array}{ll}
 S \longrightarrow aABb \mid AE & C \longrightarrow ABC \mid c \\
 A \longrightarrow aAbE \mid \epsilon & D \longrightarrow AAB \\
 B \longrightarrow abB \mid \epsilon & E \longrightarrow CA
 \end{array}$$

The normalizations steps must be carried on in the **correct order** and the **algorithms** used for each step must become evident to get full mark.