

Verification of Data-Aware Processes

Exploiting DCDSs: models, methods, concrete systems

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Outline

- 1 The story so far
- 2 Checking/ensuring state boundedness
- 3 Boundedness and resources
- 4 Unbounded systems
- 5 Towards concrete systems

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The story so far, with main references

- The need of **combining (business) processes and data**.
[Calvanese, De Giacomo, and Montali 2013]
- A pristine formalism for data-aware business processes: **DCDS**.
[Bagheri Hariri, Calvanese, De Giacomo, et al. 2013; Montali and Calvanese 2016]
- Suitable **verification logics** for data-aware processes.
[Bagheri Hariri, Calvanese, De Giacomo, et al. 2013; Calvanese, De Giacomo, Montali, and Patrizi 2017]
- Corresponding **characterization theorems**.
[Calvanese, De Giacomo, Montali, and Patrizi 2017]
- A **decidability map**, with an unexpected dichotomy between $\mu\mathcal{L}_A$ and $LTL-FO_A$.
[Bagheri Hariri, Calvanese, De Giacomo, et al. 2013; Calvanese, De Giacomo, Montali, and Patrizi 2017]

Note: Incorrect results in [Bagheri Hariri, Calvanese, De Giacomo, et al. 2013; Okamoto 2010] fixed in [Calvanese, De Giacomo, Montali, and Patrizi 2017].

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How to check/ensure state boundedness?

Theorem

Checking whether a DCDS is **state-/run-bounded** is:

- **Decidable** for a **given** bound.
- **Undecidable** for an **unknown** bound.

Three possible strategies:

- Single out **classes of DCDSs** for which checking state-/run-boundedness is decidable.
- Identify **sufficient syntactic conditions** that are decidable to check, and that guarantee state-/run-boundedness
 - cf. syntactic conditions for chase termination in data exchange.
- Devise **modeling methodologies** that guarantee state boundedness.

DCDSs with decidable state-boundedness

Fact

DCDSs using only **unary relations** correspond to variants of **Petri nets**.

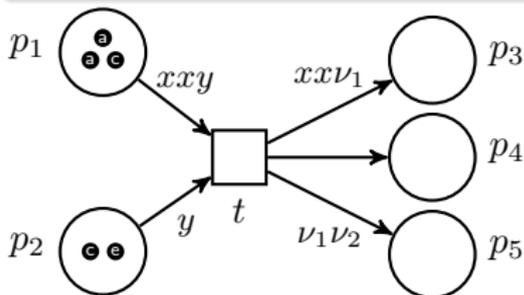
- The specific variant depends on the features used in the DCDS.

Note: **State-boundedness** relate to **boundedness** in Petri nets.

Petri nets with name management

Decidable boundedness.

[Rosa-Velardo and Frutos-Escrig 2011]



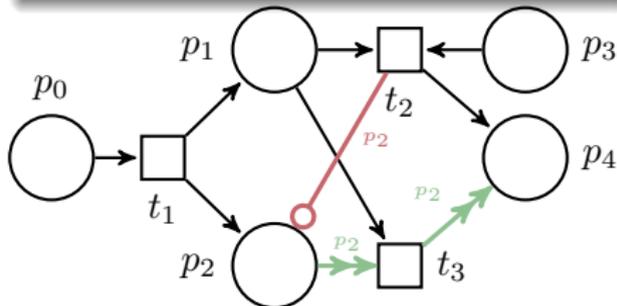
[Montali and Rivkin 2016]

Translation to DCDSs and $\mu\mathcal{L}_P$ verification.

Reset-Transfer Nets

Undecidable boundedness.

[Dufourd, Jancar, and Schnoebelen 1999]



[Bagheri Hariri, Calvanese, Deutsch, et al. 2014]

“Lossy” correspondence with DCDSs.

Attacking state-boundedness

The class of DCDSs with decidable state-boundedness **very restrictive**

These variants of Petri nets corresponds to DCDSs with only unary relations, limited use of negation, no or limited joins, ...

How to **check/guarantee** that a DCDS is state-bounded?

Sufficient, **syntactic conditions**:

- Extract a data flow graph from the DCDS.
- Check sources of unboundedness through this graph.

See [Bagheri Hariri, Calvanese, De Giacomo, et al. 2013] and [Bagheri Hariri, Calvanese, Deutsch, et al. 2014].

State-boundedness by design:

Design methods for state-bounded DCDSs. In [Solomakhin et al. 2013]:

- Processes are bound to evolving business objects (artifacts).
- Each business object manipulate boundedly many data.
- (New) business objects pick their names from a fixed pool of ids.

More sophisticated techniques in [Montali and Calvanese 2016; Calvanese, Montali, et al. 2014].

State-boundedness in concrete process modeling languages

Classical BPM languages/suites

- Central notion of **case** representing a process instance.
- Each case carries its own **case data**, in isolation to the other cases (e.g., order details, customer address, ...).
- Cases interact by accessing a central, **persistent data storage**.

Artifact-centric approaches:

- Central notion of **business object** gluing data and behaviour together.
- **All data** relevant to a business object are attached to it.
- Processes may query **multiple business objects** at once, to determine the possible next steps.

External and internal stakeholders...

- New cases/business objects are created upon **events** issued by **external stakeholders** (e.g., new order request).
- But then they are bound to **internal resources**, responsible for progressing the corresponding process instances.

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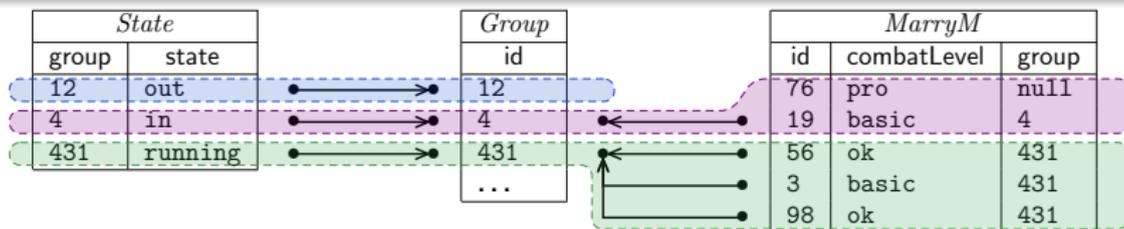
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Data isolation and case unboundedness

What if the number of simultaneously active cases cannot be bounded?

In [Montali and Calvanese 2016; Calvanese, Montali, et al. 2014], we show that **decidability** of model checking can be retained, if the system obeys to:

- **relative boundedness** (each case manipulates boundedly many data);
- **data isolation** (cases interact very weakly).



Modeling guidelines to guarantee data isolation and relative boundedness:

- 1 Queries must be navigational (no arbitrary access to relations).
- 2 1-to-many relations require a number restriction on the “many” side.
- 3 Each case cannot create a chain of tuples of unbounded length.
- 4 Cases can share tuples only in a controlled way (no construction of chains).

Beyond State-Boundedness

Question

Are there classes of DCDSs that are **unbounded**, but still **amenable to verification**?

Key result in [Abdulla et al. 2016].

Recency-bounded data-aware processes

Unbounded DB, but only the latest inserted/accessed values can bound to parameters.

Verification via under-approximation

Decidability by focusing only on runs that are k-recency-bounded for an explicitly given key.

Open problem

Investigate the relationships between all such results and those where the initial DB is **not fixed**, and verification is studied **for every possible** initial DB.

Incorporation of datatypes

Databases have **datatypes**

Numeric domains, domain-specific predicates, arithmetic.

- Many coordination algorithms and auctions require dense orders.
- Processes with costs and payment policies require integers and arithmetic.

Dense orders combine well with state-boundedness

Data-aware, state-bounded distributed systems with reals [Calvanese, Delzanno, and Montali 2015]:

- **OK** to include **dense linear orders**: minor extension to the standard DCDS abstraction technique. Intuition...

Rigid $>$ relation
over the entire domain

→

Non-rigid *GreaterThan* relation
over active domain elements.

- **No hope** to include the **successor** relation (or integers):
2 data slots are sufficient to encode two counters.

Discrete orders and arithmetic combine well with run-boundedness

Ongoing work...

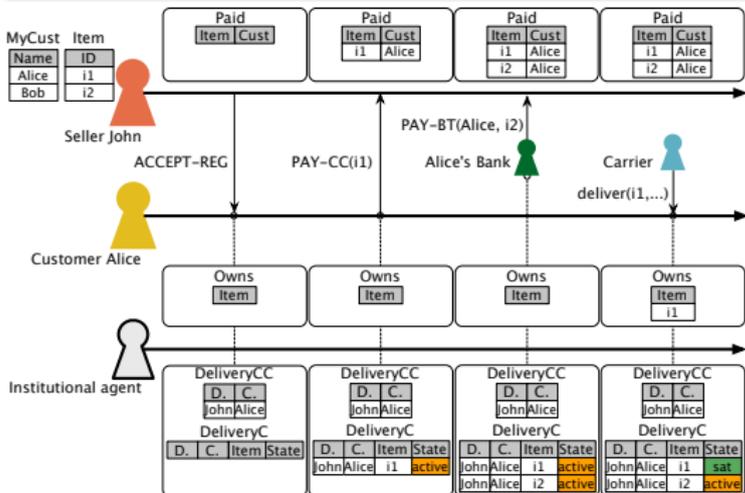
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Relational multiagent systems and commitments

Relational MAS [Montali, Calvanese, and De Giacomo 2014]

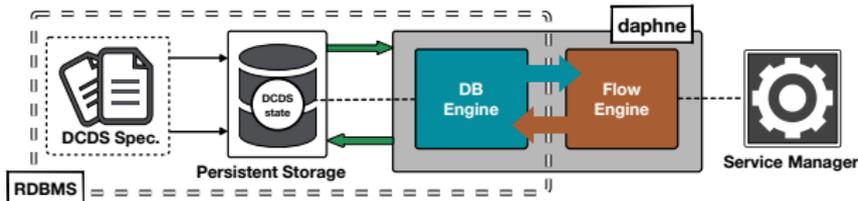
- Agents have **names** and hold/manipulate **local, state-bounded DBs**.
- Agents **exchange data** using their names for addressing.
- An **institutional agent** manages agent creation and deletion.
 - Due to state-boundedness: **unboundedly** many agents can dynamically enter into the system, but at each moment only **boundedly many** are active.



Relational commitments

In the same work: first proposal for modeling and verifying **interaction protocols** based on **relational commitments**, i.e., commitments with data payload and multiple instances.

daphne: implementing DCDSs with relational technology



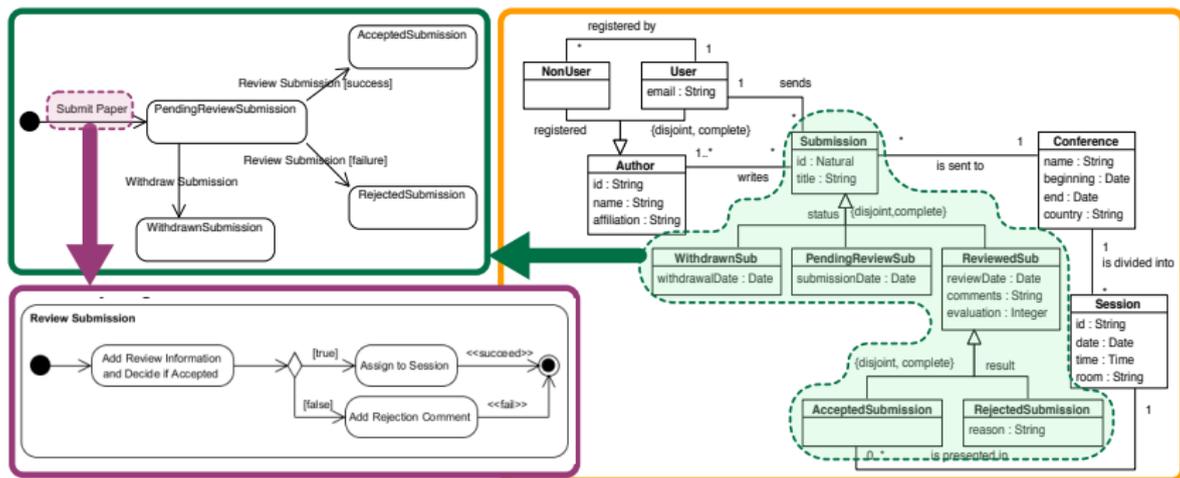
Native modeling and execution of DCDSs using relational DBMSs:

- **SQL-like syntax** for DCDSs with datatypes.
- Automated translation into **relational DBMSs**, as (temporal) **tables**, **constraints**, and **stored procedures**.
- **Java APIs** to support **enactment** and integration with concrete **services**.

Native explicit model checking of DCDSs using relational DBMSs:

- **Same model for execution and verification!**
- **Special tables** for **storing the RTS** induced by a DCDSs.
- **Factoring of tables** into **temporal** and **atemporal** parts.
- **Computation** of **equality commitments** and **value recycling** in services.
- **Java APIs** for **RTS construction and search**.

BAUML: artifact-centric processes with UML

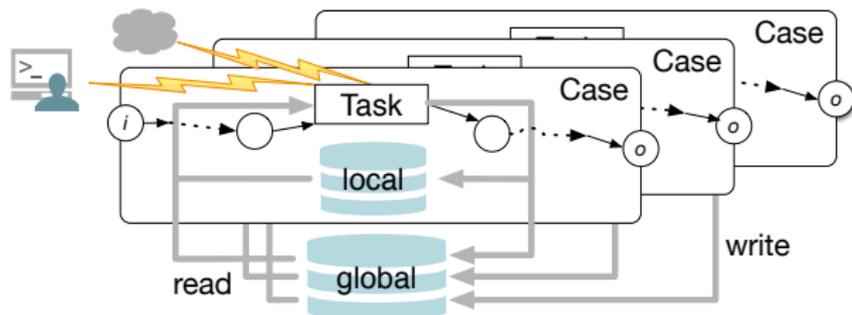


BAUML approach

- **Business objects**, states, associations and attributes: **UML class diagrams**.
- Business object **lifecycle**: **UML statechart diagram**.
- **Complex event** triggering a lifecycle transition: **UML activity diagram**.
 - **Tasks** modeled as **OCL operation contracts**.

In [Calvanese, Montali, et al. 2014]: **methodology to guarantee decidability of model checking** (see before). Estanol PhD thesis: BAUML to DCDS!

raw-sys: marrying workflow nets and databases



raw-sys **model** [De Masellis et al. 2017]:

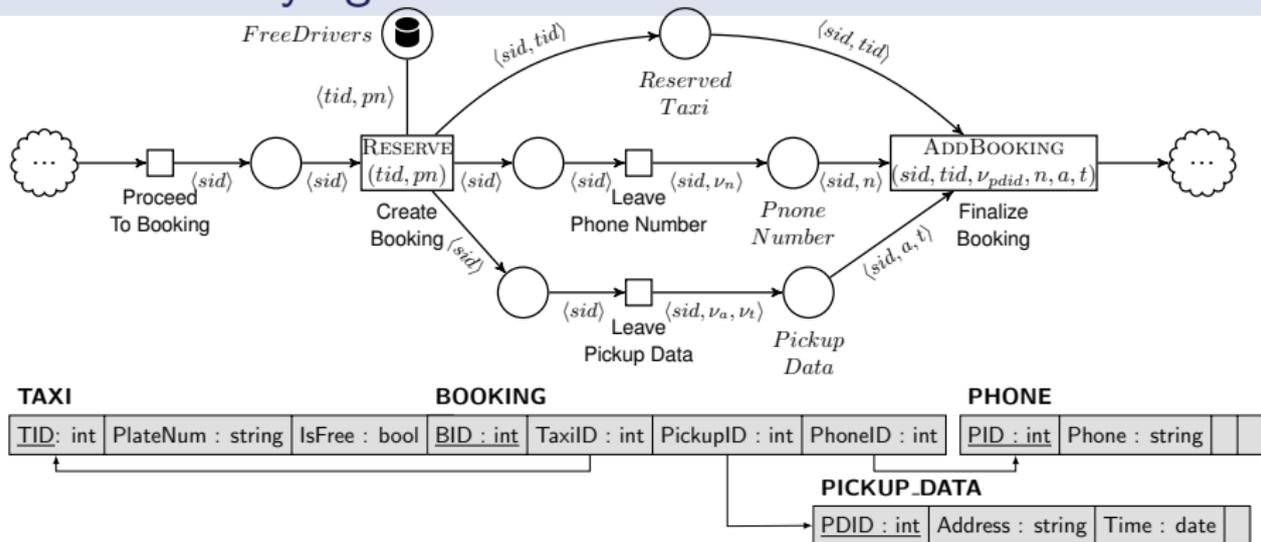
Data-aware processes using **well-known formalisms**:

- **Data**: global and local **relational databases**.
- **Process control-flow**: **workflow nets**, enriched with:
 - **Guards** (queries over the DBs).
 - **STRIPS-like actions** with external inputs from an infinite domain, invoked upon firing net transitions.

raw-sys **verification** [De Masellis et al. 2017]:

- Map of **(un)decidability**, exploiting **translation to DCDSs**.
- **Encoding into planning** systems to handle **reachability problems**.

db-nets: marrying colored Petri nets and databases

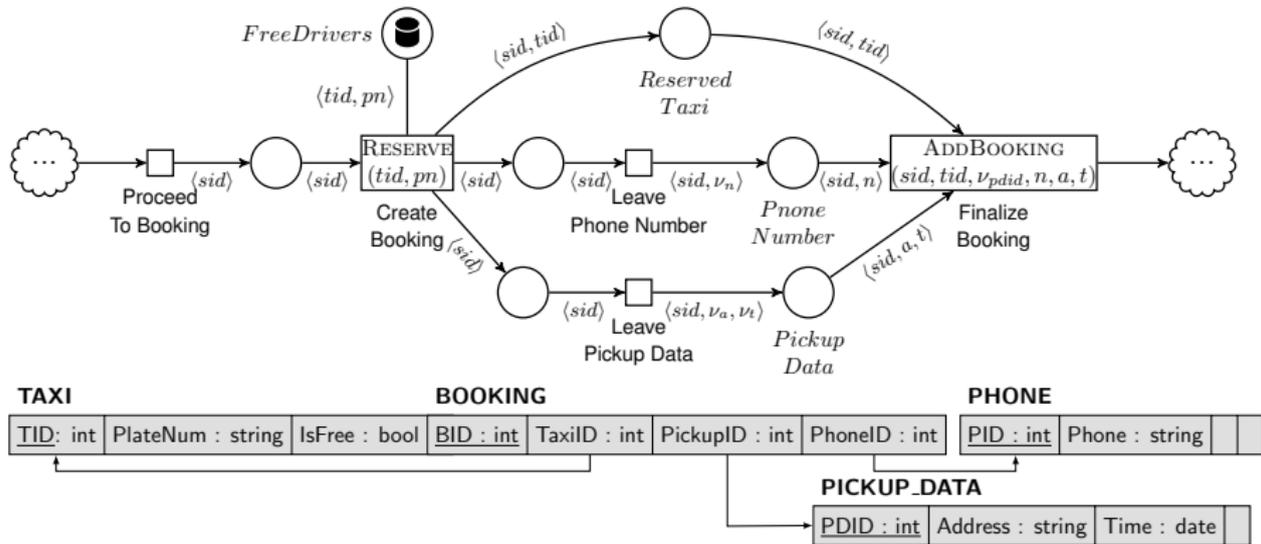


db-net model [Montali and Rivkin 2017], three layers:

- 1 **Persistence:** relational database with constraints.
- 2 **Data logic:** queries and actions over the persistence layer.
- 3 **Control:** colored Petri net with ν -variables, enriched with **view places** and **transition-action bindings** to inspect/update the persistence layer.

Note: Natural formalization of contemporary process modeling suites!

db-nets: marrying colored Petri nets and databases



db-nets execution, simulation, verification [Montali and Rivkin 2017]:

- Foundational results thanks to **translation to DCDSs**.
- Ongoing implementation effort inside www.cpntools.org.

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