

# Programming Paradigms

## Unit 5 — Recursion

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# Outline

- 1 Recursion
- 2 Writing Recursive Functions

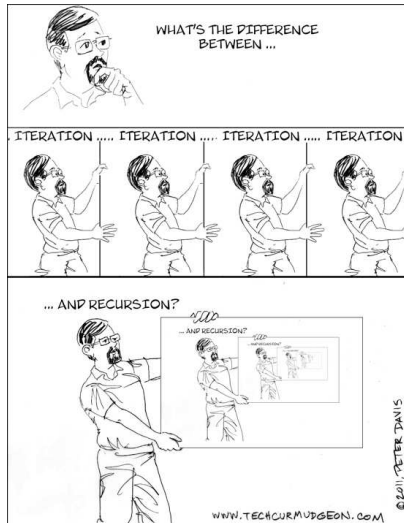
# Outline

1 Recursion

2 Writing Recursive Functions

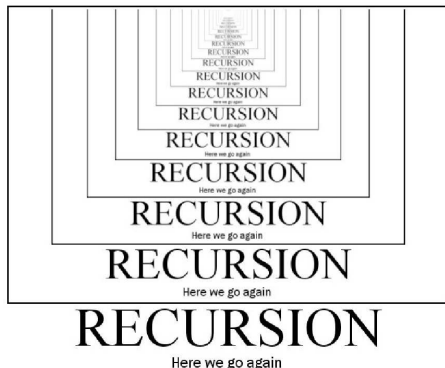
# A Different Kind of Loop/1

- The languages we are covering next have a **lack of iterative constructs**
  - That means, they have no loops
- That does not mean that they are less expressive, they use **recursion**, instead



## A Different Kind of Loop/2

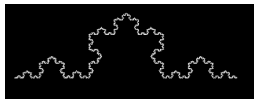
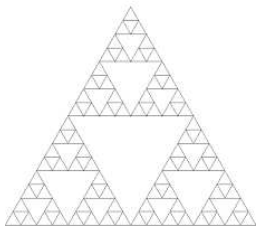
- It depends on the language how easy and efficient one or the other is
  - Some languages lack recursion: Fortran77, Assembler
  - Some languages allow recursion, but aren't very efficient with it: C++, Java
  - Languages we cover next are optimized for recursion



# Basic Idea of Recursion

- The basic idea of recursion is “*Divide et Impera*”, and is composed of three elements:
  - **Recursive case:** Divide a problem  $P$  into **subproblems** with the **same structure**, but **smaller**
  - **Base case:** At some point, the subproblem is **small enough** to solve it
  - **Composition:** Once the subproblems are solved, they can be **composed** to solve  $P$
- Many problems can be expressed very elegantly and naturally with recursion

*To iterate is human, to recurse divine (Anonymous)*



$$n! = n * (n - 1)!$$

# General Structure of a Recursive Solution/Program

```
solve( problem  $p$  )  
  if (  $p$  is simple )  
    Solve the problem  $p$  directly  
  else  
    Divide  $p$  into new sub-problems  $p_1, p_2, \dots$   
     $r_1 = \text{solve}(p_1)$   
     $r_2 = \text{solve}(p_2)$   
     $\vdots$   
    Reassemble  $r_1, r_2, \dots$  to solve  $p$   
  end  
end
```

## Example: The Handshake Problem

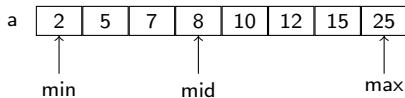
- There are  $n$  people in a room, and each person shakes hands once with every other person.
- What is the total number  $h(n)$  of handshakes?
- Recursive solutions seems very natural
  - 2 persons:  $h(2) = 1$
  - 3 persons:  $h(3) = h(2) + 2$
  - $\vdots$
  - $n$  persons:  $h(n) = h(n - 1) + (n - 1)$
- That is, the  $n$ -th person shakes  $n-1$  hands in addition to the number of handshakes of the previous  $n-1$  people
- Same as sum of  $1 + 2 + \dots + (n - 1) = \frac{n \cdot (n-1)}{2}$





## Example: Binary Search

- Binary search for an element  $v$  in a sorted array  $a$ 
  - Compare  $v$  with the middle element of  $a$
  - If not equal, apply binary search to half of  $a$  where  $v$  would be



```

bsearch(a, i, j, v)
  if ( a[mid] = v )
    return mid
  else if ( i ≥ j )
    return not found
  else if ( v ≤ a[mid] )
    return bsearch(a, min, mid-1, v)
  else
    return bsearch(a, mid+1, max, v)
  end
end

```

# Iterative Sum Example

- Let's convert a simple loop into recursion
- We're going to stay with Ruby for a while longer and write a function that computes a sum:

```
def sum(n)
  total = 0
  while(n != 0)
    total += n
    n -= 1
  end
  return total
end
```

IN ORDER TO UNDERSTAND RECURSION, ONE MUST FIRST UNDERSTAND RECURSION.

# Explaining Recursion

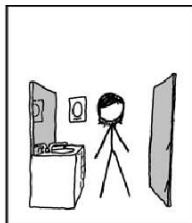
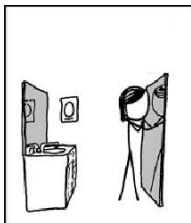
- The function from the previous slide in plain words:
  - 1 You have `n`, set `total` to 0
  - 2 If `n` is not 0 yet:
    - (a) Add `n` to `total`
    - (b) Decrement `n` by 1
    - (c) Repeat Step 2
  - 3 Done, return `total`
- Explaining Step 2 in **recursive fashion**:
  - 2 If `n` is not 0 yet, **repeat this same step** with
    - (a) `total + n` as new value for `total`
    - (b) `n-1` as new value for `n`

# Recursive Sum Example/1

- How would this look like in Ruby?

```
def sum(n,total)
  if n != 0
    sum(n-1,total+n)
  end
end
```

- Something is still missing ...



A fatal error has occurred.

0x00000539 0x4641494C  
0x4F4F5053 0x78686364

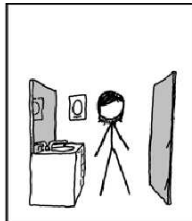
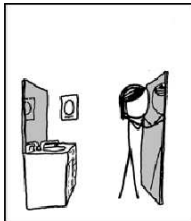
Press Ctrl+Alt+Delete to  
restart the universe.

# Recursive Sum Example/1

- How would this look like in Ruby?

```
def sum(n,total)
  if n != 0
    sum(n-1,total+n)
  end
end
```

- Something is still missing ... when and how do we **stop**?



## Recursive Sum Example/2

- When  $n$  has reached 0, we return total:

```
def sum(n,total)
  if n != 0
    sum(n-1,total+n)
  else
    total
  end
end
```

- This isn't exactly the same as the iterative version

## Recursive Sum Example/3

- To obtain the same signature as in the iterative version, we need a second function that initializes `total = 0`
  - Be careful that this cannot be done in the recursive function!

```
def do_sum(n)
  sum(n,0)
end
```

```
def sum(n,total)
  if n != 0
    sum(n-1,total+n)
  else
    total
  end
end
```

# Tail Recursion

- A very important sub-class of recursive functions are **tail recursive** functions
- This means, there is nothing left to do when the recursive call returns
- The example on the previous slide is tail recursive
- Why are these functions so important?





# Execution of Tail-Recursive Function

Recursive call	n	total	n != 0
sum(3,0)	3	0	yes
sum(2,3)	2	3	yes
sum(1,5)	1	5	yes
sum(0,6)	0	6	no
→ 6			
→ 6			
→ 6			
→ 6			

```
def sum(n,total)
  if n != 0
    sum(n-1,total+n)
  else
    total
  end
end
```

# Non-Tail Recursion

- We implemented `sum` as a tail recursive function
- It could have also been implemented in a **non-tail recursive** way:

```
def sum(n)
  if n != 0
    sum(n-1) + n
  else
    0
  end
end
```

- After returning from the recursive call we still have to add `n`

# Execution of Non-Tail-Recursive Function

Recursive call	n	n != 0
sum(3)	3	yes
sum(2) + 3	2	yes
sum(1) + 2	1	yes
sum(0) + 1	0	no
↪ 0 + 1		
→ 1		
↪ 1 + 2		
→ 3		
↪ 3 + 3		
→ 6		
→ 6		

```

def sum(n)
  if n != 0
    sum(n-1) + n
  else
    0
  end
end

```

# Tail Recursion vs. Non-Tail Recursion/1

- Non-tail recursive function calls put parameters on the stack
  - Every call grows the stack
  - On return, the parameters are needed to compute the result (together with the partial result returned)
- In tail recursive functions, the parameters from the call before are not needed anymore
  - Instead, the result is directly handed to the parent
  - Hence, no parameters need to be put on the stack
- Hence, the main difference is
  - Tail recursion composes the (partial) result before the recursive call
  - Non-tail recursion composes the (partial) result after the recursive call
- Languages that use tail recursion optimization realize this and don't grow the stack
- The languages we cover next are optimized in this way
  - So they are much more efficient when using tail recursion

# Tail Recursion vs. Non-Tail Recursion/2

- So we always use tail recursion and everything is fine?
- Unfortunately, it is not that simple:
  - **Not every recursive function** can be formulated in a tail recursive way
  - **Non-tail recursive** functions are usually easier to write: they store everything on the **stack**
  - **Tail recursive** functions have to track information and **collect partial results in accumulator** parameters, e.g. `total` in the `sum` function
- If a recursive function “**loops**” **forever**, it has to be tail recursive for obvious reasons

# Outline

1 Recursion

**2 Writing Recursive Functions**

# Writing Recursive Functions

- If you have no experience with recursive functions, writing them may seem difficult, but there are a few tricks
- Let's have a look at a concrete example: reversing an array
- First of all, it helps to **look at examples**

```
[]          -> []  
[1]         -> [1]  
[1,2]      -> [2,1]  
[1,2,3]    -> [3,2,1]
```

- This will help you get a “feel” for the problem
- You may even be able to recognize some pattern

# Base Cases

- Next, try to figure out the **base cases**
  - These are the cases that don't need a recursive call (e.g., empty list)

```
def rev(a)
  if a.length == 0 or a.length == 1
    return a
  else
    puts "not implemented yet"
  end
end
```

- You can already test this function by calling it with different parameters

```
rev([])           -> []
rev([1])         -> [1]
rev(['abc'])     -> ["abc"]
rev([1,2,3])    -> not implemented yet
rev([[1,2,3]])  -> [[1,2,3]]
```



# Recursive Cases/1

- Now, you have to consider the **recursive case**, which is a bit more difficult
- What do we have?
  - We know there are at least two elements in the array (and possibly some rest)
  - We have to add a recursive call to `rev` somewhere
- Why not imagine you already have a working version?
- Summing up, we have
  - first two elements: `a[0]` and `a[1]`
  - the rest: `a.drop(2)`  
(`drop(n)` drops the first `n` elements, here 2)
  - a working function: `old_rev`

## Recursive Cases/2

- How do we put this together?

```
def rev(a)
  if a.length == 0 or a.length == 1
    return a
  else
    old_rev(a.drop(2)).push(a[1]).push(a[0])
  end
end
```

- Basically, we reverse the rest of the array ...
- ... and append the first two elements in reverse order

## Recursive Cases/3

- This should work now
- But if it works, then it is as good as `old_rev`
  - So you can replace `old_rev` with a **recursive call** `rev` and you're done!

```
def rev(a)
  if a.length == 0 or a.length == 1
    return a
  else
    rev(a.drop(2)).push(a[1]).push(a[0])
  end
end
```

- Well, we're not quite done yet ...
  - We have to check that the recursion stops
  - We may be able to simplify the function

# Termination/1

- **Termination** is crucial in recursive functions
- For simple functions it may be easy to see it won't get stuck in an endless loop
- For more complicated ones, you can check that its arguments are **monotonically decreasing/increasing**
  - and will eventually reach one of your **base cases**

# Termination/2

- The function `rev` terminates
  - We keep dropping items from the array, making it smaller and smaller
  - Eventually it will contain only one or no item, i.e., base case
- However, checking the function `sum` we have overlooked a case
  - What happens if we call it with a negative number?



## Termination of the Sum Function

- To make the sum function always terminate, we have to check for negative numbers
- Let's change the condition to  $n > 0$

```
def sum(n,total)
  if n > 0
    sum(n-1,total+n)
  else
    return total
  end
end
```

- Alternatively, we could check for negative numbers in the initialization function `do_sum`

```
def do_sum(n)
  return sum(n,0) if n > 0
  0
end
```

# Simplification/1

- If you have **multiple base cases**, check if you actually need all of them
- If we can handle empty arrays, do we need arrays with one element as a base case?
  - The case with one element can be rewritten into:  
`[1] -> rev([]).push(1)`
  - So we only need the empty array as base case

## Simplification/2

- The simplified function looks like this:

```
def rev(a)
  return a if a == []
  rev(a.drop(1)).push(a[0])
end
```

- Was not that difficult, was it?





# Just One More Flaw/1

- We now have a recursive function that reverses an array
- However, it is not tail recursive
  - We append an element to the return value
- Can you make it tail recursive?

# Just One More Flaw/2

- We need a second parameter, which keeps the **paritally reversed array** (partial result) when going down the recursive calls

```
def rev(a,b)
  return b if a == []
  rev(a.drop(1), [a[0]] + b)
end
```

# Summary

- **Recursion** is just a different kind of loop, but as **expressive** as loops
- Some programming languages are heavily based on recursion, others do not offer recursion at all
- Three important steps in writing recursive programs
  - **Base** cases
  - **Recursive** cases
  - **Termination**
- Often recursion allows you to write elegant code
- With the right language, it is even efficient
- **Tail recursion** is important to make recursive programs efficient
  - They essentially don't need to store any data on the stack